

A PRESENTATION FOR A GROUP OF AUTOMORPHISMS OF A SIMPLICIAL COMPLEX

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Introduction. The Bass–Serre theorem supplies generators and relations for a group of automorphisms of a tree. Recently K. S. Brown [2] has extended the result to produce a presentation for a group of automorphisms of a simply connected complex, the extra ingredient being relations which come from the 2-cells of the complex. Suppose G is the group, K the complex and L the 1-skeleton of K . Then an extension of $\pi_1(L)$ by G acts on the universal covering space of L (which is of course a tree) and Brown's technique is to apply the work of Bass and Serre to this action. Our aim is to give a direct elementary proof of Brown's theorem which makes no use of covering spaces, deals with the Bass–Serre theorem as a special case and clarifies the roles played by the various generators and relations.

1. Preliminaries. By a group of automorphisms of a simplicial complex K we mean a group of homeomorphisms of the underlying polyhedron $|K|$ whose elements permute the simplexes of K .

A directed edge e of K is a physical edge plus a choice of one of its vertices as initial vertex $i(e)$. The second vertex is then written $t(e)$ and called the terminal vertex. Making the other choice for $i(e)$ produces the reverse \bar{e} of e . From now on we refer to directed edges simply as edges. Let V denote the set of vertices and E the set of edges of K . These two sets together with the functions

$$\begin{aligned} E &\rightarrow E, & e &\rightarrow \bar{e} \\ E &\rightarrow V \times V, & e &\rightarrow (i(e), t(e)) \end{aligned}$$

form a *graph* X in the sense of Serre [3] because we clearly have $\bar{\bar{e}} = e$, $\bar{e} \neq e$ and $i(\bar{e}) = t(e)$ for each $e \in E$.

If G is a group of automorphisms of K we have a natural action of G on V and on E such that

$$gi(e) = i(ge) \quad \text{and} \quad gt(e) = t(ge)$$

for each $g \in G$ and $e \in E$. We shall assume that edges of K are *never reversed* by the action of G . More formally, if $g \in G$ and $e \in E$ then $ge \neq \bar{e}$. Thus the quotient X/G has the structure of a graph. Because X comes from a simplicial complex the initial and terminal vertices of an edge are always different. Of course this property may well be lost on passage to X/G .

A path in K (and in X) joining vertex u to vertex v is an ordered string of edges $e_1 e_2 \dots e_n$ such that $i(e_1) = u$, $i(e_{k+1}) = t(e_k)$ for $1 \leq k \leq n-1$, and $t(e_n) = v$. A path of

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the form $e\bar{e}$ will be called a *round trip*. Let $e_1 \dots e_k e_{k+1} \dots e_n$ be a path and suppose $i(e_k), t(e_k), t(e_{k+1})$ are the three vertices of a triangle of K . Let e be the edge of this triangle determined by $i(e) = i(e_k), t(e) = t(e_{k+1})$. Replacing the pair $e_k e_{k+1}$ by e in our original path is called taking a *short cut*.

We shall work with a complex which is both connected and simply connected. Each of these notions has a combinatorial description. A complex is connected if any two distinct vertices can be joined by a path. It is simply connected if two paths which join the same pair of vertices u, v are always homotopic. That is to say we can convert one path into the other (keeping a path from u to v throughout) by a finite number of steps each of which involves the addition or removal of either a round trip or a short cut.

We adopt the usual notation whereby G_u denotes the stabilizer of the vertex u . If $g \in G$ happens to fix u we write g_u for the element g thought of as a member of G_u . Of course G_e denotes the stabilizer of the edge e . If u is a vertex of e then G_e is a subgroup of G_u .

Recall that a graph is a *tree* if any two of its vertices may be joined by a path, and any path which joins a vertex to itself must contain a round trip.

