## CONGRUENCES ON FREE MONOIDS AND SUBMONOIDS OF POLYCYCLIC MONOIDS

#### JOHN MEAKIN and MARK SAPIR

(Received 20 August 1990)

Communicated by P. G. Trotter

#### Abstract

We establish a one-to-one "group-like" correspondence between congruences on a free monoid  $X^*$  and so-called positively self-conjugate inverse submonoids of the polycyclic monoid P(X). This enables us to translate many concepts in semigroup theory into the language of inverse semigroups.

1991 Mathematics subject classification (Amer. Math. Soc.): 20 M 8, 20 M 10.

### 1. Introduction

For basic background on semigroups we refer the reader to Clifford and Preston [2] or Lyapin [3]: standard information about inverse semigroups may be found in Petrich [8]. Of particular concern to us in the present paper is the *polycyclic monoid* P(X) on a set X, introduced into the literature by Nivat and Perrot [7]. The monoid P(X) may be defined as the inverse hull of the free monoid  $X^*$  on the non-empty set X: it arises naturally as the syntactic monoid of the restricted Dyck language on a set of cardinality |X| and also in connection with the operation of a push-down automaton, which may be realized by a rational transducer with values in P(X). We refer the reader to the papers [6, 7, 9], for further details and applications of the polycyclic monoid in formal language theory. For the reader's convenience, we briefly

The research of the first author was partially supported by NSF grant DMS-8702019. (c) 1993 Australian Mathematical Society 0263-6115/93 \$A2.00 + 0.00

indicate below a few elementary facts about the polycyclic monoid which we shall use in the present paper.

If X is a non-empty set and  $X^{-1}$  a disjoint set in one-one correspondence with X (by means of the mapping  $x \to x^{-1}$ ,  $x \in X$ ) then P(X) may be viewed as the monoid presented by the set  $X \cup X^{-1}$  of generators and relations of the form  $xx^{-1} = 1$ ,  $xy^{-1} = 0$ ,  $x \neq y$ ,  $x, y \in X$ . If |X| = 1, then P(X) is just the bicyclic monoid (Clifford and Preston [2]). As usual, with every word  $w = x_1 \dots x_n \in (X \cup X^{-1})^*$  we associate a word  $w^{-1} = x_n^{-1} \dots x_1^{-1}$ . From the presentation of P(X) it follows that each element of P(X) may be written uniquely in the form  $u^{-1}v$  for some words  $u, v \in X^*$ . The multiplication in P(X) is described as follows:

$$(u_1^{-1}v_1)(u_2^{-1}v_2) = \begin{cases} u_1^{-1}wv_2 & \text{if } v_1 = wu_2, \text{ some } w \in X^* \\ u_1^{-1}w^{-1}v_2 & \text{if } u_2 = wv_1, \text{ some } w \in X^* \\ 0 & \text{otherwise} \end{cases}$$

Idempotents of P(X) are of the form  $u^{-1}u$  ( $u \in X^*$ ) and the semilattice of idempotents of P(X) is just the usual partially ordered set of  $X^*$  (with a 0 adjoined if |X| > 1).

Of course P(X) is an inverse monoid: we show in this paper that every monoid is coded by an inverse submonoid of an appropriate polycyclic monoid.

### 2. The semigroups $P_X$ , $Q_X$ , $I_X$

Let X be a finite non-empty set and let  $X^{Z}$  be the set of doubly infinite words (sequences) over X. An element  $\underline{x}$  of  $X^{Z}$  will usually be denoted by  $\underline{x} = (\dots, x_{-2}, x_{-1}, x_0, x_1, x_2, \dots)$  or  $\underline{x} = (x_n)_{n \in \mathbb{Z}}$ : here  $x_n$  denotes the entry in position  $n(n \in \mathbb{Z})$  of the doubly infinite word  $\underline{x}$  and  $\mathbb{Z}$  denotes the set of integers. It may be convenient for the reader to view the word  $\underline{x}$  as being printed on a doubly infinite tape and to consider position 0 of the tape as being located by means of a fixed reading head that scans  $x_0$ . Let Y = Xif |X| > 1 and  $Y = X \cup \{t\}$  where  $t \notin X$  if |X| = 1. We define |X| + 1(partial) mappings from  $Y^{Z}$  to  $Y^{Z}$ .

The first mapping is the shift T that shifts each sequence one space to the right. More precisely, if  $\underline{x} = (x_n)_{n \in \mathbb{Z}}$ , then T maps  $\underline{x}$  to  $\underline{y} = (y_n)_{n \in \mathbb{Z}}$  where  $y_n = x_{n-1}$  (for  $n \in \mathbb{Z}$ ). Note that if we view  $\underline{x}$  as being printed on a doubly infinite tape with fixed reading head then T has the effect of shifting the tape one space to the right, so that the symbol scanned by the reading head becomes  $x_{-1}$ , instead of  $x_0$ . Clearly T is a permutation of  $Y^{\mathbb{Z}}$  (with

domain all of  $Y^{Z}$  and inverse  $T^{-1}$ , the map that shifts each sequence one space to the left).

For each letter  $x \in X$  we denote by  $\alpha_x$  the mapping that deletes  $x_1$  from  $(x_n)_{n \in \mathbb{Z}}$  if  $x_1 = x$ . Thus  $\alpha_x$  is a partial mapping from  $Y^{\mathbb{Z}}$  to  $Y^{\mathbb{Z}}$  with domain being set of all sequences  $(x_n)_{n \in \mathbb{Z}}$  for which  $x_1 = x$ : the mapping  $\alpha_x$  maps  $(x_n)_{n \in \mathbb{Z}}$  to  $(y_n)_{n \in \mathbb{Z}}$  where

$$y_n = \begin{cases} x_n & \text{if } n \le 0\\ x_{n+1} & \text{if } n \ge 1 \end{cases}$$

It is clear that  $\alpha_x(x \in X)$  is a partial one-one mapping of  $Y^Z$  to  $Y^Z$  with inverse map  $\alpha_x^{-1}$ , the map that inserts the letter x in  $(x_n)_{n \in \mathbb{Z}}$  between  $x_0$  and  $x_1$ , thus creating the new sequence  $(z_n)_{n \in \mathbb{Z}}$  with

$$z_n = \begin{cases} x_n & \text{for } n \le 0\\ x & \text{for } n = 1\\ x_{n-1} & \text{for } n > 1 \,. \end{cases}$$

Since T and  $\alpha_x (x \in X)$  are partial one-one maps of  $Y^Z$  to  $Y^Z$ , we may view them as elements of the symmetric inverse monoid SIM( $Y^Z$ ) on the set  $Y^Z$  (see Clifford and Preston [2] or Petrich [8] for background on the symmetric inverse monoid). We now define three submonoids  $P_X$ ,  $I_X$  and  $Q_X$  of SIM( $Y^Z$ ). We use the notation  $\langle U \rangle$  to denote the inverse submonoid of SIM( $Y^Z$ ) generated by the subset  $U \subseteq$  SIM( $Y^Z$ ): define

$$P_{\mathbf{X}} = \langle \{\alpha_x : x \in \mathbf{X}\} \rangle,$$
  
$$I_{\mathbf{X}} = \langle \{\alpha_x : x \in \mathbf{X}\} \cup \{\mathbf{T}\} \rangle,$$

and let  $Q_X$  be the set of elements of  $I_X$  that can be written as a product of (positive and negative) powers of T and the  $\alpha_x (x \in X)$  for which the sum of powers of T involved is zero.

**REMARK.** We shall see later that no element of  $Q_X$  can be written as a product of powers of T and the  $\alpha_x (x \in X)$  in such a way that the sum of powers of T involved is non-zero.

Note that  $I_X$  is a finitely generated inverse monoid and that  $Q_X$  is a finitely generated algebra of type  $\langle \cdot, -1 \rangle$ , T where T is an isomorphism  $a \mapsto T^{-1} a T$ .

For every word  $u \in X^*$ , it is convenient to fix the notation  $\alpha_u$  for the mapping from  $Y^Z$  to  $Y^Z$  that deletes the word u from a sequence  $(x_n)_{n \in \mathbb{Z}}$  if  $u = x_1 x_2 \dots x_k$ . Clearly  $\alpha_u$  is a partial one-one map of  $Y^Z$  to  $Y^Z$  with domain the set of sequences  $(x_n)_{n \in \mathbb{Z}}$  for which  $u = x_1 x_2 \dots x_k$ ; also  $\alpha_u$  is the product  $\alpha_u = \alpha_{x_1} \alpha_{x_2} \dots \alpha_{x_k} \in P_X$ . We establish a number of basic properties and facts concerning the monoids  $P_X$ ,  $I_X$ , and  $Q_X$ .

