

LETTERS TO THE EDITOR

Dear Editor,

The extremal index in 10 seconds

Introduction. In a recent paper, Smith [3] introduces a method to calculate the extremal index of a stationary Harris chain $\{X_n\}$. Loosely speaking, a stationary sequence $\{X_n\}$ with marginal distribution F has extremal index θ if

$$P\{\max(X_1, \dots, X_n) \leq u_n\} - (F(u_n))^{n\theta} \rightarrow 0,$$

as $n \rightarrow \infty$, for sequences u_n with $F^n(u_n) = c \in (0, 1)$. The main assumption of Smith is that the transition density q of the Harris chain satisfies

$$\lim_{u \rightarrow \infty} q(u, u+x) = h(x),$$

for some limiting function h , with $h(x) \geq 0$ and $\int h(x)dx \leq 1$.

In this letter we show that there is a simple way to compute the extremal index. The numerical method given here is adapted from the Wiener–Hopf algorithm developed by Grübel [2], designed to calculate the distribution of the stationary waiting time of a stable $G/G/1$ queue.

Implementation. From (2.6)–(2.8) of [3],

$$(1) \quad \theta = \int_{-\infty}^0 e^x P\{S_1 < x, S_2 < x, \dots \mid S_0 = 0\} dx,$$

where $S_0 = 0, S_1, S_2, \dots$ is a random walk with stepsize density h . In order to facilitate the use of Grübel’s algorithm, consider the random walk $S'_k = -S_k$, with density $g, g(x) = h(-x)$. To avoid trivial cases, assume that

$$(2) \quad \int xg(x)dx > 0.$$

As h may be defective, with missing mass transferred to $-\infty$, g may be defective with mass at ∞ , in which case the expectation in (2) is infinite.

Define

$$M = \inf_{k \geq 0} \{S'_k \mid S'_0 = 0\};$$

condition (2) implies that M is finite. Its distribution can be computed with Grübel’s algorithm. From (1), we obtain

$$\begin{aligned}
 \theta &= \int_0^\infty e^{-x} \mathbf{P} \left\{ \inf_{k \geq 1} S'_k > x \mid S'_0 = 0 \right\} dx \\
 (3) \quad &= \int_0^\infty e^{-x} \int_x^\infty \mathbf{P} \{ M > x - y \} g(y) dy dx \\
 &= \int_0^\infty e^{-x} \int_{-\infty}^\infty \mathbf{P} \{ M > x - y \} g(y) dy dx;
 \end{aligned}$$

the last equality follows from $\mathbf{P} \{ M > 0 \} = 0$.

Note that θ can be expressed as the probability of an event: let Z , Y , and M be independent random variables, with Z exponentially distributed with mean 1, M as defined, and Y a copy of the stepsize of the random walk S'_0, S'_1, \dots . Then

$$(4) \quad \theta = \mathbf{P} \{ Y + M - Z > 0 \}.$$

This leads to the following algorithm for the computation of θ . Steps 3–6 below are steps (iii)–(viii) in Grübel’s algorithm; for details we refer to [2]. Note that steps 1 and 2 differ from the first two steps in Grübel’s algorithm: for the $G/G/1$ case the stepsize distribution first has to be computed as the difference of two independent random variables representing an interarrival time and a service time.

Step 1. Discretize the distribution of the stepsize Y . For a large positive integer m the distribution of Y is approximated by the vector p of length $2m$ with

$$p(k) = \mathbf{P} \left\{ (k - \frac{1}{2})h < Y \leq (k + \frac{1}{2})h \right\}, \quad k = -m, -m + 1, \dots, m - 1.$$

The gridsize h should be as small as possible, whereas m should be chosen so that $(-mh, mh)$ gives a fair coverage of the range of both Y and Z . For computational efficiency it is advised to take m equal to a power of 2.

Step 2. Calculate the discrete Fourier transform (fft) fp on $2m$ points of the vector p :

$$fp(k) = \sum_{n=-m}^{m-1} p(n) e^{2\pi i k n / 2m}, \quad k = 0, \dots, 2m - 1.$$

For the non-defective case we need the fft of the tailvector r ,

$$\begin{aligned}
 r(k) &= \mathbf{P} \{ Y > h(k + \frac{1}{2}) \}, & k = 0, 1, \dots, m - 1, \\
 &= -\mathbf{P} \{ Y \leq h(k + \frac{1}{2}) \}, & k = -m, \dots, -1,
 \end{aligned}$$

given by

$$fr(k) = \sum_{n=-m}^{m-1} r(n) e^{2\pi i k n / 2m}, \quad k = 0, \dots, 2m - 1.$$

Step 3. Calculate $fs = -\log(fr)$, where $x \rightarrow \log x$ denotes the complex logarithm.

Step 4. Calculate the inverse Fourier transform of fs :

$$s(k) = \frac{1}{2m} \sum_{n=-m}^{m-1} fs(n) e^{-2\pi i k n / 2m}, \quad k = -m, -m+1, \dots, m-1,$$

(in shorthand $s = \text{ifft}(fs)$), and define

$$\begin{aligned} sm(k) &= s(k), & k \leq 0, \\ &= 0, & k > 0. \end{aligned}$$

The vector sm is an approximation to the harmonic renewal function of the descending ladder height H^- , corresponding to the random walk generated by Y .

