# THE NUMBER OF FACTORS IN A PAPERFOLDING SEQUENCE 

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We prove that the number of factors of length $k$ in any paperfolding sequence is equal to $4 k$ once $k \geqslant 7$.

## 1. Introduction

A factor of an infinite sequence $u=(u(n))_{n \geqslant 0}$ with values in $A$ is a word on $A$ occurring as $u(n) u(n+1) \cdots u(n+k-1)$ for some $n ; k$ is called the length of the factor.

The study of factors of infinite sequences goes back at least to Thue $[\mathbf{1 5}, \mathbf{1 6}]$ and has interested mathematicians and computer scientists working in combinatorics, symbolic dynamics, finitely generated groups, number theory, formal languages ... .

Among the questions which have been addressed is the problem of computing for a given finite sequence $u$ its complexity function $P_{u}$, where $P_{u}(k)$ is the number of factors of length $k$ in $u$. We quote here some results:

- if for some $k$ one has $P_{u}(k) \leqslant k$, then $u$ is ultimately periodic (see [13] for example),
- the sequences with minimal complexity which are not ultimately periodic satisfy $P_{u}(k)=k+1$; these are called Sturmian sequences (see [7] for example),
- if $u$ is an automatic sequence (in the sense of [5]), then one has $P_{u}(k) \leqslant$ $C k$ for some constant $C[6] ;$
- if $u$ is the Thue-Morse sequence, then $P_{u}(k)$ has been computed [4,10]; it depends upon the digits of the binary expansion of $k$. More precisely the sequence $\left(P_{u}(k+1)-P_{u}(k)\right)_{k \geqslant 0}$ has only finitely many values and is an automatic sequence;
- if $u$ is an automatic sequence satisfying some technical requirements, then $\left(P_{u}(k+1)-P_{u}(k)\right)_{k \geqslant 0}$ is also automatic [14].

[^0]Very recently Shallit and the author proved in [2] that, for $u$ a generalised RudinShapiro sequence in the sense of [1], ( $u$ counts the parity of the number of blocks $1 * \cdots * 1$ in the binary expansion of $n$ ), the function $P_{u}(k)$ is ultimately affine. This result is somewhat surprising when compared to the complicated case of the Thue-Morse sequence.

We will prove here that the number of factors of length $k$ of any paperfolding sequence (see [ 8 ] for instance for a definition) is equal to $4 k$, provided $k \geqslant 7$. As a corollary we obtain that all generalised Rudin-Shapiro sequences in the sense of [12] (which except for the classical Rudin-Shapiro sequence are different from the sequences studied in [2]) have an ultimately affine complexity.

## 2. A quick survey of paperfolding

We recall that a paperfolding sequence is the sequence of ridges and valleys obtained by unfolding a sheet of paper which has been folded infinitely many times (see [3, 8, $11,12]$ ). In other words the sequence $(u(n))_{n \geqslant 0}$ is a paperfolding sequence if and only if

$$
\begin{aligned}
u(4 n)=0 \quad & (\text { respectively } 1) \\
u(4 n+2)=1 \quad & (\text { respectively } 0) \\
u(2 n+1) & \text { is a paperfolding sequence. }
\end{aligned}
$$

Another way of generating these sequences is to view them as Toeplitz sequences [9]: given an infinite binary sequence $i=(i(n))_{n \geqslant 0}$ (sequence of "folding instructions"), one defines the paperfolding sequence $u_{i}=\left(u_{i}(n)\right)_{n \geqslant 0}$ with folding instructions $i$ by successively "filling holes":

- first step, one writes down the sequence $(i(0) \overline{i(0)})^{\infty}$ at the even places, which gives

$$
i(0) \bullet \overline{i(0)} \bullet i(0) \bullet \overline{i(0)} \bullet i(0) \bullet \overline{i(0)} \bullet \ldots
$$

- second step, one writes down the sequence $(i(1) \overline{i(0)})^{\infty}$ at the even holes, which gives

$$
i(0) i(1) \overline{i(0)} \bullet i(0) \overline{i(1)} \overline{i(0)} \bullet i(0) i(1) \overline{i(0)} \bullet \cdots
$$

and so on; the limit obtained after an infinite number of steps is the sequence $u_{i}$.