# **LEMMA 2.1.** The monoid $P_X$ is isomorphic to the polycyclic monoid P(X).

**PROOF.** Note that for every  $x, y \in X$ ,  $\alpha_x^{-1}\alpha_x = 1$  (identity map of SIM(Y<sup>Z</sup>)) and  $\alpha_x^{-1}\alpha_y = 0$  if  $x \neq y$ . It follows that every element of  $P_X$  is expressible in the form  $\alpha_u \alpha_v^{-1}$  for some  $u, v \in X^*$ . The element  $\alpha_u \alpha_v^{-1}$  of  $P_X$  has as domain the set of all sequences  $(x_n)_{n \in \mathbb{Z}}$  with  $x_1 \dots x_k = u$  (some k) and has the effect of replacing the subword  $x_1 \dots x_k = u$  of such a sequence by the subword v. It follows that if  $u, v, s, t \in X^*$  then  $\alpha_u \alpha_v^{-1} = \alpha_s \alpha_t^{-1}$  iff u = s and v = t. Hence every element of  $P_X$  is uniquely expressible in the form  $\alpha_u \alpha_v^{-1}$  for some  $u, v, \in X^*$ . From the description of the polycyclic monoid P(X) given in section 1 it is now clear that the map  $u^{-1}v \mapsto \alpha_{r(u)} \alpha_{r(u)}^{-1}$  is an isomorphism from P(X) onto  $P_X$ . (Here r(u) denotes the reverse of the word u: that is, if  $u = x_1 x_2 \dots x_k$  then  $r(u) = x_k \dots x_2 x_1$ .)

Denote the set of idempotents of  $I_X$ ,  $Q_X$  and  $P_X$  by  $EI_X$ ,  $EQ_X$  and  $EP_X$  respectively. Note that  $EI_X$  and  $EP_X$  are semilattices and  $EQ_X$  is a semilattice and subalgebra of  $Q_X$  (since every automorphism preserves idempotents).

**LEMMA 2.2.** (a) The semilattice  $EP_X$  consists of zero and all elements of the form  $\alpha_u \alpha_u^{-1}$  for some  $u \in X^*$ ; (b) the semilattice  $EI_X$  is generated by the set  $\{T^{-k} \alpha_x \alpha_x^{-1} T^k : x \in X, k \in \mathbb{Z}\} \cup \{1\}$ ; (c) the subalgebra  $EQ_X$  coincides with  $EI_X$  and is generated by the set  $\{\alpha_x \alpha_x^{-1} : x \in X\} \cup \{1\}$ .

PROOF. Part (a) is obvious from the description given in section 1 of the semilattice of idempotents of the polycyclic monoid. To prove part (b), suppose first that  $\alpha$  is any element of  $I_X$  with  $\alpha \neq 0$  and  $\alpha \neq 1$ . Then  $\alpha$  must be expressible as a finite product of elements of the form T,  $T^{-1}$ ,  $\alpha_x$  and  $\alpha_x^{-1}$  for  $x \in X$ . Now T,  $T^{-1}$  and  $\alpha_x^{-1}$  all have domain  $X^Z$  and  $\alpha_x$  has domain consisting of all sequences  $(x_n)_{n \in \mathbb{Z}}$  with  $x_1 = x$ . It follows that there must exist integers  $i_1 < i_2 < \ldots < i_k \in \mathbb{Z}$  and fixed elements  $a_1, a_2, \ldots, a_k \in X$  such that  $(x_n)_{n \in \mathbb{Z}}$  is in the domain of  $\alpha$  if and only if  $x_{i_i} = a_i$  for  $1 \le t \le k$ . Now if  $\alpha$  is an idempotent of  $I_X$  then  $\alpha$  must be the restriction of the identity map on  $X^Z$  to its domain, so  $\alpha = \beta_{i_1}\beta_{i_2}\ldots\beta_{i_k}$  where  $\beta_{i_i}$  is the identity map on the set of sequences  $(x_n)_{n \in \mathbb{Z}}$  with  $x_{i_i} = a_i$  (for  $1 \le t \le k$ ). Since  $\beta_{i_i} = T^{-i_i+1}\alpha_{a_i}\alpha_{a_i}^{-1}T^{i_i-1}$  the result of part (b) follows. From part (b) it follows that  $EI_X \subseteq EQ_X$ : since the converse is

obvious it follows immediately that EQ<sub>X</sub> is generated (as an algebra) by  $\{\alpha_x \alpha_x^{-1} : x \in X\} \cup \{1\}$ .

LEMMA 2.3. For all words  $u, v, s, t \in X^*$ , all  $k \in \mathbb{Z}$  and all integers  $l \ge 0$  we have

(1) 
$$\alpha_u \operatorname{T}^{-l} \alpha_s \alpha_t^{-1} \operatorname{T}^{l} = \operatorname{T}^{-(l+|u|)} \alpha_s \alpha_t^{-1} \operatorname{T}^{(1+|u|)} \alpha_u$$

and hence

(2) 
$$(\mathbf{T}^{-k} \alpha_{u} \alpha_{v}^{-1} \mathbf{T}^{k}) (\mathbf{T})^{(k+|v|+l)} \alpha_{s} \alpha_{t}^{-1} \mathbf{T}^{(k+|v|+l)})$$
$$= (\mathbf{T}^{-(k+|u|+l)} \alpha_{s} \alpha_{t}^{-1} \mathbf{T}^{(k+|u|+l)}) (\mathbf{T}^{-k} \alpha_{u} \alpha_{v}^{-1} \mathbf{T}^{k}).$$

In particular,

(3) 
$$\alpha_{u}\alpha_{s}\alpha_{t}^{-1} = \mathbf{T}^{-|u|}\alpha_{s}\alpha_{t}^{-1}\mathbf{T}^{|u|}\alpha_{u}.$$

**PROOF.** Let  $\alpha$  (respectively  $\beta$ ) denote the mapping on the left side (respectively right side) of equation (1). A sequence  $(x_n)_{n \in \mathbb{Z}}$  is in the domain of  $\alpha$  if and only if  $x_1 \dots x_{|u|} = u$  and  $x_{|u|+l+1} \dots x_{|u|+l+|s|} = s$ : similarly one checks that  $\beta$  has this domain. The effect of  $\alpha$  when applied to such a sequence is to first erase the segment  $x_1 \dots x_{|u|}$  and then replace the segment  $x_{|u|+l+1} \dots x_{|u|+l+|s|}$  by t, while  $\beta$  had the same effect, but in the opposite order. Hence  $\alpha = \beta$ . Clearly (3) follows from (1) with l = 0. To see that (2) also follows from (1) note first that, on replacing u by v and s by t in (1) and taking inverses, we have

$$\alpha_v^{-1} \mathbf{T}^{-(l+|v|)} \alpha_s \alpha_t \mathbf{T}^{(l+|v|)} = \mathbf{T}^{-l} \alpha_s \alpha_t^{-1} \mathbf{T}^{l} \alpha_v^{-1},$$

SO

$$\begin{aligned} \alpha_{u} \alpha_{v}^{-1} \operatorname{T}^{-(l+|v|)} \alpha_{s} \alpha_{t}^{-1} \operatorname{T}^{(l+|v|)} &= \alpha_{u} \operatorname{T}^{-l} \alpha_{s} \alpha_{t}^{-1} \operatorname{T}^{l} \alpha_{v}^{-1} \\ &= \operatorname{T}^{-(l+|u|)} \alpha_{s} \alpha_{t}^{-1} \operatorname{T}^{(l+|u|)} \alpha_{u} \alpha_{v}^{-1}, \quad \text{by (1).} \end{aligned}$$

Equation (2) follows immediately from this by premultiplying by  $T^{-k}$ , postmultiplying by  $T^{k}$  and noting that  $T^{k}T^{-k} = 1$ .

LEMMA 2.4. For all words  $u, v, w \in X^*$  with  $|w| \ge 1$  and for every positive integer  $k \le |w|$  we have

(4) 
$$\mathbf{T}^{-k} \alpha_u \alpha_v^{-1} \mathbf{T}^k \alpha_w = \alpha_{w(k)u} \alpha_{w(k)v}^{-1} \alpha_w;$$

(5) 
$$\alpha_w^{-1} \operatorname{T}^{-k} \alpha_u \alpha_v^{-1} \operatorname{T}^{k} = \alpha_w^{-1} \alpha_{w(k)u} \alpha_{w(k)v}^{-1}$$

where w(k) is the prefix of w of length k.