Step 5. Apply the Fourier transform and then the transformation $y \rightarrow 1 - \exp(-y)$ to obtain the Fourier transform of the (defective) probability mass function fh of the ladder height:

$$fh = 1 - \exp(-(\text{fft}(sm))).$$

Step 6. Compute the Fourier transform of M from the Fourier transform of the ladder height by

$$fm = (1 - fh(0)) / (1 - fh).$$

Step 7. On the same grid $-mh, \dots, (m-1)h$ make a discretization of an exponentially distributed random variable Z , with mean 1, and calculate the discrete Fourier transform $fminz$ on $2m$ points of $-Z$. Compute the product $fm \cdot fr \cdot fminz$, this is an approximation to the Fourier transform of $Y + M - Z$.

Step 8. Apply the inverse Fourier transform and sum the probabilities corresponding to positive subscripts to get the extremal index θ . Adding half of the probability at zero generally improves the accuracy.

Note. In the defective case the vector of tail probabilities r is not needed and fs in step 3 can be computed directly from the defective vector p by $fs = -\log(1 - fp)$.

Runtimes. We illustrate the method with Example 2 from Smith [3]. In this example $H(z) = \int_{-\infty}^z h(x) dx = (1 + e^{-rz})^{1/r-1}$, where $r > 1$. For m we use powers of 2: $m = 2^k$, $k = 8, 9, \dots$; $h = 15/m$ gives adequate coverage for r between 2 and 5.

Using a 386 20 MHz personal computer and 386-MATLAB, we obtained the values of θ , for $r=2$ shown in Table 1.

The approximation can be improved considerably by applying a simple extrapolation method. It is conjectured by Grübel in [2] that the approximation of M_h , with gridsize h , is of the form

$$P\{M_h \leq t\} = P\{M \leq t\} + ch + o(h),$$

for $h \rightarrow 0$. If this is true and the density g of the stepsize of the random walk S'_k is sufficiently smooth, then the same discretization error is present in θ . If we call θ_k the

TABLE 1

| m | Computer time | θ | m | Computer time | θ |
|----------|---------------|----------|----------|---------------|----------|
| 2^8 | 2 sec | 0.32148 | 2^{12} | 25 sec | 0.32809 |
| 2^9 | 4 sec | 0.32498 | 2^{13} | 59 sec | 0.32831 |
| 2^{10} | 6 sec | 0.32675 | 2^{14} | 173 sec | 0.32842 |
| 2^{11} | 12 sec | 0.32764 | | | |

approximation of θ , based on $m=2^k$, then $2 * \theta_{k+1} - \theta_k$ should have a discretization error of the order $o(h)$, since the discretization parameter h used to calculate θ_{k+1} is half the value of the discretization parameter for θ_k . Embrechts *et al.* [1] give a rigorous mathematical derivation of this method, called Richardson's deferred approach to the limit, and apply it to compound distributions.

We found as a rule of thumb that the maximum of h^2 and 10 times the maximum of the missing probability masses can be used as an error upper bound, where h is the smallest of the two grid sizes used for the extrapolation.

The calculation with extrapolation for $m=2^8$ takes only 6 seconds of computing time. It can even be done on an 80286 or 8086 personal computer using PC-MATLAB (the computing time is then ≈ 30 seconds). Table 2 presents the values of θ for $r=2, 3, 4, 5$, which are also included in Smith [3]; as before, we took $h=15/m$.

TABLE 2

| r | m | θ | m | θ |
|-----|-------|----------|----------|----------|
| 2 | 2^8 | 0.32848 | 2^{13} | 0.32853 |
| 3 | 2^8 | 0.15794 | 2^{13} | 0.15806 |
| 4 | 2^8 | 0.09218 | 2^{13} | 0.09234 |
| 5 | 2^8 | 0.06024 | 2^{13} | 0.06043 |

The value 0.0616 for $r=5$ given by Smith seems to be slightly off.

Remark. As noted by the referee, for $r \rightarrow 1$ the value of mh has to be increased considerably to obtain a fair coverage of the distribution of the stepsize, hence for r close to 1 the algorithm becomes unstable.

An alternative approximation for θ can be given through

$$\theta' := \lim_{u \rightarrow \infty} P\{X_2 < u \mid X_1 > u\}.$$

It is easy to show that θ' is an upper bound for the extremal index θ : $\theta \leq \theta' \leq 1$. It depends on the length of the arrays that can be stored efficiently in the working memory at which point one should abandon the algorithm and switch to the approximation θ' . It seems that the algorithm can safely be used in the range $r \geq 1.01$, provided that mh is chosen in such a way that the missing probability mass of Y and Z is small; this missing mass directly affects the error. For $r=1.01$, $mh=800$ and extrapolating on the values $m=2^{13}$ and $m=2^{14}$ we obtain: $\theta=0.98629$. In this case $\theta'=2^{1/r}-1=0.98632$.

References

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