Finally we mention the generation of a paperfolding sequence by "perturbed symmetry" (see [3]).

If $a$ is a letter ( 0 or 1 ), we define the operator $T_{a}$ by: for every word $M, T_{a}(M)=$ $M a \widetilde{M}$ (where $\widetilde{M}$ is obtained from $M$ by reading $M$ backwards, then replacing the 0 's by 1 's and the 1 's by 0 's).

Given a sequence $a_{0} a_{1} \cdots$, one then obtains a paperfolding sequence by starting, say, from 0 , and successively applying the operators $T_{a_{j}}$ :

$$
\begin{aligned}
T_{a_{0}}(0) & =0 a_{0} 1 \\
T_{a_{1}} T_{a_{0}}(0) & =0 a_{0} 1 a_{1} 0 \overline{a_{0}} 1 \\
T_{a_{2}} T_{a_{1}} T_{a_{0}}(0) & =0 a_{0} 1 a_{1} \overline{a_{0}} 1 a_{2} 0 a_{0} 1 \overline{a_{1}} 0 \overline{a_{0}} 1
\end{aligned}
$$

... .

Note that the word obtained at each step is of length $2^{r}-1$ for some $r \geqslant 1$.

## 3. O-factors and E-factors in a paperfolding sequence

In the sequel we say that a factor of the paperfolding sequence $u_{i}=\left(u_{i}(n)\right)_{n \geqslant 0}$ (with folding instructions i) is an O-factor (respectively an E-factor) of $u_{i}$ if it occurs in $u_{i}$ as $u_{i}(n) u_{i}(n+1) \cdots u_{i}(n+k-1)$ with $n$ odd (respectively $n$ even). Note that a factor can simultaneously be an O-factor and an E -factor (for instance the factor 0 and the factor 1 are simultaneously $O$-factors and $E$-factors of any paperfolding sequence).

We take a sequence of folding instructions beginning with $i(0)=\alpha, i(1)=\beta$, and write it as $\alpha \beta j$ (so $j$ is defined by $j(n)=i(n+2)$ ).

Applying the Toeplitz process twice, we obtain:

$$
\begin{equation*}
\alpha \beta \bar{\alpha} \bullet \alpha \bar{\beta} \bar{\alpha} \bullet \alpha \beta \bar{\alpha} \bullet \alpha \bar{\beta} \bar{\alpha} \bullet \alpha \beta \bar{\alpha} \bullet \cdots \tag{*}
\end{equation*}
$$

We are now ready to state two lemmata:
Lemma 1. The E-factors of length $\geqslant 4$ of a paperfolding sequence $u_{\alpha \beta j}$ are in one of the following four disjoint classes:

$$
\begin{array}{ll}
\alpha \beta \bar{\alpha} \bullet & \cdots \\
\bar{\alpha} \bullet \alpha \bar{\beta} & \cdots \\
\alpha \bar{\beta} \bar{\alpha} \bullet & \cdots \\
\bar{\alpha} \bullet \alpha \beta & \ldots
\end{array}
$$

The $O$-factors of length $\geqslant 4$ of a paperfolding sequence $u_{\alpha \beta j}$ are in one of the following four disjoint classes:

$$
\begin{array}{cl}
\beta \bar{\alpha} \bullet \alpha & \cdots \\
\bullet \alpha \bar{\beta} \bar{\alpha} & \cdots \\
\bar{\beta} \bar{\alpha} \bullet \alpha & \cdots \\
\bullet \alpha \beta \bar{\alpha} & \cdots
\end{array}
$$

Proof: Inspection of (*) reveals that the E-factors (respectively the O-factors) begin as written. Moreover the classes one obtains are disjoint regardless of $\alpha, \beta$ and the holes.

LEmMA 2. If a factor of a paperfolding sequence has length greater than or equal to 7 , then it cannot be simultaneously an o-factor and an e-factor.