**PROOF.** Write w in the form w = w(k)s for some  $s \in X^*$ . By Lemma 2.3, (3) we have

$$\mathbf{T}^{-k} \alpha_{u} \alpha_{v}^{-1} \mathbf{T}^{k} \alpha_{w(k)} = \alpha_{w(k)} \alpha_{u} \alpha_{v}^{-1}$$
$$= \alpha_{w(k)u} \alpha_{v}^{-1}.$$

Hence

$$\mathbf{T}^{-k} \alpha_{u} \alpha_{v}^{-1} \mathbf{T}^{k} \alpha_{w} = \mathbf{T}^{-k} \alpha_{u} \alpha_{v}^{-1} \mathbf{T}^{k} \alpha_{w(k)} \alpha_{s}$$
$$= \alpha_{w(k)u} \alpha_{v}^{-1} \alpha_{s}$$
$$= \alpha_{w(k)u} \alpha_{v}^{-1} \alpha_{w(k)}^{-1} \alpha_{w(k)} \alpha_{s}$$
$$= \alpha_{w(k)u} \alpha_{w(k)v}^{-1} \alpha_{w},$$

so equality (4) is established. Of course equality (5) follows from equality (4) by interchanging u and v and taking inverses of both sides of the resulting equality.

If  $p = \alpha_u \alpha_v^{-1} \in P_X$  for some words  $u, v \in X^*$ , it is convenient to denote |v| by |p| and |u| - |v| by d(p). We also use the notation A + B to denote the join of A and B in the lattice of inverse submonoids of  $I_X$ . If  $w = \dots x_{-1}x_0x_1\dots$  is an infinite word and k, l are integers,  $k \leq l$ , then w[k, l] will denote the word  $x_k \dots x_l$ .

**LEMMA 2.5.** (a) Every element  $q \in Q_X$  may be uniquely represented in the form

(6) 
$$q = \mathbf{T}^{-k_1} p_1 \mathbf{T}^{k_1} \mathbf{T}^{-k_2} p_2 \mathbf{T}^{k_2} \dots \mathbf{T}^{-k_n} p_n \mathbf{T}^{k_n}$$

for some integers  $k_1, k_2, ..., k_n$  and elements  $p_1, ..., p_n \in P_X \setminus \{1\}$  with  $k_2 \ge k_1 + |p_1|, k_3 \ge k_2 + |p_2|, ..., k_n \ge k_{n-1} + |p_{n-1}|;$ 

(b) if an element q is represented in the form (6) then the domain of q consists of all infinite words w such that

$$w[k_1 + m_1, k_1 + m_1 + |u_1|] = u_1, \quad w[k_2 + m_2, k_2 + m_2 + |u_2|] = u_2, \dots,$$
  
$$w[k_n + m_n, k_n + m_n + |u_n|] = u_n,$$

where  $m_n$  may be defined inductively by:

(7) 
$$m_1 = 1, ..., m_i = m_{i-1} + d(p_{i-1}), \quad i > 1,$$

and the effect of applying q to such a word is to replace each segment  $w[k_i + m_i, k_i + m_i + |u_i|]$  by  $v_i$  for i = 1, ..., n.

(c) if  $p_i = \alpha_{u_i} \alpha_{v_i}^{-1}$  for some words  $u_i$ ,  $v_i$  (i = 1, ..., k), then  $(\langle p_i : i = 1, 2, ..., k \rangle + \langle T \rangle) \cap Q_X$  consists of elements of the type

(8) 
$$\mathbf{T}^{-k_1} q_1 \mathbf{T}^{k_1} \mathbf{T}^{-k_2} q_2 \mathbf{T}^{k_2} \dots \mathbf{T}^{-k_n} q_n \mathbf{T}^{k_n}$$

where each  $q_i$  is a product of elements of the form  $\alpha_{wu}\alpha_{wv}^{-1}$  or  $\alpha_{wv}\alpha_{wu}^{-1}$  for some words  $u, v, w \in X^*$  with  $(u, v) \in \{(u_j, v_j) : j = 1, ..., k\}$  and some integers  $k_i$  with  $k_{i+1} > k_i + |q_i|$  for i = 1, ..., n-1.

PROOF. Using the definition of  $Q_X$  and the fact that  $TT^{-1} = T^{-1}T = 1$ , it is easy to see that every element  $q \in Q_X$  may be written in the form  $q = T^{-k_1}p_1T^{k_1}T^{-k_2}p_2T^{k_2}\dots T^{-k_n}p_nT^{k_n}$  for some integers  $k_i$  and elements  $p_i$  of  $P_X$ . Successive application of equalities (2) of Lemma 2.3 and (4) of Lemma 2.4 then enables us to rewrite the product for q above in such a way that we may assume that  $k_1 < k_2 < \dots < k_n$ . Now write  $p_i = \alpha_{u_i} \alpha_{v_i}^{-1}$  for some words  $u_i, v_i \in X^*$  and let  $l_i = k_{i+1} - k_i$  (so that each  $l_i > 0$ ). Thus

$$q = \mathbf{T}^{-k_1} \alpha_{u_1} \alpha_{v_1}^{-1} \mathbf{T}^{-l_1} \alpha_{u_2} \alpha_{v_2}^{-1} \mathbf{T}^{-l_2} \alpha_{u_3} \alpha_{v_3}^{-1} \dots \mathbf{T}^{-l_{n-1}} \alpha_{u_n} \alpha_{v_n}^{-1} \mathbf{T}^{k_n}$$

If  $l_1 \le |v_1| = |p_1|$  we may rewrite this as

$$q = \mathbf{T}^{-k_1} \alpha_{u_1} \alpha_{v_1}^{-1} \mathbf{T}^{-l_1} \alpha_{u_2} \alpha_{v_2}^{-1} \mathbf{T}^{-(l_1+l_2)} \alpha_{u_3} \alpha_{v_3}^{-1} \dots \alpha_{u_n} \alpha_{v_n}^{-1} \mathbf{T}^{k_n}$$

and we may apply equality (5) (Lemma 2.4) to obtain

$$q = \mathbf{T}^{-k_1} \alpha_{u_1} \alpha_{v_1}^{-1} \alpha_{v_1(l_1)u_2} \alpha_{v_1(l_1)v_2}^{-1} \mathbf{T}^{-(l_1+l_2)} \alpha_{u_3} \alpha_{v_3}^{-1} \dots \alpha_{u_n} \alpha_{v_n}^{-1} \mathbf{T}^{k_n}.$$

Similar reductions occur if  $l_i \le |v_i| = |p_i|$  (for 1 < i < n). Then by induction on *n* we obtain (6). Equality (8) also follows from the above proof.

Condition (b) is an immediate consequence of the definition of the semigroup  $I_X$ . Notice that the replacements that are described in condition (b) are independent. Indeed by (6) we have

$$(k_{i+1} + m_{i+1}) - (k_i + m_i + |u_i|) = k_{i+1} - k_i + m_{i+1} - m_i - |u_i|$$
  
=  $k_{i+1} - k_i - |v_i| > 0.$ 

This means that each replacement does not change the segment we obtain as a result of any other replacement. This gives us the uniqueness of the representation (6).

The preceding results enable us to give a presentation for the semigroup  $I_X$ . We denote by Inv(X : R) the *inverse* monoid presented by a set X of generators and a set R of relations. Thus we may think of R as a subset of  $(X \cup X^{-1})^* \times (X \cup X^{-1})^*$  where  $X^{-1}$  is a set in one-one correspondence

with X (and disjoint from X) and  $Inv\langle X : R \rangle = (X \cup X^{-1})^* / \tau$ , where  $\tau$  is the congruence on  $(X \cup X^{-1})^*$  generated by  $\rho \cup R$ . (Here  $\rho$  is the Vagner congruence on  $(X \cup X^{-1})^*$  – see Petrich [8].) Presentations of inverse monoids have received considerable recent attention in the literature - see for example the papers of Stephen [10] and Margolis and Meakin [4].

**THEOREM 2.6.** (a) The semigroup  $P_X$  admits a presentation of the form  $P_X = \text{Inv}\langle \tilde{X} : R_1 \rangle$  where  $\tilde{X} = \{\alpha_x : x \in X\}$  and  $R_1$  consists of the relations  $\alpha_x^{-1}\alpha_x = 1$ ,  $\alpha_x^{-1}\alpha_y = 0$  for  $x, y \in X$ ,  $x \neq y$ .

(b) The semigroup  $I_X$  is a semidirect product of  $Q_X$  and Z (the group of integers).

(c) The semigroup  $I_X$  admits a presentation of the form  $I_X = Inv\langle \overline{X} \cup \{T\}$ :  $R_2\rangle$  where  $R_2 = R_1 \cup R_T \cup R_C$ ,  $R_T = \{(TT^{-1}, 1), (T^{-1}T, 1)\}$  and  $R_C$  consists of all relations listed in equation (1), (Lemma 2.3).