With G, K, X as above choose a maximal tree M in X/G and lift it [3, Proposition I.14] to a subtree T of X . The vertices of T form a set of representatives for the action of G on the vertices of X . For each pair of edges f, \bar{f} from $X/G - M$, select one, say f , and lift it to an edge e of X which has its initial vertex x in T . Exactly one vertex z of T lies in the same orbit as $t(e)$ and we choose an element γ_f from G that maps z onto $t(e)$. We can now lift \bar{f} to $(\gamma_f)^{-1}\bar{e}$. This has its initial vertex z in T and $\gamma_f^{-1} = (\gamma_f)^{-1}$ sends the vertex x of T to its terminal vertex (Fig. 1). Finally we extend the correspondence $f \rightarrow \gamma_f$ over the edges of M by setting $\gamma_f = 1$ (the identity element of G) whenever $f \in M$.

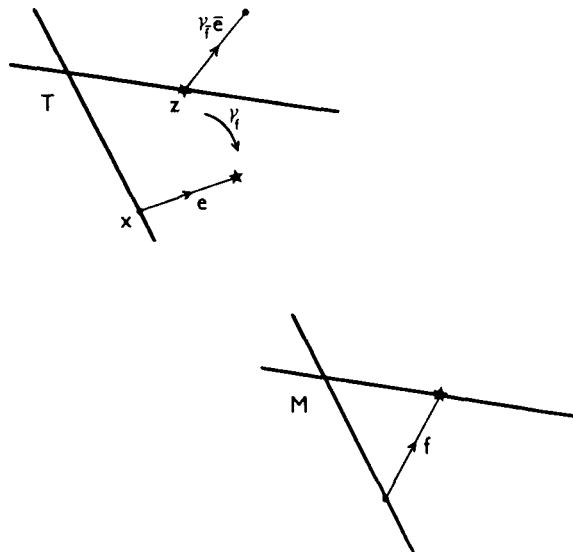


Figure 1

The elements of all the G_w , where w is a vertex of T , and the γ_f , where f is an edge of X/G , together generate G . They satisfy the relation

$$\gamma_{\bar{f}} g_x \gamma_f = (\gamma_{\bar{f}} g \gamma_f)_z$$

whenever g fixes the canonical lift e of f , plus relations which come from the triangles of K and which will be described in the next section.

2. Tail wagging. Let $g \in G$ and let $e_1 e_2 \dots e_n$ be a path which joins a vertex v of T to gv . If the path lies entirely in T then $gv = v$ because no two vertices of T lie in the same orbit. Therefore $g = g_v$, where as usual g_v denotes the element g interpreted as a member of G_v . Otherwise there is a first edge e_m that is not in T . The path $e_m e_{m+1} \dots e_n$ will be called the *tail* of $e_1 e_2 \dots e_n$. Let x_1 be the initial vertex of e_m . Project e_m into X/G to give an edge f_1 . The canonical lift e^1 of f_1 into X has its initial vertex in T , so $i(e^1) = x_1$. Choose an element $a_{x_1} \in G_{x_1}$ which sends e^1 to e_m . Let

$$e_k^1 = (\gamma_{\bar{f}_1} a_{x_1}^{-1}) e_k$$

for $m + 1 \leq k \leq n$ and replace $e_1 e_2 \dots e_n$ by the new path $e^1_{m+1} e^1_{m+2} \dots e^1_n$. We call this process *tail wagging*. Our new path begins at

$$z_1 = t(\gamma_{\bar{f}_1} e^1) = i(e^1_{m+1})$$

which is a vertex of T and ends at $(\gamma_{\bar{f}_1} a_{x_1}^{-1} g)v$; see Fig. 2. We walk along it to the first point x_2 where it quits T and repeat the above procedure. Since we shorten the tail at each step we eventually obtain a path which lies entirely in T and ends at, say,

$$(\gamma_{\bar{f}_r} a_{x_r}^{-1} \dots \gamma_{\bar{f}_2} a_{x_2}^{-1} \gamma_{\bar{f}_1} a_{x_1}^{-1} g)v.$$

