LETTERS TO THE EDITOR

ON THE ATTAINED WAITING TIME

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Abstract

By using properties of up- and downcrossings of the sample functions of the workload process and of the attained waiting-time process for a G/G/1 queueing model, a direct proof of a theorem proved by Sakasegawa and Wolff is given.

WORKLOAD PROCESS; SAMPLE FUNCTIONS; STATIONARY DISTRIBUTIONS

Sakasegawa and Wolff (1990) show by using sample function arguments that for the FIFO G/G/1 queueing model the workload process v_t and the attained waiting-time process η_t possess the same stationary distribution, if such distributions exist. However their proof is somewhat artificial (see their use of preemptive LIFO).

A direct proof of their Theorem 1 proceeds as follows. Consider a busy cycle c with n the number of customers served; τ_1, \ldots, τ_n are the service times of these customers, w_1, \ldots, w_n their successive actual waiting times, i the idle time, so

(1)
$$\boldsymbol{c} = \boldsymbol{\tau}_1 + \cdots + \boldsymbol{\tau}_n + \boldsymbol{i}.$$

The attained service time η_t at epoch t is by definition the time between t and the arrival epoch of the customer being served at epoch t. In Figure 1 the sample function of the workload process v_t and the corresponding η_t -process during the busy cycle c are shown, with n = 4.

Define for $v \ge 0$,

(2)

$$d(v) := \# \text{ downcrossings of } v_t, \ 0 \le t \le c \text{ with level } v, \ (*)$$

$$u(v) := \# \text{ upcrossings of } v_t, \ 0 \le t \le c \text{ with level } v, \ (*)$$

$$\delta(v) := \# \text{ upcrossings of } \eta_t, \ 0 \le t \le c \text{ with level } v, \ (*)$$
(3)

(5)
$$\omega(v) := \# \text{ downcrossings of } \eta_t, \ 0 \le t \le c \text{ with level } v, \ (^\circ).$$

Note that in the figure d(v) = 3; the upcrossings are there indicated by °, the downcrossings by *. It is immediately evident from the geometry of the sample functions (see Cohen (1977), (1982)) that with probability 1, for $v \ge 0$,

(4)
$$d(v) = u(v), \qquad \delta(v) = \omega(v),$$

(5)
$$\boldsymbol{u}(\boldsymbol{v}) = \boldsymbol{\delta}(\boldsymbol{v});$$

and

(6)
$$\boldsymbol{d}(\boldsymbol{v}) = \frac{d}{d\boldsymbol{v}} \int_0^c (\boldsymbol{v}_t < \boldsymbol{v}) \, dt, \qquad \boldsymbol{\delta}(\boldsymbol{v}) = \frac{d}{d\boldsymbol{v}} \int_0^c (\boldsymbol{\eta}_t < \boldsymbol{v}) \, dt,$$

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where we use the notation

(7)
$$(\boldsymbol{v}_t < \boldsymbol{v}) \equiv 1_{\boldsymbol{v}_t < \boldsymbol{v}} \text{ and } \int_0^c (\boldsymbol{v}_t < \boldsymbol{v}) dt \equiv \int_0^\infty (\boldsymbol{v}_t < \boldsymbol{v}, \boldsymbol{c} \ge t) dt,$$

for the indicator function and the integral. Since

(8)
$$\mathbf{i} = \left\{ \int_0^c (\mathbf{v}_i < \mathbf{v}) \, dt \right\}_{\mathbf{v}=0+} = \left\{ \int_0^c (\mathbf{\eta}_i < \mathbf{v}) \, dt \right\}_{\mathbf{v}=0+},$$

integration of (6), using the boundary conditions (8) yields, via (4) and (5), that with probability 1

(9)
$$\int_0^{\epsilon} (\boldsymbol{v}_t < \boldsymbol{v}) dt = \int_0^{\epsilon} (\boldsymbol{\eta}_t < \boldsymbol{v}) dt, \qquad \boldsymbol{v} \ge 0.$$

Because

$$(\boldsymbol{v}_t < \boldsymbol{v}) = 1 - (\boldsymbol{v}_t \ge \boldsymbol{v}),$$

we have from (9)

$$\int_0^c (\boldsymbol{v}_t \geq \boldsymbol{v}) dt = \int_0^c (\boldsymbol{\eta}_t \geq \boldsymbol{v}) dt,$$

which is Theorem 1 of Sakasegawa and Wolff (1990).

References

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