**PROOF.** Part (a) was already established in Lemma 2.1. To prove part (b) note first that application of the relations  $TT^{-1} = T^{-1}T = 1$  easily enables us to represent every element  $y \in I_X$  in the form  $y = qT^k$  for some  $q \in Q_X$ ,  $k \in \mathbb{Z}$ . If  $q_1 T^{k_1} = q_2 T^{k_2}$  for some  $q_1$ ,  $q_2 \in Q_X$  and  $k_1$ ,  $k_2 \in \mathbb{Z}$  then  $q_1 T^{k_1-k_2} = q_2$ . We prove that  $q_1 = q_2$  and  $k_1 = k_2$ . This will in particular provide a justification for the remark made immediately after the definition of  $Q_X$ , namely that an element y of  $I_X$  is in  $Q_X$  only if the sum of powers of T involved in the expression for y, as a product of powers of T and the  $\alpha_x(x \in X)$ , is zero. Represent  $q_1$  and  $q_2$  in the forms

$$q_1 = \mathbf{T}^{-k_1} \alpha_{u_1} \alpha_{v_1}^{-1} \mathbf{T}^{k_1} \dots \mathbf{T}^{-k_n} \alpha_{u_n} \alpha_{v_n}^{-1} \mathbf{T}^{k_n}$$

and

$$q_2 = \mathbf{T}^{-l_1} \alpha_{s_1} \alpha_{t_1}^{-1} \mathbf{T}^{l_1} \dots \mathbf{T}^{-l_m} \alpha_{s_m} \alpha_{t_m}^{-1} \mathbf{T}^{l_m}$$

as in equation (6). Then  $q_1$ ,  $q_2$  and  $q_1 T^{k_1-k_2}$  all have the same domains, so m = n,  $k_i = l_i$  and  $u_i = s_i$  for all *i*. The effect of applying  $q_1$  (respectively  $q_2$ ) to a sequence  $\underline{x} = (x_n)_{n \in \mathbb{Z}}$  in this domain is to replace various segments of  $\underline{x}$  as specified in Lemma 2.5(b), while  $T^{k_1-k_2}$  translates  $\underline{x}q_i$  by  $(k_1 - k_2)$  spaces to the right. Since  $q_1 T^{k_1-k_2} = q_2$ , this forces  $k_1 - k_2 = 0$ , so  $k_1 = k_2$ , whence  $q_1 = q_2$  and the representation of  $y \in I_X$  in the form  $q T^k$  is unique. It is now easy to see that  $I_X$  is isomorphic to the semidirect product of  $Q_X$  and  $\langle T \rangle \cong \mathbb{Z}$  defined by the action  $T^k \cdot q = T^k q T^{-k}$  of  $\langle T \rangle$  on  $Q_X$ .

To prove part (c), let  $Y = \widetilde{X} \cup \{T\}$  and let  $\tau_1$  be the congruence on  $(Y \cup Y^{-1})^*$  generated by  $\rho \cup R_2$ . Then  $Inv \langle \widetilde{X} \cup \{T\} : R_2 \rangle = (Y \cup Y^{-1})^* / \tau_1$ ,

and since  $I_X$  satisfies all the relations in  $R_2$  we may write  $I_X \cong (Y \cup Y^{-1})^* / \tau_2$ for some congruence  $\tau_2$  on  $(Y \cup Y^{-1})^*$  with  $\tau_1 \subseteq \tau_2$ . If q is a word in  $(Y \cup Y^{-1})^*$  for which the sum of powers of T occurring in q is zero, then the proofs of Lemmas 2.3, 2.4 and 2.5 show that application of the relations  $R_2$  alone reduces q to a form as represented in the right-hand side of equation (6). Thus every element  $y \in (Y \cup Y^{-1})^*$  is  $\tau_1$ -related to a

side of equation (6). Thus every element  $y \in (Y \cup Y^{-1})^*$  is  $\tau_1$ -related to a word of the form  $q T^k$  where q is represented as in the right-hand side of equation (6). If  $(q_1 T^{k_1})\tau_1(q_2 T^{k_2})$ , then  $(q_1 T^{k_1})\tau_2(q_2 T^{k_2})$  so by the proof of part (b) of the present theorem and part (a) of Lemma 2.5 we see that  $k_1 = k_2$  and  $q_1 = q_2$ , so each  $\tau_1$ -class has a unique representative of the form  $q T^k$  (with q represented as in the right-hand side of (6)). Since each  $\tau_2$ -class has a unique representative of the same form it easily follows that  $\tau_1 = \tau_2$ . Hence  $I_X \cong \text{Inv}(\tilde{X} \cup \{T\} : R_2)$ .

We close this section by recording another structural property of the semigroups  $P_X$ ,  $I_X$  and  $Q_X$ . We define an inverse semigroup S with zero 0 to be 0-*E*-unitary if  $E(S) - \{0\}$  is a unitary subset of S: that is if  $e^2 = e$ ,  $ex \neq 0$  and  $(ex)^2 = ex$  for some e,  $x \in S$ , then  $x^2 = x$ . (Here, as usual, E(S) denotes the semilattice of idempotents of S).

**PROPOSITION 2.7.** The semigroups  $P_X$ ,  $I_X$  and  $Q_X$  are 0-E-unitary inverse monoids.

**PROOF.** This result is well-known (and easy to prove) for the polycyclic monoid, and hence for  $P_X$ . Let  $\alpha = \alpha^2 \in I_X$  and  $\gamma \in I_X$ . From (the proofs of) Lemmas 2.2 and 2.6 we may write

$$\gamma = \mathbf{T}^{-k_1} \alpha_{u_1} \alpha_{v_1}^{-1} \mathbf{T}^{k_1} \mathbf{T}^{-k_2} \alpha_{u_2} \alpha_{v_2}^{-1} \mathbf{T}^{k_2} \dots \mathbf{T}^{-k_m} \alpha_{u_m} \alpha_{v_m}^{-1} \mathbf{T}^{k_m} \mathbf{T}^{k_m}$$

(where the  $k_i$  satisfy the constraints of equality (6)) and  $\alpha = \beta_{i_1} \beta_{i_2} \dots \beta_{i_k}$ where  $\beta_{i_i}$  is the identity map on the set of sequences  $(x_n)_{n \in \mathbb{Z}}$  with  $x_{i_i} = a_t$ (a fixed element of X) for each t. Now  $\alpha \gamma$  is the restriction of  $\gamma$  to the intersection of the domain of  $\alpha$  with the domain of  $\gamma$ , so by Lemma 2.5(b), its domain consists of all those sequences w with  $w[k_i + m_i, k_i + m_i + |u_i|] =$  $u_i, x_{i_i} = a_i$  (for  $i = 1, \dots, m$  and  $t = 1, \dots, k$ ). If  $\alpha \gamma$  is a non-zero idempotent of  $I_X$  then  $\alpha \gamma$  must be the identity map on its domain so  $v_i = u_i$ for  $i = 1, \dots, m$  and so  $\gamma$  is also an idempotent of  $I_X$ . Hence  $I_X$  is 0-Eunitary. From this and the fact that  $EI_X = EQ_X$  it follows that  $Q_X$  is 0-E-unitary.

#### 3. Correspondence between semigroups and inverse semigroups

For every relation  $\Sigma = \{(u_i, v_i) : i \in I\} \subseteq X^* \times X^*$  on the free monoid  $X^*$  define

$$\begin{split} I(\Sigma) &= \langle \alpha_{u_i} \alpha_{v_i}^{-1}, \, \mathrm{T} : i \in I \rangle + \mathrm{EI}_{\mathrm{X}} \leq I_{\mathrm{X}} \, ; \\ P(\Sigma) &= I(\Sigma) \cap P_{\mathrm{X}} \, ; \\ Q(\Sigma) &= I(\Sigma) \cap Q_{\mathrm{X}} \, . \end{split}$$

Thus  $I(\Sigma)$ ,  $P(\Sigma)$  and  $Q(\Sigma)$  are all inverse submonoids of  $I_X$ . (Here, as before, + denotes join in the lattice of inverse submonoids of  $I_X$ ). In the present section we show that the map  $\Sigma \to P(\Sigma)$  sets up an isomorphism from the lattice of congruences on  $X^*$  onto a certain sublattice of the lattice of the lattice of inverse submonoids of  $P_X$ . We first need one preliminary concept.

The set  $P_X^+ = \{\alpha_u : u \in X^*\}$  is a submonoid of  $P_X$  isomorphic to the free monoid  $X^*$ : elements of  $P_X^+$  will be called *positive* elements of  $P_X$  in this paper (although some readers may prefer to label such elements "negative" in view of the isomorphism developed in Lemma 2.1 between  $P_X$  and the polycyclic monoid). An inverse submonoid R of  $P_X$  will be called a *positively self-conjugate* (PSC) submonoid of  $P_X$  if  $p R p^{-1} \subseteq R$  for every positive element  $p \in P_X^+$ .