Then $\gamma_{\bar{f}_r} a_{x_r}^{-1} \dots \gamma_{\bar{f}_1} a_{x_1}^{-1} g$ must fix v , say $\gamma_{\bar{f}_r} a_{x_r}^{-1} \dots \gamma_{\bar{f}_1} a_{x_1}^{-1} g = a_v \in G_v$. We now have

$$g = a_{x_1} \gamma_{f_1} \dots a_{x_r} \gamma_{f_r} a_v.$$

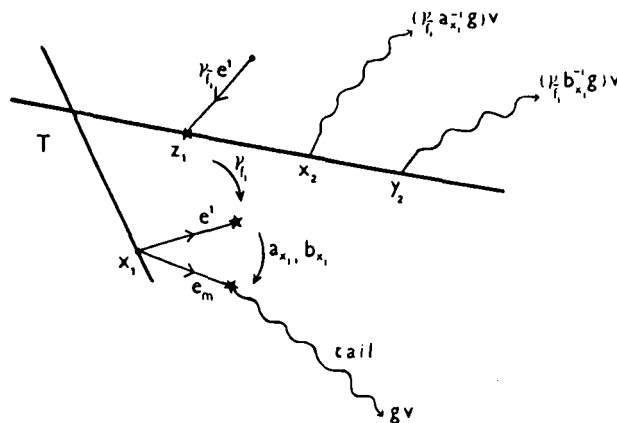


Figure 2

This shows that the elements of the G_w , $w \in T$, together with the γ_f , f an edge of X/G , do indeed generate G .

Now for the extra relations mentioned in §1. Our group G acts on the collection of all triangles in K . From each orbit we choose a triangle which has a vertex in T . Let Δ be such a triangle and let v be a vertex of Δ which belongs to T . Walking round the boundary of Δ (there is a choice of direction here) produces a path which joins v to $v = 1v$, and tail wagging this path expresses the identity element of G as a word r_Δ in our generators. The missing relations have the form

$$r_\Delta = 1,$$

one for each orbit of triangles in K .

With the notation established above let $*G_w$ denote the free product of the stabilizers of the vertices of T , and F the free group generated by symbols λ_f , one for each edge of X/G . Let R be the normal consequence in $(*G_w) * F$ of the words

$$\begin{aligned} &\lambda_f \text{ (} f \text{ an edge of } M), \\ &\lambda_{\bar{f}}\lambda_f \text{ (for each edge of } X/G), \\ &\lambda_{\bar{f}}g_x\lambda_f(\gamma_{\bar{f}}g\gamma_f)^{-1} \text{ (when } g \text{ fixes the canonical lift } e \text{ of } f), \text{ and} \\ &r_\Delta^\lambda \text{ (obtained from } r_\Delta \text{ by changing each occurrence of } \gamma_f \\ &\text{ to } \lambda_f, \text{ one such for each orbit of triangles).} \end{aligned}$$

PRESENTATION THEOREM. *If G is a group of automorphisms of a connected simply connected complex K , and if no edge of K is reversed by the action of G , then G is isomorphic to $[(*G_w) * F]/R$.*

If K is one-dimensional, so that X is a tree, this is the Bass–Serre Theorem [3, Théorème I.13]. For dimension two or more the extra relations were provided by K. S. Brown [2].

3. Homotopy. We shall produce an isomorphism

$$\psi : G \rightarrow [(*G_w) * F]/R$$

as follows. Take a vertex v of T as base point. Given $g \in G$, choose a path α in K which joins v to gv . By tail wagging α we express g as a word $a_{x_1}\gamma_{f_1} \dots a_{x_r}\gamma_{f_r}a_v$ and we define

$$\psi(g) = a_{x_1}\lambda_{f_1} \dots a_{x_r}\lambda_{f_r}a_vR.$$

Of course various choices are involved here and we must show that ψ is well defined. For a particular path α joining v to gv the first ambiguity occurs in the choice of the element $a_{x_1} \in G_{x_1}$ which maps e^1 to e_m . That a different choice at this and subsequent stages gives the same coset for $\psi(g)$ is verified exactly as in [1]. As to the choice of the actual path α we need only check that altering α by the addition or removal of a single round trip or short cut makes no difference to the value of $\psi(g)$.