Proof: Once again inspecting $\left(^{*}\right)$ one sees that a factor of $u_{\alpha \beta j}$ of length greater than or equal to 7 begins in one of the following ways:

1- $\alpha \beta \bar{\alpha} \bullet \alpha \bar{\beta} \bar{\alpha} \quad$.
2- $\beta \bar{\alpha} \bullet \alpha \bar{\beta} \bar{\alpha} \quad .$.
3- $\bar{\alpha} \bullet \alpha \bar{\beta} \bar{\alpha} \bullet \alpha \quad \cdots$
4- $\quad \bar{\beta} \bar{\alpha} \bullet \alpha \beta \quad \ldots$
5- $\alpha \bar{\beta} \bar{\alpha} \bullet \alpha \beta \bar{\alpha} \quad \cdots$
6- $\bar{\beta} \bar{\alpha} \bullet \alpha \beta \bar{\alpha} \bullet \ldots$
7- $\bar{\alpha} \bullet \alpha \beta \bar{\alpha} \bullet \alpha \quad \cdots$
8- $\bullet \alpha \bar{\alpha} \bullet \alpha \bar{\beta} \quad \cdots$.
The cited words of length 7 are different (whatever the values of $\alpha, \beta$, and the holes). In particular the E-factors (numbered 1, 3, 5, 7) and the O-factors (numbered $2,4,6,8$ ) (!) are different.

## 4. A recurrence relation for the number of O-factors and E-factors

For $u_{i}=\left(u_{i}(n)\right)_{n \geqslant 0}$ a paperfolding sequence with folding instructions $i$, one defines:
$g_{i}(k)$ is the number of $E$-factors of length $k$ in $u_{i}$,
$h_{i}(k)$ is the number of $O$-factors of length $k$ in $u_{i}$.
How can we obtain an E-factor of length $4 k$ in $u_{\alpha \beta j}$ ? Once again inspecting (*) we see that an E-factor of length $4 k$ of $u_{\alpha \beta j}$ is of one of the following types:

1: $\alpha \beta \bar{\alpha} \bullet \alpha \bar{\beta} \bar{\alpha} \bullet \alpha \beta \bar{\alpha} \bullet \ldots$
2: $\bar{\alpha} \bullet \alpha \bar{\beta} \bar{\alpha} \bullet \alpha \beta \bar{\alpha} \bullet \alpha \bar{\beta} \quad \ldots$
3: $\alpha \bar{\beta} \bar{\alpha} \bullet \alpha \beta \bar{\alpha} \bullet \alpha \bar{\beta} \bar{\alpha} \bullet \ldots$
4: $\bar{\alpha} \bullet \alpha \beta \bar{\alpha} \bullet \alpha \bar{\beta} \bar{\alpha} \bullet \alpha \beta \quad \ldots$
and these types are disjoint for $4 k \geqslant 4$ (Lemma 1). Moreover if one considers the subsequent steps in the Toeplitz process, one sees that the remaining holes are filled exactly by the E-factors of $u_{j}$ in Cases 1 and 2 and by the O-factors of $u_{j}$ in Cases 3 and 4.

Hence:

$$
\forall k \geqslant 1 \quad g_{\alpha \beta_{j}}(4 k)=2 g_{j}(k)+2 h_{j}(k)
$$

In the same way we compute the quantities $g_{\alpha \beta j}(4 k+r)$ and $h_{\alpha \beta j}(4 k+r)$ for $r=$ $0,1,2,3$, obtaining the following proposition:

Proposition. One has the following relations for $k \geqslant 1$ :

$$
\begin{aligned}
g_{\alpha \beta j}(4 k) & =2 g_{j}(k)+2 h_{j}(k), \\
h_{\alpha \beta j}(4 k) & =2 g_{j}(k)+2 h_{j}(k), \\
g_{\alpha \beta j}(4 k+1) & =2 g_{j}(k)+2 h_{j}(k), \\
h_{\alpha \beta j}(4 k+1) & =g_{j}(k)+g_{j}(k+1)+h_{j}(k)+h_{j}(k+1), \\
g_{\alpha \beta j}(4 k+2) & =g_{j}(k)+g_{j}(k+1)+h_{j}(k)+h_{j}(k+1), \\
h_{\alpha \beta j}(4 k+2) & =g_{j}(k)+g_{j}(k+1)+h_{j}(k)+h_{j}(k+1), \\
g_{\alpha \beta j}(4 k+3) & =g_{j}(k)+g_{j}(k+1)+h_{j}(k)+h_{j}(k+1), \\
h_{\alpha \beta j}(4 k+3) & =2 g_{j}(k+1)+2 h_{j}(k+1) .
\end{aligned}
$$