**LEMMA 3.1.** Let R be a PSC submonoid of  $P_x$ . Then

(a) **R** is a full inverse submonoid of  $P_X$  (that is  $EP_X \subseteq R$ );

(b) if  $\alpha_{u}\alpha_{v}^{-1} \in \mathbb{R}$  and  $w \in X^{*}$  then  $\alpha_{uv}\alpha_{uv}^{-1}$ ,  $\alpha_{uv}\alpha_{uv}^{-1} \in \mathbb{R}$ .

**PROOF.** To prove part (a) note that since  $1 \in \mathbb{R}$  we have  $\alpha_u \alpha_u^{-1} = \alpha_u 1 \alpha_u^{-1} \in \mathbb{R}$  for all  $u \in X^*$ , so the result follows from Lemma 2.2(a). To prove part (b) note first that  $\alpha_{wu} \alpha_{wv}^{-1} = \alpha_w \alpha_u \alpha_v^{-1} \alpha_w^{-1} \in \mathbb{R}$  by the definition of PSC submonoid, while

$$\begin{aligned} \alpha_{uw} \alpha_{vw}^{-1} &= \alpha_u \alpha_w \alpha_w^{-1} \alpha_v^{-1} = \alpha_u \alpha_w \alpha_w^{-1} \alpha_v^{-1} \alpha_v \alpha_v^{-1} \\ &= \alpha_u \alpha_v^{-1} \alpha_v \alpha_w \alpha_w^{-1} \alpha_v^{-1} \\ &= \alpha_u \alpha_v^{-1} \alpha_{vw} \alpha_{vw}^{-1} \in \mathbb{R} \quad \text{by part (a)}. \end{aligned}$$

**LEMMA 3.2.** If R is a PSC submonoid of  $P_X$  then in  $I_X$  we have  $R = (R + \langle T \rangle + EI_X) \cap P_X$ .

**PROOF.** We need only prove that  $(R + \langle T \rangle + EI_X) \cap P_X \subseteq R$ . Notice first that  $EP_X \subseteq R$  by Lemma 3.1. From Lemma 2.2 it is clear that  $EI_X \subseteq$ 

 $EP_X + \langle T \rangle$ , so it suffices to prove that  $R = (R + \langle T \rangle) \cap P_X$ . By Lemma 2.5(b) and Lemma 3.1(b), every element of  $(R + \langle T \rangle) \cap Q_X$  can be represented in the form  $T^{-k_1}q_1T^{k_1}T^{-k_2}q_2T^{k_2}...T^{-k_m}q_mT^{k_m}$  where  $k_{i+1} > k_i + |q_i|$  for i = 1, ..., m-1 and  $q_i \in R$ . By the uniqueness of this representation, such an element is in  $P_X$  if and only if  $k_1 = 0$  and m = 1. Thus, if such an element belongs to  $P_X$  it also belongs to R.

**PROPOSITION 3.3.** If  $\Sigma$  is a relation on the free monoid  $X^*$  then  $P(\Sigma)$  is a PSC submonoid of  $P_X$ . Conversely, if R is a PSC submonoid of  $P_X$  then there is some relation  $\Sigma$  on  $X^*$  (in fact, some congruence  $\Sigma$  on  $X^*$ ) such that  $\mathbf{R} = P(\Sigma)$ .

PROOF. Let  $\Sigma = \{(u_i, v_i) : i \in I\}$  be a relation on  $X^*$ . Note that equation (3) (Lemma 2.3) implies that  $\alpha_u p \alpha_u^{-1} = T^{-|u|} p T^{|u|} \alpha_u \alpha_u^{-1}$  for all words  $u \in X^*$  and all  $p \in P_X$ . If  $p \in P(\Sigma)$  then p is a product of elements in  $EI_X$  and elements in  $I(\Sigma)$ , each of which is a product of powers of T and elements of the form  $\alpha_{u_i} \alpha_{v_i}^{-1}$ ,  $i \in I$ . It follows that  $T^{-|u|} p T^{|u|}$  is of the same form and hence  $\alpha_u p \alpha_u^{-1} \in I(\Sigma) \cap P_X = P(\Sigma)$ . So  $P(\Sigma)$  is a PSC submonoid of  $P_X$ .

Conversely, let R be a PSC submonoid of  $P_X$ . Define a relation  $\Sigma'$  on  $X^*$  by

(8) 
$$(u, v) \in \Sigma'$$
 if and only if  $\alpha_u \alpha_v^{-1} \in \mathbb{R}$ .

It is clear from Lemma 3.1 that  $\Sigma'$  is a congruence on  $X^*$ . Now  $I(\Sigma') = \langle \{\alpha_u \alpha_v^{-1} : (u, v) \in \Sigma'\} \rangle + EI_X + \langle T \rangle = R + \langle T \rangle$ , so  $P(\Sigma') = (R + \langle T \rangle) \cap P_X = R$ , by Lemma 3.2.

The main result of the paper is the following.

**THEOREM 3.4.** Let  $\Sigma$  be a relation on the free monoid  $X^*$  and let  $\Sigma^c$  denote the congruence on  $X^*$  generated by  $\Sigma$ . Then

- (a) for all  $u, v \in X^*$  we have  $(u, v) \in \Sigma^c$  if and only if  $\alpha_u \alpha_v^{-1} \in P(\Sigma)$ ;
- (b)  $P(\Sigma) = P(\Sigma^c);$
- (c) the correspondence  $\Sigma^c \to P(\Sigma)$  is an isomorphism between the lattice of congruences on  $X^*$  and the lattice of PSC submonoids of  $P_X$ .

**PROOF.** To prove part (a), suppose first that  $(s, t) \in \Sigma$  and w = psq,

z = ptq for some  $p, q \in X^*$ . Then

$$\begin{aligned} \alpha_w \alpha_z^{-1} &= \alpha_{psq} \alpha_{ptq}^{-1} = \alpha_p \alpha_s \alpha_q \alpha_q^{-1} \alpha_t^{-1} \alpha_p^{-1} \\ &= \mathbf{T}^{-|ps|} \alpha_q \alpha_q^{-1} \mathbf{T}^{|ps|} \alpha_p \alpha_s \alpha_t^{-1} \alpha_p^{-1} \quad \text{(by equation (3))} \\ &= \mathbf{T}^{-|ps|} \alpha_q \alpha_q^{-1} \mathbf{T}^{|ps|} \mathbf{T}^{-|p|} \alpha_s \alpha_t^{-1} \mathbf{T}^{|p|} \alpha_p \alpha_p^{-1} \quad \text{(by equation (3))} \\ &\in I(\mathbf{\Sigma}). \end{aligned}$$

Also,  $\alpha_w \alpha_z^{-1} \in P_X$ , so  $\alpha_w \alpha_z^{-1} \in P(\Sigma)$ . Suppose now that u = v is a consequence of the relations  $\Sigma$  (that is,  $(u, v) \in \Sigma^c$ ). Then there exists a sequence of words  $u = w_1, w_2, \ldots, w_n = v$  such that  $w_i = p_i s_i q_i$  and  $w_{i+1} = p_i t_i q_i$  for some words  $p_i, s_i, q_i, t_i$  with  $(s_i, t_i) \in \Sigma$  and  $i = 1, \ldots, n-1$ . Then by what was just proved,  $\alpha_{w_i} \alpha_{w_{i+1}}^{-1} \in P(\Sigma)$  for  $i = 1, \ldots, n-1$ . It follows that

$$\alpha_{u}\alpha_{v}^{-1} = \alpha_{w_{1}}\alpha_{w_{2}}^{-1}\alpha_{w_{2}}\alpha_{w_{3}}^{-1}\dots\alpha_{w_{n-1}}^{-1}\alpha_{w_{n}}^{-1} \in P(\Sigma).$$

