

ON POLYNOMIALS WITH RELATED LEVEL SETS

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If p is a polynomial in one real variable and $p(x) = p(-x)$ then p has only even powers of x and is thus a polynomial in x^2 . If p is a polynomial in n variables and $p(x_1, \dots, x_n) = p(y_1, \dots, y_n)$ when $x_1^2 + \dots + x_n^2 = y_1^2 + \dots + y_n^2$ then p is a polynomial in q where $q(x_1, \dots, x_n) = x_1^2 + \dots + x_n^2$.

The problem considered in this note is this: For which polynomials q is it true that if $p(x) = p(y)$ whenever $q(x) = q(y)$ then p is a polynomial in q ? Such polynomials q will be said to satisfy (*). If the problem is posed for polynomials with complex variables, the answer is simple: any polynomial in n complex variables satisfies (*) (Theorem 1). However the problem is not as simple for polynomials with real variables. We give two classes of polynomials in one variable satisfying (*), neither class containing the other: if q is a polynomial of degree n and q has a level set containing n points, then q satisfies (*) (Theorem 2). If q is a polynomial such that the polynomial $Q(x, y) = [q(x) - q(y)] / (x - y)$ is irreducible and q is not 1:1, then q satisfies (*) (Theorem 3). Of course, x^3 , being 1:1, doesn't satisfy (*) and more generally the composition of two polynomials $q_0 \circ q_1$ does not satisfy (*) if q_0 is 1:1 on the range of q_1 (of course q_1 not being a constant). Thus $x^3 + 3x^2 + 3x (= (x + 1)^3 - 1)$ doesn't satisfy (*) yet $x^3 + 4x^2 + 3x (= x(x + 1)(x + 3))$ does satisfy (*).

THEOREM 1. *Let $q(z_1, \dots, z_n)$ be a polynomial in n complex variables. Let $p(z_1, \dots, z_n)$ be another such that p is constant on the level sets of q . Then p is a polynomial in q .*

Proof. We may assume q is not a constant. The function f from the range of q to C (the complex plane) defined by $f(q(z_1, \dots, z_n)) = p(z_1, \dots, z_n)$ is well defined by hypothesis. We show that it is a polynomial. It is possible to specialize all but one of the variables of q so that q defines a nonconstant polynomial, say q_0 , in just one variable. Its range is C . By the same specialization p defines a polynomial p_0 . If $q'_0(z) \neq 0$, then f is differentiable at $q_0(z)$ with derivative $p'_0(z) / q'_0(z)$. Since q_0 is open and p_0 continuous, $f = p_0 \circ q_0^{-1}$ is continuous on C and since f is analytic except at a finite number of points, f is an entire function. Since f has a pole at ∞ , $(q_0(z) \rightarrow \infty \text{ implies } z \rightarrow \infty \text{ implies } p_0(z) \rightarrow \infty)$, f is a polynomial.

THEOREM 2. *Let q be a polynomial in one real variable and of degree n . If q has some level set containing n points, then q satisfies (*).*

Proof. The hypothesis guarantees that q has an infinite number of level sets with n points. We show by induction on the degree of p that if p is constant on an infinite number of those level sets of q containing n points, then p is a polynomial in q . It is clearly true if p has degree 0. Thus suppose the assertion is known to be true

for any polynomial of degree less than the degree of p . Let $S = \{r_1, r_2, \dots, r_n\}$ be a level set of q on which p is constant. Then $p - p(r_1)$ is divisible by $(x - r_1)(x - r_2) \cdots (x - r_n)$ and therefore by $q - q(r_1)$. Thus $p(x) - p(r_1) = p_1(x)[q(x) - q(r_1)]$, p_1 is of lower degree than p , and p_1 is constant on all those level sets of q , other than S , that p is constant on. By our inductive hypothesis p_1 is a polynomial in q and hence so is p .

THEOREM 3. *Let q be a polynomial in one real variable such that $Q(x, y) = [q(x) - q(y)]/(x - y)$ is irreducible. If q is not 1:1 then q satisfies (*).*

Proof. The hypotheses guarantee that $S = \{y: q(x) - q(y) = 0 \text{ has a solution other than } x = y\}$ is infinite. Let p be a polynomial which is constant on an infinite number of the level sets of q which meet S . We show by induction on the degree of p that p is a polynomial in q . It is clear if p has degree 0. Suppose p has degree k and that if p_0 has degree less than k and is constant on an infinite number of level sets of q meeting S , then p_0 is a polynomial in q . Let $P(x, y) = [p(x) - p(y)]/(x - y)$. By an application of the Euclidean algorithm one may, as follows, show that Q divides P . (This argument is modeled after one appearing in [1, p. 291].) Define inductively polynomials r_k, s_k, R_k by:

$$\begin{aligned} r_1(y)P(x, y) &= q_1(x, y)Q(x, y) + R_2(x, y) \\ r_2(y)Q(x, y) &= q_2(x, y)R_2(x, y) + R_3(x, y) \\ &\vdots \\ r_{n-1}(y)R_{n-1}(x, y) &= q_{n-1}(x, y)R_n(x, y) + R_{n+1}(x, y) \end{aligned}$$

where d_k , the degree of $R_k(x, y)$ considered as a polynomial in x over the field of rational functions in y , becomes progressively smaller, $d_{n+1} = 0$ and $d_n \neq 0$ (let $Q = R_1$). There are infinitely many numbers y such that $q(x) = q(y)$ and $p(x) = p(y)$ have a common solution $x(y)$ not equal to y . Hence $P(x(y), y)$ and $Q(x(y), y)$ both vanish for infinitely many y . This means $R_{n+1}(y)$ has an infinite number of zeros and must be zero. Any irreducible factor of $R_n(x, y)$ which is of positive degree in x must divide both $Q(x, y)$ and $P(x, y)$. Since Q is irreducible, this means Q divides P . Thus $p(x) - p(y) = R(x, y)[q(x) - q(y)]$ and letting $p_1(x) = R(x, 0)$, $p(x) - p(0) = p_1(x)[q(x) - q(0)]$. The degree of p_1 is less than k and p_1 is constant on those level sets of q on which p is constant (other than the level set containing 0) and therefore p_1 is a polynomial in q and consequently p is a polynomial in q .

Let $q(x) = (x^2 - 1)(x^2 - 4)$. Then q meets the hypotheses of Theorem 2, but not Theorem 3, since $q(x) - q(y)$ is divisible by $x^2 - y^2$ and so $[q(x) - q(y)]/(x - y)$ is divisible by $x + y$. Let $q(x) = x^4 - x$. Then q meets the hypotheses of Theorem 3 but not Theorem 2. For a straightforward calculation shows that $x^3 + x^2y + xy^2 + y^3 - 1$ is irreducible.

REFERENCE

1. L. V. Ahlfors, *Complex analysis*, McGraw-Hill, New York (second edition), 1966.
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