## 5. Counting the factors of a paperfolding sequence

We are now able to prove the theorem:
THEOREM. For any paperfolding sequence $u_{i}=\left(u_{i}(n)\right)_{n \geqslant 0}$, the number of factors of length $k, P_{u_{i}}(k)$, is given by:

$$
P_{u_{i}}(1)=2, P_{u_{i}}(2)=4, P_{u_{i}}(3)=8, P_{u_{i}}(4)=12, P_{u_{i}}(5)=18, P_{u_{i}}(6)=23
$$

and for all $k \geqslant 7, P_{u_{i}}(k)=4 k$.
Proof: Define

$$
V_{i}(k)=\left(\begin{array}{c}
g_{i}(k) \\
h_{i}(k) \\
g_{i}(k+1) \\
h_{i}(k+1)
\end{array}\right)
$$

The recurrence relations in the previous paragraph can be rewritten as:

$$
V_{\alpha \beta j}(4 k+4)=A_{r} V_{j}(k) \quad \forall k \geqslant 1
$$

where the matrices $A_{r}$ are given by

$$
\begin{array}{ll}
A_{0}=\left(\begin{array}{llll}
2 & 2 & 0 & 0 \\
2 & 2 & 0 & 0 \\
2 & 2 & 0 & 0 \\
1 & 1 & 1 & 1
\end{array}\right), & A_{1}=\left(\begin{array}{llll}
2 & 2 & 0 & 0 \\
1 & 1 & 1 & 1 \\
1 & 1 & 1 & 1 \\
1 & 1 & 1 & 1
\end{array}\right) \\
A_{2}=\left(\begin{array}{llll}
1 & 1 & 1 & 1 \\
1 & 1 & 1 & 1 \\
1 & 1 & 1 & 1 \\
0 & 0 & 2 & 2
\end{array}\right), & A_{3}=\left(\begin{array}{llll}
1 & 1 & 1 & 1 \\
0 & 0 & 2 & 2 \\
0 & 0 & 2 & 2 \\
0 & 0 & 2 & 2
\end{array}\right)
\end{array}
$$

We then notice (still using (*) and considering all possibilities for $\alpha, \beta$, and the holes) that, for every sequence of folding instructions $j$, one has:

$$
g_{j}(1)=h_{j}(1)=2, g_{j}(2)=h_{j}(2)=4, g_{j}(3)=4, h_{j}(3)=8, g_{j}(4)=h_{j}(4)=8
$$

Now we claim that for every sequence of instructions $j$, one has:

$$
\begin{aligned}
& \forall k \geqslant 1, k \text { even, } \quad V_{j}(k)=\left(\begin{array}{c}
2 k \\
2 k \\
2 k \\
2 k+4
\end{array}\right), \\
& \forall k \geqslant 1, k \text { odd }, \quad V_{j}(k)=\left(\begin{array}{c}
2 k-2 \\
2 k+2 \\
2 k+2 \\
2 k+2
\end{array}\right) .
\end{aligned}
$$

The proof that this is true for every sequence $j$ and for every $k$ in $[1,4 n-1]$ follows easily by induction on $n$ and is left as an exercise for the reader.