Suppose conversely that  $\alpha_{\mu}\alpha_{\nu}^{-1} \in P(\Sigma)$ . Then  $\alpha_{\mu}\alpha_{\nu}^{-1} \in Q(\Sigma) + EQ_{\chi}$ , so  $\alpha_{\nu} \alpha_{\nu}^{-1}$  can be written as a product of elements in  $Q(\Sigma)$  and idempotents of  $I_x$ . Using Lemma 2.2(b) and the fact that  $Q(\Sigma) = (\langle \{\alpha_s \alpha_t^{-1} : (s, t) \in \mathcal{O}\} \}$  $\Sigma$ } +  $\langle T \rangle$  + EI<sub>X</sub>)  $\cap Q_X$ , we may write  $\alpha_u \alpha_v^{-1}$  (not necessarily uniquely) as a product of the form  $\alpha_u \alpha_v^{-1} = q_1 q_2 \dots q_m$  where  $q_i = T^{-k_i} \alpha_s \alpha_t^{-1} T^{k_i}$  for some  $k_i \in \mathbb{Z}$  and  $(s_i t_i) \in \Sigma \cup \Sigma^{-1} \cup \{i_X\}$ . (Here  $i_X$  denotes the identity map on X and  $\Sigma^{-1} = \{(s, t) : (t, s) \in \Sigma\}$ .) Now Dom $(q_i)$  consists of those sequences  $(x_n)_{n \in \mathbb{Z}}$  with  $x_{k_i+1} \dots x_{k_i+|s_i|} = s_i$  and  $\text{Dom}(\alpha_u \alpha_v^{-1})$  consists of those sequences within  $x_1 \dots x_{|u|} = u$ . Since  $\alpha_u \alpha_v^{-1} = q_1(q_2 \dots q_m)$  and  $\text{Dom}(q_1 \dots q_m) \subseteq \text{Dom} q_1$  we must have  $u = a_1 s_1 b_1$  for some  $a_1, b_1 \in$  $X^*$ . The effect of applying  $q_1$  to a sequence in its domain is to replace  $x_1 \dots x_{|u|} = a_1 s_1 b_1$  by  $a_1 t_1 b_1$ . Again, since  $\text{Dom}(q_1 q_2 \dots q_m) \subseteq \text{Dom}(q_1 q_2)$ we must have  $a_1t_1b_1 = a_2s_2b_2$  for some  $a_2, b_2 \in X^*$ . Continuing in this manner (by induction on m) and using the fact that  $q_1 q_2 \dots q_m = \alpha_n \alpha_n^{-1}$ , we obtain  $u = a_1 s_1 b_1$ ,  $a_1 t_1 b_1 = a_2 s_2 b_2$ ,  $a_2 t_2 b_2 = a_3 s_3 b_3$ , ...,  $a_{m-1} t_{m-1} b_{m-1} = a_m t_m b_m = v$ , for some  $a_i$ ,  $b_i \in X^*$ . It follows that  $(u, v) \in \Sigma^c$  (that is, u = v is a consequence of the relations  $\Sigma$ ). This proves part (a). Note also that  $\Sigma^c = \Sigma'$ , the congruence defined by equality (8) with  $R = P(\Sigma)$ . It is clear that  $P(\Sigma) \subseteq P(\Sigma^c)$  since  $\Sigma \subseteq \Sigma^c$ . Conversely, if  $\alpha_{\mu} \alpha_{\nu}^{-1} \in P(\Sigma^c)$  then  $(u, v) \in (\Sigma^c)^c = \Sigma^c$ , so by part (a),  $\alpha_u \alpha_u^{-1} \in P(\Sigma)$ . Hence  $P(\Sigma^c) \subseteq P(\Sigma)$ and part (b) is proved.

Suppose now that  $P(\Sigma_1) = P(\Sigma_2)$  for some congruences  $\Sigma_1$ ,  $\Sigma_2$  on  $X^*$ . Then if  $(u, v) \in \Sigma_1$  we have  $\alpha_u \alpha_v^{-1} \in P(\Sigma_1) = P(\Sigma_2)$  by part (a), so  $(u, v) \in \Sigma_1$ 

[13]

 $\Sigma_2$ , again by part (a). Hence  $\Sigma_1 \subseteq \Sigma_2$  and dually  $\Sigma_2 \subseteq \Sigma_1$ . It follows by Proposition 3.3 that the map  $\Sigma \to P(\Sigma)$  is a one-one map from the lattice of congruences on  $X^*$  onto the lattice of PSC submonoids of  $P_X$ . Again, a routine argument using part (a) shows that, for all congruences  $\Sigma_1$ ,  $\Sigma_2$  on  $X^*$ ,  $P(\Sigma_1 \cap \Sigma_2) = P(\Sigma_1) \cap P(\Sigma_2)$  and  $P(\Sigma_1 + \Sigma_2) = P(\Sigma_1) + P(\Sigma_2)$ , so the map  $\Sigma \to P(\Sigma)$  above is a lattice isomorphism. This proves part (c).

The following result was obtained by the authors for free monoids. Francis Pastijn noticed that it is true for arbitrary finitely generated monoids.

**COROLLARY 3.5.** The congruence lattice of every finitely generated monoid is isomorphic to a subalgebra lattice of some finitely generated algebra of finite type.

**PROOF.** If M is a finitely generated monoid, then  $M = X^* / \Sigma$ , for some finite set X and some congruence  $\Sigma$  on  $X^*$ . The congruence lattice on M is isomorphic to the lattice of congruences on  $X^*$  containing  $\Sigma$ . By Theorem 3.4 this lattice is isomorphic to the lattice of PSC submonoids of  $P_X$  containing  $P(\Sigma)$ . We shall introduce the following new unary operations on the inverse monoid  $P_X$ :

- (i) the unary operations  $\overline{\alpha}_x$ ,  $x \in X$ , where  $\overline{\alpha}_x(p) = \alpha_x p \alpha_x^{-1}$  for all  $p \in P_X$ ;
- (ii) the unary operation  $\alpha$ , such that  $\alpha: p \to \alpha(p)$  is a permutation of  $P_{\mathbf{x}}$  whose only nontrivial orbit is  $P(\Sigma)$ ;
- (iii) the unary operation  $\beta$ , such that  $\beta: p \to \beta(p)$  is the inverse of  $\alpha$ .

We observe that 1 belongs to every subalgebra of  $P_X$  and that  $P(\Sigma)$  is the least subalgebra of  $P_X$ . In fact, the subalgebras of  $P_X$  are precisely the PSC submonoids containing  $P(\Sigma)$ . Thus the lattice of congruences of M is isomorphic to the lattice of submonoids of the monoid  $P_X$  with additional unary operations.

REMARK 3.6. From an old result of Hanf (see [11] for an exposition of this) it follows that the congruence lattice of any finitely generated algebra A with finitely many operations is isomorphic to a subalgebra lattice of some finitely generated Moufang loop. This result is more general than Corollary 3.5. Nevertheless we included Corollary 3.5 in our paper because in order to construct the loop in Hanf's proof one needs the lattice of congruences of A, whereas one needs only the defining relations of A to construct our algebra  $P_X$ . Among the new operations of  $P(\Sigma)$  only one is not polynomial. Notice that in the group case the lattice of congruences of every finitely generated group  $A = \langle X \rangle$  is isomorphic to the lattice of subalgebras of the group A with

additional unary operations  $\alpha_x$  which provide conjugations by  $x, x \in X$ . This observation allows us to formulate the following.

**PROBLEM 3.7.** Is it possible to associate with every finitely generated monoid  $A = \langle X \rangle$  an inverse monoid  $B = \langle X \rangle$  in such a way that the congruence lattice of A is isomorphic to a subalgebra lattice of the monoid B with additional unary operations  $\alpha_x$  which provide conjugations by  $x, x \in X$ ?

Theorem 3.4 presents a solution to this problem in the case of free monoids.

### 4. Connections with other problems

Theorem 3.4 enables us to translate many important decision problems and properties of arbitrary monoids into equivalent problems and properties of inverse submonoids of the polycyclic monoid. In this section we provide a few remarks and observations along these lines. We first indicate how several standard properties of monoids may be reformulated. The notation is as in Section 3.

**PROPOSITION 4.1.** The monoid  $I(\Sigma)$  (respectively  $Q(\Sigma)$ ) is finitely generated if and only if all relations in  $\Sigma$  are consequences of a finite subset of the relations in  $\Sigma$  (that is,  $\Sigma^c = \Sigma_1^c$  for some finite set  $\Sigma_1 \subseteq \Sigma$ ).

**PROOF.** Suppose first that  $I(\Sigma) = \langle \{\alpha_u \alpha_v^{-1} : (u, v) \in \Sigma\} \rangle + \langle T \rangle$  is finitely generated. Then there is a finite set  $\{q_1, q_2, \ldots, q_n\}$  of generators for  $I(\Sigma)$ , each of which is a product of a (finite) number of powers of T and elements of the form  $\alpha_s \alpha_t^{-1}$  for some  $(s, t) \in \Sigma \cap \Sigma^{-1}$ . So there is a finite set  $\Sigma_1 \subseteq \Sigma$  such that  $I(\Sigma) = \langle \{\alpha_s \alpha_t^{-1} : (s, t) \in \Sigma_1\} \rangle + EI_X + \langle T \rangle = I(\Sigma_1)$ . It follows that  $P(\Sigma) = P(\Sigma_1)$  and hence by Theorem 3.4 that  $\Sigma^c = \Sigma_1^c$ .