Let α be the string $e_1 e_2 \dots e_n$ leading to the word $a_{x_1} \lambda_{f_1} \dots a_{x_r} \lambda_{f_r} a_v$ in $(*G_w) * F$. Suppose we add the *round trip* $e_0 \bar{e}_0$. Let f be the image of e_0 in X/G and e the canonical lift of f back into X , so that $i(e) = x$ lies in T . Tail wagging the new path gives a word of the form

$$a_{x_1} \lambda_{f_1} \dots a_{x_k} \lambda_{f_k} (a_x \lambda_f) (\lambda_{\bar{f}}) (a_x^{-1}) a_{x_{k+1}} \lambda_{f_{k+1}} \dots a_{x_r} \lambda_{f_r} a_v$$

if x is distinct from all the x_i , or possibly

$$a_{x_1} \lambda_{f_1} \dots a_{x_k} \lambda_{f_k} (a'_x \lambda_f) (\lambda_{\bar{f}}) (a_x^{-1} a_{x_{k+1}} \lambda_{f_{k+1}}) \dots a_{x_r} \lambda_{f_r} a_v$$

if $x = x_{k+1}$. We have used parentheses to emphasise groups of symbols which correspond to a *single* tail wag. Removing a round trip has the opposite effect and clearly neither process changes the coset of the original word.

Suppose now we start with α and take a *short cut*, calling the new path β . Tail wagging β proceeds as for α until we reach a vertex w of T from which the modified α runs along two sides of a triangle Δ_0 , whereas what remains of β short cuts along the third side. Clearly α and β give the same R -coset if $r_{\Delta_0}^\lambda$ belongs to R . The following lemma shows that this is the case.

LEMMA. Let Δ be a triangle which has a vertex u in T and suppose that $r_\Delta^\lambda \in R$. If g sends Δ to a triangle Δ_0 which also has a vertex in T then $r_{\Delta_0}^\lambda \in R$.

Proof. We may as well assume that the situation is represented by Fig. 3, the only possible ambiguity being the direction in which we choose to go round one or other of the triangles. As is easily checked a reversal of direction produces a word which is R -equivalent to the inverse of the original word, and does not affect our lemma.

For $m = 1, 2, 3$ let f_m be the projection of e_m into X/G and e^m the canonical lift of f_m back into X with $i(e^m) \in T$. Tail wagging $e_1 e_2 e_3$ gives, say,

$$r_\Delta^\lambda = a_u \lambda_{f_1} a_{x_2} \lambda_{f_2} a_w \lambda_{f_3} b_u \tag{1}$$

and taking the boundary of Δ_0 in the direction shown leads to

$$r_{\Delta_0}^\lambda = (a_w \lambda_{f_3}) (b_u c_u a_u \lambda_{f_1}) (a_{x_2} \lambda_{f_2}) d_w \tag{2}$$