Finally, using Lemma 2, we obtain that, for every sequence of instructions $i$, one has:

$$
\forall k \geqslant 7 \quad P_{u_{i}}(k)=g_{i}(k)+h_{i}(k)=4 k
$$

The values of $P_{u_{i}}(k)$ for $1 \leqslant k \leqslant 6$ are computed by hand using (*) for a final time.
6. The number of factors of the generalised Rudin-Shapiro sequences
(a) The sequences we consider here were introduced in [12] and are defined as follows: if $u_{i}$ is a paperfolding sequence, one defines $w_{i}$ by

$$
\begin{aligned}
w_{i}(0) & =0 \\
w_{i}(n) & =\sum_{i=0}^{n-1} u_{i}(t) \text { modulo } 2, \text { for } n \geqslant 1
\end{aligned}
$$

These sequences have the Rudin-Shapiro property that

$$
\left\|\sum_{n=0}^{N-1}(-1)^{w_{i}(n)} e^{2 i \pi n x}\right\|_{\infty} \leqslant C_{i} \sqrt{N} .
$$

We shall prove the following theorem (compare with the "other" generalised RudinShapiro sequences studied in [2]):

Theorem 2. For any generalised Rudin-Shapiro sequence (in the sense of [12]) $w_{i}$ one has:

$$
\begin{aligned}
& P_{w_{i}}(1)=2, P_{w_{i}}(2)=4, P_{w_{i}}(3)=8, P_{w_{i}}(4)=16, P_{w_{i}}(5)=24, P_{w_{i}}(6)=36, \\
& P_{w_{i}}(7)=46
\end{aligned}
$$

and for all $k \geqslant 8, P_{w_{i}}(k)=8 k-8$.
We first need a lemma:
Lemma 3. Let $u_{i}=\left(u_{i}(n)\right)_{n \geqslant 0}$ be a paperfolding sequence. If $F=u_{i}(n) u_{i}(n+1)$ $\cdots u_{i}(n+k-1)$ is a factor of $u_{i}$ of length $k$, such that $\sum_{t=0}^{n-1} u_{i}(t)=a$, then there exists $n^{\prime}$ such that:

* $\quad F=u_{i}\left(n^{\prime}\right) u_{i}\left(n^{\prime}+1\right) \cdots u_{i}\left(n^{\prime}+k-1\right)$
* $\sum_{t=0}^{n^{\prime}-1} u_{i}(t)=1+a$ modulo 2.

Here we use the definition of paperfolding by means of "perturbed symmetry". If $M$ is a factor of $u_{i}$, then there exist two factors $X$ and $Y$ such that $X M Y$ is a left factor of $u_{i}$ (that is, beginning at place 0 ) of length $2^{s}-1$ for some $s \geqslant 1$.

Now applying perturbed symmetry operators three times we see that $u_{i}$ begins with:

## $X M Y \alpha \widetilde{Y} \widetilde{M} \tilde{X} \beta X M Y \bar{\alpha} \tilde{Y} \widetilde{M} \tilde{X} \gamma X M Y \alpha \tilde{Y} \widetilde{M} \tilde{X} \bar{\beta} X M Y \bar{\alpha} \tilde{Y} \widetilde{M} \tilde{X}$

( $\bar{\alpha}$ means $1+\alpha$ modulo 2 ).
$M$ occurs four times. Denoting by $s(X)$ the sum modulo 2 of the letters of $X$ one sees that:

* the first occurrence of $M$ is preceded by a word of $\operatorname{sum} s(X)$,
* the second occurrence of $M$ is preceded by a word of sum $s(X)+$ $s(X M Y)+s(\tilde{Y} \widetilde{M} \tilde{X})+\alpha+\beta$; but for every word $Z, s(Z)+s(\tilde{Z})=$ length of $Z$ modulo 2 ; hence this sum is equal to $s(X)+1+\alpha+\beta$;
* the third occurrence of $M$ is preceded by a word of $\operatorname{sum} s(X)+1+\alpha+$ $\beta+$ length $(X M Y)+\bar{\alpha}+\gamma=s(X)+1+\beta+\gamma ;$
* the fourth occurrence of $M$ is preceded by a word of sum $s(X)+1+\beta+$ $\gamma+$ length $(X M Y)+\alpha+\bar{\beta}=s(X)+1+\alpha+\gamma$.