Conversely suppose that  $\Sigma^c = \Sigma_1^c$  for some finite set  $\Sigma_1 \subseteq \Sigma$ . By Theorem 3.4 we have  $P(\Sigma) = P(\Sigma_1)$ , whence  $I(\Sigma) \cap P_X = I(\Sigma_1) \cap P_X$ . Thus if  $(u, v) \in \Sigma$  then  $\alpha_u \alpha_v^{-1} \in I(\Sigma) \cap P_X$ , so  $\alpha_u \alpha_v^{-1} \in I(\Sigma_1)$ . It follows that  $I(\Sigma) = \langle \alpha_u \alpha_v^{-1} : (u, v) \in \Sigma \rangle + EI_X + \langle T \rangle \subseteq I(\Sigma_1)$ , whence  $I(\Sigma) = I(\Sigma_1)$ , and  $I(\Sigma)$  is finitely generated.

The statements about the semigroup  $Q(\Sigma)$  follow from those about  $I(\Sigma)$ . Indeed, if  $Q(\Sigma)$  is finitely generated (as an inverse monoid with additional unary operation T) then  $I(\Sigma)$  is finitely generated as an inverse semigroup and so  $\Sigma^c = \Sigma_1^c$  for some finite subset  $\Sigma_1$  of  $\Sigma$ . Conversely, if  $\Sigma^c = \Sigma_1^c$  with finite  $\Sigma_1 \subseteq \Sigma$ , then  $I(\Sigma) = I(\Sigma_1)$  and so  $Q(\Sigma) = I(\Sigma) \cap Q_X = I(\Sigma_1) \cap Q_X =$  $Q(\Sigma)$ .

[14]

REMARK 4.2. Theorem 3.4 provides us with an analogue, for arbitrary congruences on the free monoid  $X^*$ , of the classical coset decomposition of a group relative to a subgroup. Let  $\Sigma$  be a relation on  $X^*$  and  $P(\Sigma)$  the associated PSC submonoid of  $P_X$ . A subset of  $P_X^+$  of the form  $P(\Sigma)\alpha_u \cap P_X^+$ for some  $u \in X^*$  will be called a *(right)* coset of  $P(\Sigma)$  in  $P_X^+$ . From Theorem 3.4 it is clear that if u and v are any words in  $X^*$  then  $(u, v) \in \Sigma^c$  iff  $\alpha_u \alpha_v^{-1} \in P(\Sigma)$  iff  $P(\Sigma) \alpha_u = P(\Sigma) \alpha_v$  and that the  $\Sigma^c$ -class [u] containing u is given by  $[u] = \{v \in X^* : \alpha_v \in P(\Sigma)\alpha_u\}$ . It follows that  $P_X^+$  decomposes as a disjoint union of right cosets of  $P(\hat{\Sigma})$  in the usual way. The map  $[u] \rightarrow$  $P(\Sigma)\alpha_{\mu} \cap P_{X}^{+}$  is a well-defined bijection from the set of  $\Sigma^{c}$ -classes of X<sup>\*</sup> onto the set of right cosets of  $P(\Sigma)$  in  $P_X^+$ . Since  $P(\Sigma)$  is a PSC-submonoid of  $P_{\mathbf{X}}, \ \alpha_u P(\Sigma) \alpha_u^{-1} \subseteq P(\Sigma) \text{ and so } \alpha_u P(\Sigma) \subseteq P(\Sigma) \alpha_u \text{ for all } u \in \mathbf{X}^*$ . It follows that  $P(\Sigma) \alpha_u P(\Sigma) \alpha_v \subseteq P(\Sigma) P(\Sigma) \alpha_u \alpha_v = P(\Sigma) \alpha_u \alpha_v$  for all  $u, v \in \mathbf{X}^*$ , so the right cosets of  $P(\Sigma)$  in  $P_X^+$  form a monoid in the obvious way and the map  $[u] \to P(\Sigma)\alpha_u \cap P_X^+$  is an isomorphism of  $M_{\Sigma} = Mon(X : \Sigma)$  onto the monoid of right cosets of  $P(\Sigma)$  in  $P_X^+$ . The cardinality of the set of right cosets of  $P(\Sigma)$  in  $P_X^+$  will be called the *index* of  $P(\Sigma)$  in  $P_X^+$ . It is clear from the above that  $M_{\Sigma} = \text{Mon}(X : \Sigma)$  is finite if and only if  $P(\Sigma)$  has finite index in  $P_{\mathbf{x}}^+$ .

**PROPOSITION 4.3.** Let  $\Sigma$  be a relation on  $X^*$  and  $M_{\Sigma} = Mon(X : \Sigma)$ . Then

- (a)  $M_{\Sigma}$  is left cancellative if and only if  $pP(\Sigma)p^{-1} \subseteq P(\Sigma)$  for all  $p \in P_{X}$ ; (b)  $M_{\Sigma}$  is right cancellative if and only if  $P(\Sigma) \{0\}$  is a unitary subset of  $P_{\mathbf{x}}$ .

**PROOF.** By Theorem 3.4 it follows that  $M_{\Sigma}$  is left cancellative if

$$\alpha_u \alpha_s \alpha_t^{-1} \alpha_u^{-1} \in P(\Sigma)$$
 implies  $\alpha_s \alpha_t^{-1} \in P(\Sigma)$ .

Since  $\alpha_u^{-1}\alpha_u = 1$ , it is clear that this implication holds if  $pP(\Sigma)p^{-1} \subseteq P(\Sigma)$  for all  $p \in P_X$ . Suppose on the other hand that  $M_{\Sigma}$  is left cancellative,  $\alpha_s \alpha_t^{-1} \in P(\Sigma)$  and  $v \in X^*$ . In order to show that  $pP(\Sigma)p^{-1} \subseteq P(\Sigma)$  for all  $p \in P_{\mathbf{X}}$ , it clearly suffices to show that  $\alpha_v^{-1} \alpha_s \alpha_t^{-1} \alpha_v \in P(\Sigma)$ , since  $P(\Sigma)$  is a PSC submonoid of  $P_{\mathbf{X}}$ . It is clear that  $\alpha_v^{-1} \alpha_s \alpha_t^{-1} \alpha_v \in P(\Sigma)$  if  $\alpha_v^{-1} \alpha_s \alpha_t^{-1} \alpha_v = 0$ , so suppose that  $\alpha_v^{-1} \alpha_s \alpha_t^{-1} \alpha_v \neq 0$ . There are several cases to consider. If  $|v| \leq |s|$  and  $|v| \leq |t|$ , then s = vw and t = vu for some  $w, u \in X^*$ , so  $\alpha_s \alpha_t^{-1} = \alpha_v \alpha_w \alpha_u^{-1} \alpha_v^{-1} \in P(\Sigma)$ , and  $(vw, vu) \in \Sigma^c$ , and hence  $(w, u) \in \Sigma^c$ by left cancellation, that is,  $\alpha_w \alpha_u^{-1} \in P(\Sigma)$ . Thus in this case,  $\alpha_v^{-1} \alpha_s \alpha_t^{-1} \alpha_v =$  $\alpha_{u}\alpha_{u}^{-1} \in P(\Sigma)$ , as required. If  $|v| \ge |s|$  and  $|v| \ge |t|$ , then v = sw = tu

for some w,  $u \in X^*$ , so  $\alpha_v^{-1} \alpha_s \alpha_t^{-1} \alpha_v = \alpha_w^{-1} \alpha_u \neq 0$ . Hence we must have either u = wp or w = up for some  $p \in X^*$ . Suppose that u = wp, some  $p \in X^*$ , so that  $\alpha_v^{-1} \alpha_s \alpha_t^{-1} \alpha_v = \alpha_p$ . Now v = sw = twp, so  $(sw, twp) \in \Sigma^c$ . Also,  $(s, t) \in \Sigma^c$ , so  $(sw, tw) \in \Sigma^c$ . Since  $M_{\Sigma}$  is left cancellative this implies that  $(1, p) \in \Sigma^c$ , so  $\alpha_p \in P(\Sigma)$ . The case w = upis similar. Hence  $\alpha_v^{-1} \alpha_s \alpha_t^{-1} \alpha_v \in P(\Sigma)$  in this case as well. Suppose next that  $|v| \leq |s|$  and  $|v| \geq |t|$ , so that s = vw and v = tu for some  $u, w \in X^*$ . Then  $\alpha_v^{-1} \alpha_s \alpha_t^{-1} \alpha_v = \alpha_w \alpha_u$ . Since  $(s, t) = \Sigma^c$  we have  $(tuw, t) \in \Sigma^c$ , so  $(uw, 1) \in \Sigma$ , by left cancellativity of  $M_{\Sigma}$ . Thus  $(wuwu, wu) \in \Sigma^c$  and so  $(wu, 1) \in \Sigma^c$ , again by left cancellation. This implies that  $\alpha_w \alpha_u \in P(\Sigma)$ , so again,  $\alpha_v^{-1} \alpha_s \alpha_t^{-1} \alpha_v \in P(\Sigma)$  in this case. The final case  $(|v| \geq |s|, |v| \leq |t|)$ is handled in a similar fashion, so  $\alpha_v^{-1} \alpha_s \alpha_t^{-1} \alpha_v \in P(\Sigma)$  in all cases. This establishes part (a).