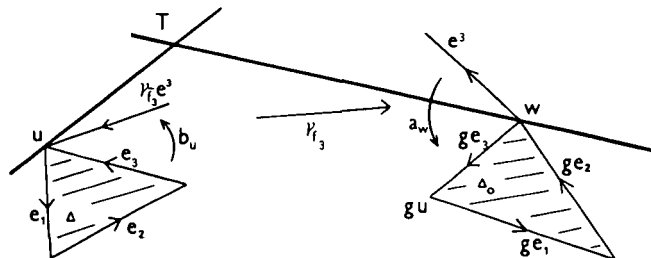


Figure 3

where c_u fixes the edge e_3 . Now

$$\begin{aligned} d_w &= [(a_w \gamma_{f_3} b_u c_u a_u \gamma_{f_1} a_{x_2} \gamma_{f_2})^{-1}]_w \\ &= [(a_u \gamma_{f_1} a_{x_2} \gamma_{f_2})^{-1} (a_w \gamma_{f_3} b_u c_u)^{-1}]_w \\ &= [(a_w \gamma_{f_3} b_u) (a_u \gamma_{f_3} b_u c_u)^{-1}]_w \text{ by (1)} \\ &= a_w (\gamma_{f_3} b_u c_u^{-1} b_u^{-1} \gamma_{f_3})_w a_w^{-1} \\ &= a_w \lambda_{f_3} b_u c_u^{-1} b_u^{-1} \lambda_{f_3} a_w^{-1} \end{aligned}$$

because $b_u c_u b_u^{-1}$ fixes $\gamma_{f_3} \bar{e}_3$. Substitution in (2) gives

$$r_{\Delta_0}^\lambda = (a_w \lambda_{f_3} b_u c_u) r_\Delta^\lambda (a_w \lambda_{f_3} b_u c_u)^{-1}$$

and shows that $r_{\Delta_0}^\lambda$ belongs to R .

4. Composite paths. To complete the proof of the Presentation Theorem we show that ψ is a homomorphism and a bijection.

Given elements $g, h \in G$ join v to gv and hv by paths α, β respectively. Tail wag α and β to give say

$$\begin{aligned} \psi(g) &= a_{x_1} \lambda_{f_1} \dots a_{x_r} \lambda_{f_r} a_v R, \\ \psi(h) &= b_{x'_1} \lambda_{f'_1} \dots b_{x'_r} \lambda_{f'_r} b_v R. \end{aligned}$$

Now α followed by the image of β under g joins v to $(gh)v$. Using this composite path to compute $\psi(gh)$ gives

$$\psi(gh) = (a_{x_1} \lambda_{f_1}) \dots (a_{x_r} \lambda_{f_r}) (a_v) (b_{x'_1} \lambda_{f'_1}) \dots (b_{x'_r} \lambda_{f'_r}) b_v R$$

provided x'_1 is not the same as v , or

$$\psi(gh) = (a_{x_1} \lambda_{f_1}) \dots (a_{x_r} \lambda_{f_r}) (a_v b_{x'_1} \lambda_{f'_1}) \dots (b_{x'_r} \lambda_{f'_r}) b_v R$$

if $x'_1 = v$. Therefore $\psi(gh) = \psi(g)\psi(h)$ and ψ is a homomorphism.

Our construction of ψ ensures that if $\psi(g) = R$ then $g = 1$. So ψ is injective. The cosets $h_w R$ (w a vertex of T and $hw = w$) and $\lambda_f R$ (f an edge of X/G) together generate $[(\ast G_w) \ast F]/R$. We show ψ surjective by checking that $\psi(h) = h_w R$ and $\psi(\gamma_f) = \lambda_f R$. If x, y are vertices of T write \overline{xy} for the geodesic (shortest path) in T which joins x to y . This geodesic is unique because T is a tree. Suppose h fixes the vertex w of T . Let x be the closest vertex of T to v such that \overline{xw} is left fixed by h , and note that \overline{ux} followed by the image of \overline{xv} under h joins v to hv . This path leaves T for the first time at x and a single tail wag using h_x^{-1} shows that $\psi(h) = h_x R$. But h fixes all of the geodesic \overline{xw} . Hence $h_x R = h_w R$ as required.

Finally, if f is an edge of X/G with canonical lift e in X then (with the usual notation) the composite path \overline{vx} followed by e followed by $\gamma_f(\overline{zv})$ joins v to $\gamma_f v$. This path leaves T for the first time at x and a single tail wag by $\gamma_{\bar{f}}$ shows $\psi(\gamma_f) = \lambda_f R$. This completes the argument.

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