As one cannot simultaneously have

$$
1+\alpha+\beta=1+\beta+\gamma=1+\alpha+\gamma=0
$$

one of these sums is equal to 1 , proving the lemma.
0
Proof of Theorem 2: (b) Let $u_{i}=\left(u_{i}(n)\right)_{n \geqslant 0}$ be a paperfolding sequence with sequence of folding instructions $i$, and let $w_{i}$ be defined by:

$$
\begin{aligned}
w_{i}(0) & =0 \\
w_{i}(n) & =\sum_{t=0}^{n-1} u_{i}(t) \quad \text { for } \quad n \geqslant 1 .
\end{aligned}
$$

Let $F_{u_{i}}(k)$ be the set of factors of $u_{i}$ of length $k$ and similarly define $F_{w_{i}}(k)$. Now define on $F_{w_{i}}(k)(k \geqslant 2)$ the map $\psi_{k}$ by

$$
\psi_{k}\left(e_{0}, e_{1}, \cdots, e_{k-1}\right)=\left(e_{0}, e_{0}+e_{1}, \cdots, e_{k-2}+e_{k-1}\right)
$$

(the sums being taken modulo 2).
If $n$ is such that $w_{i}(n+t)=e_{t}$, for $0 \leqslant t \leqslant k-1$, we see that $u_{i}(n+t)=$ $w_{i}(n+t)+w_{i}(n+t+1)=e_{t}+e_{t+1}$ for $0 \leqslant t \leqslant k-2$, hence $\psi_{k}$ maps $F_{w_{i}}(k)$ to $\{0,1\} \times F_{u_{i}}(k-1)$.

Clearly $\psi_{k}$ is one-to-one. To see that $\psi_{k}$ is onto, we note that given $\left(a_{0}, a_{1}, \cdots, a_{k-1}\right)$ in $\{0,1\} \times F_{u_{i}}(k-1)$, there exists $n$ such that $u_{i}(n+t)=a_{t+1}$ for $0 \leqslant t \leqslant k-2$.

Hence:

$$
\begin{aligned}
w_{i}(n+1) & =w_{i}(n)+a_{1} \\
w_{i}(n+2) & =w_{i}(n)+a_{1}+a_{2} \\
\cdots & \\
w_{i}(n+k-1) & =w_{i}(n)+a_{1}+a_{2}+\cdots+a_{k-1} .
\end{aligned}
$$

So if $w_{i}(n)=a_{0}$, then $w_{i}(n+t)=a_{0}+a_{1}+a_{2}+\cdots+a_{t}$ for every $t$ in $[0, k-1]$, and this gives an element in $F_{w_{i}}(k)$ such that

$$
\psi_{k}\left(w_{i}(n), w_{i}(n+1), \cdots, w_{i}(n+k-1)\right)=\left(a_{0}, a_{1}, \cdots, a_{k-1}\right)
$$

If now $w_{i}(n)=1+a_{0}$, (that is $\left.\sum_{t=0}^{n-1} u_{i}(t)=1+a_{0}\right)$, then by Lemma 3 there exists an integer $n^{\prime}$ such that $\sum_{t=0}^{n^{\prime}-1} u_{i}(t)=a_{0}$, that is, $w_{i}\left(n^{\prime}\right)=a_{0}$ and $u_{i}\left(n^{\prime}+t\right)=u_{i}(n+t)=$ $a_{t+1}$ for $0 \leqslant t \leqslant k-2$.

Hence $w_{i}\left(n^{\prime}+t\right)=a_{0}+a_{1}+\cdots+a_{t}$ for $0 \leqslant t \leqslant k-1$, and

$$
\psi_{k}\left(w_{i}\left(n^{\prime}\right), w_{i}\left(n^{\prime}+1\right), \cdots, w_{i}\left(n^{\prime}+k-1\right)\right)=\left(a_{0}, a_{1}, \cdots, a_{k-1}\right) .
$$

Thus finally $\psi_{k}$ is a bijection from $F_{w_{i}}(k)$ onto $\{0,1\} \times F_{w_{i}}(k-1)$, which proves that:

$$
P_{w_{i}}(k)=2 P_{u_{i}}(k-1), \quad \text { for } \quad k \geqslant 2,
$$

and our Theorem 2 is now nothing but a reformulation of Theorem 1.

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