To prove (b), suppose first that  $P(\Sigma) - \{0\}$  is unitary in  $P_X$  and that  $(su, tu) \in \Sigma^c$  for some  $u, s, t \in X^*$ . Then

$$\alpha_s \alpha_u \alpha_u^{-1} \alpha_t^{-1} = \alpha_s \alpha_t^{-1} (\alpha_t \alpha_u \alpha_u^{-1} \alpha_t^{-1}) \in P(\Sigma) - \{0\}.$$

Also  $\alpha_{tu}\alpha_{tu}^{-1} \in P(\Sigma) - \{0\}$ , so it follows that  $\alpha_s \alpha_t^{-1} \in P(\Sigma)$  since  $P(\Sigma) - \{0\}$  is unitary. Hence  $M_{\Sigma}$  is right cancellative. Suppose conversely that  $M_{\Sigma}$  is right cancellative. Suppose also that  $\alpha_s \alpha_t^{-1} \alpha_u \alpha_v^{-1} \in P(\Sigma) - \{0\}$  and  $\alpha_u \alpha_v^{-1} \in P(\Sigma)$  for some  $s, t, u, v \in X^*$ . Then  $\alpha_v \alpha_u^{-1} \in P(\Sigma)$  and so  $\alpha_s \alpha_t^{-1} \alpha_u \alpha_v^{-1} \alpha_v \alpha_u^{-1} = \alpha_s \alpha_t^{-1} \alpha_u \alpha_u^{-1} \in P(\Sigma) - \{0\}$ . We must have either t = uw or u = tw for some  $w \in X^*$ . In the former case,  $\alpha_s \alpha_t^{-1} \alpha_u \alpha_u^{-1} = \alpha_s \alpha_w^{-1} \alpha_u^{-1} \in P(\Sigma)$ , so  $(s, uw) \in \Sigma^c$ , and  $(s, t) \in \Sigma^c$ : in the latter case,  $\alpha_s \alpha_t^{-1} \alpha_u \alpha_u^{-1} = \alpha_s \alpha_w \alpha_w^{-1} \alpha_t^{-1} \in P(\Sigma)$ , so  $(sw, tw) \in \Sigma^c$ , whence  $(s, t) \in \Sigma^c$  since  $M_{\Sigma}$  is right cancellative. Hence in both cases  $\alpha_s \alpha_t^{-1} \in P(\Sigma)$ , so  $P(\Sigma) - \{0\}$  is right unitary in  $P_X$ . Since  $P(\Sigma)$  is an inverse submonoid of the inverse monoid  $P_X$ ,  $P(\Sigma) - \{0\}$  is a unitary subset of  $P_X$ , as required.

We close the paper by reformulating a few of the basic decision problems about monoids into equivalent problems about the polycyclic monoid.

REMARK 4.4. Let  $\Sigma$  be a relation on  $X^*$ . Then the word problem for the monoid  $M_{\Sigma} = \text{Mon}\langle X : \Sigma \rangle$  is equivalent to the membership problem for the submonoid  $P(\Sigma)$  of  $P_X$  (that is, the problem of deciding, for words  $u, v \in X^*$ , whether  $\alpha_u \alpha_v^{-1} \in P(\Sigma)$  or not.)

**PROOF.** This is obvious from Theorem 3.4 since  $(u, v) \in \Sigma^c$  if and only if  $\alpha_u \alpha_u^{-1} \in P(\Sigma)$ .

[17]

On the one hand, this remark points out the complex nature of PSC submonoids of the polycyclic monoid. For example, if  $M_{\Sigma} = Mon\langle X : \Sigma \rangle$  is a finitely presented monoid with undecidable word problem, then the membership problem for the corresponding PSC submonoid  $P(\Sigma)$  of  $P_X$  is also undecidable. Note that  $P(\Sigma)$  is the PSC submonoid of  $P_X$  generated by  $\{\alpha_u \alpha_v^{-1} : (u, v) \in \Sigma\}$ ; that is,  $P(\Sigma)$  is the smallest full inverse submonoid of  $P_X$  containing  $\{\alpha_u \alpha_v^{-1} : (u, v) \in \Sigma\}$  and satisfying the condition  $\alpha_s \alpha_t^{-1} \in P(\Sigma)$  implies  $\alpha_{ws} \alpha_{wt}^{-1} \in P(\Sigma)$  for all  $w \in X^*$ . Hence

**REMARK** 4.5. If X is a set with at least two elements then there exists a finite set R of elements of  $I_X$  such that the membership problem for the subsemigroup of  $I_X$  generated by R is undecidable.

**PROOF.** This follows easily from the existence of a semigroup with two generators and undecidable word problem [5] (and Theorem 3.4).

On the other hand, several classical decision problems for semigroups may be translated into a rather appealing equivalent form by use of Remark 4.3 (that is, Theorem 3.4). For example, as far as the authors are aware, it is not known whether the word problem for one-relation semigroups of the form  $M = \text{Mon}\langle X : u = v \rangle$   $(u, v \text{ fixed words in } X^*)$  is decidable or not. (A positive solution, due to Adian [1], is known, for example, in the case v = 1.) The corresponding PSC-submonoid of  $P_X$  is generated (as a PSC-submonoid of  $P_X$ ) by the single element  $\alpha_u \alpha_v^{-1}$ . Thus we have

**REMARK** 4.6. The word problem for one-relation semigroups is equivalent to the membership problem for one-generator PSC submonoids of the polycyclic monoid.

Other well-known problems may be reformulated along these lines. The authors believe that progress on some of these problems may result from a detailed examination of the PSC submonoids of the polycyclic monoid.

### Acknowledgements

Joint work on this paper commenced when the authors met at a workshop on semigroups organized by John Rhodes in Berkeley in July 1989 and continued when the authors met again at the Mal'cev algebra conference in Novosibirsk in August 1989. Part of the work was completed while the first author was visiting the L.I.T.P., University of Paris VI in October-November 1989 and later Monash University in the first half of 1990. The authors would like to thank Francis Pastijn for several comments and in particular for providing the proof of Corollary 3.5, which is more general than their original result.

### References

- S. I. Adian, 'Defining relations and algorithmic problems for groups and semigroups', Trudy Mat. Inst. Steklov 85 (in Russian) Am. Math. Soc. translation 152, 1967.
- [2] A. H. Clifford and G. B. Preston, The algebraic theory of semigroups, Math. Surveys Monographs 7, vol. 1, vol. 2 (Amer. Math. Soc., Providence, 1961–1967).
- [3] E. S. Lyapin, Semigroups (Moscow, 1960) (in Russian).
- [4] S. Margolis and J. Meakin, 'Inverse monoids, trees and context-free languages', Trans. Amer. Math. Soc. (to appear).
- [5] Yu. V. Matiyasevich, 'Investigation on some algorithmic problems in algebra and number theory', Trudy MIAN SSSR, 168 (1984), 218-235.
- [6] M. Nivat, 'Sur les automates a memoire pile', in: Proceedings of the International Computing Symposium, Bonn, 1970 (ed. W. Itzeld), (North Holland, Amsterdam, 1970) pp. 655-663.
- [7] M. Nivat and J. F. Perrot, 'Une generalisation du monoide bicyclique', C.R. Acad. Sci. Paris Sér. I Math. A 271 (1970), 824–827.
- [8] M. Petrich, Inverse semigroups, (Wiley, New York, 1984).
- [9] J. Sakarovitch, Syntaxe des langages de Chomsky, (Th. Sc. Math., Univ. Paris 7, 1979).
- [10] J. Stephen, 'Presentations of inverse monoids', J. Pure and Applied Algebra 63 (1990), 81-112.
- [11] Th. P. Whaley, Algebras satisfying the descending chain condition for subalgebras, (Ph.D. Thesis, Vanderbilt University, 1968).

Department of Mathematics and Statistics University of Nebraska-Lincoln Lincoln, Nebraska 68588-0323 U.S.A.