

UNIFORM SPANNING FORESTS OF PLANAR GRAPHS

TOM HUTCHCROFT¹ and ASAF NACHMIAS²

¹ Statslab, DPMMS, University of Cambridge, Cambridge CB3 0WB, UK;
email: t.hutchcroft@maths.cam.ac.uk

² Department of Mathematical Sciences, Tel Aviv University, Tel Aviv 6997801, Israel;
email: asafnach@tauex.tau.ac.il

Received 24 January 2018; accepted 3 March 2019

Abstract

We prove that the free uniform spanning forest of any bounded degree proper plane graph is connected almost surely, answering a question of Benjamini, Lyons, Peres and Schramm. We provide a quantitative form of this result, calculating the critical exponents governing the geometry of the uniform spanning forests of transient proper plane graphs with bounded degrees and codegrees. We find that the same exponents hold universally over this entire class of graphs provided that measurements are made using the hyperbolic geometry of their circle packings rather than their usual combinatorial geometry.

2010 Mathematics Subject Classification: 60D05 (primary); 05C10 (secondary)

1. Introduction

The **uniform spanning forests** (USFs) of an infinite, locally finite, connected graph G are defined as weak limits of uniform spanning trees (USTs) of finite subgraphs of G . These limits can be taken with either **free** or **wired** boundary conditions, yielding the **free uniform spanning forest (FUSF)** and the **wired uniform spanning forest (WUSF)**, respectively. Although the USFs are defined as limits of random spanning trees, they need not be connected. Indeed, a principal result of Pemantle [43] is that the WUSF and FUSF of \mathbb{Z}^d coincide, and that they are almost surely (a.s.) a single tree if and only if $d \leq 4$. Benjamini, Lyons, Peres and Schramm [10] (henceforth referred to as BLPS) later gave a complete characterization of connectivity of the WUSF, proving that the WUSF of a graph

© The Author(s) 2019. This is an Open Access article, distributed under the terms of the Creative Commons Attribution licence (<http://creativecommons.org/licenses/by/4.0/>), which permits unrestricted re-use, distribution, and reproduction in any medium, provided the original work is properly cited.

is connected a.s. if and only if two independent random walks on the graph intersect a.s. [10, Theorem 9.2].

The FUSF is much less understood. No characterization of its connectivity is known, nor has it been proven that connectivity is a zero-one event. One class of graphs in which the FUSF is relatively well understood are the *proper plane graphs*. Recall that a **planar graph** is a graph that can be embedded in the plane, while a **plane graph** is a planar graph G together with a specified embedding of G in the plane (or some other topological disc). A plane graph is **proper** if the embedding is proper, meaning that every compact subset of the plane (or whatever topological disc G was embedded in) intersects at most finitely many edges and vertices of the drawing (see Section 2.5 for further details). For example, every tree can be drawn in the plane without accumulation points, while the product of \mathbb{Z} with a finite cycle is planar but cannot be drawn in the plane without an accumulation point (and therefore has no proper embedding in the plane).

BLPS proved that the free and wired uniform spanning forests are distinct whenever G is a transient proper plane graph with bounded degrees, and asked [10, Question 15.2] whether the FUSF is a.s. connected in this class of graphs. They proved that this is indeed the case when G is a self-dual plane Cayley graph that is rough-isometric to the hyperbolic plane [10, Theorem 12.7]. These hypotheses were later weakened by Lyons, Morris and Schramm [37, Theorem 7.5], who proved that the FUSF of any bounded degree proper plane graph that is rough-isometric to the hyperbolic plane is a.s. connected.

Our first result provides a complete answer to [10, Question 15.2], obtaining optimal hypotheses under which the FUSF of a proper plane graph is a.s. connected. The techniques we developed to answer this question also allow us to prove quantitative versions of this result, which we describe in the next section. We state our result in the natural generality of proper plane networks. Recall that a **network** (G, c) is a locally finite, connected graph $G = (V, E)$ together with a function $c : E \rightarrow (0, \infty)$ assigning a positive **conductance** to each edge of G . The **resistance** of an edge e in a network (G, c) is defined to be $1/c(e)$. The uniform spanning forest of a finite network gives each tree probability proportional to the product of the conductances of its edges. The free and wired USFs of an infinite network are obtained, as before, by taking weak limits over exhaustions. Graphs may be considered as networks by setting $c \equiv 1$. A **plane network** is a planar graph G together with specified conductances and a specified drawing of G in the plane.

THEOREM 1.1. *The FUSF is a.s. connected in any bounded degree proper plane network with edge conductances bounded above.*

In light of the duality between the free and wired uniform spanning forests of proper plane graphs (see Section 2.5.1), the FUSF of a proper plane graph G with locally finite dual G^\dagger is connected a.s. if and only if every component of the WUSF of G^\dagger is a.s. *one-ended*. Thus, Theorem 1.1 follows easily from the dual statement Theorem 1.2. (The implication is immediate when the dual graph is locally finite.) Recall that an infinite graph $G = (V, E)$ is said to be **one-ended** if, for every finite set $K \subset V$, the subgraph induced by $V \setminus K$ has exactly one infinite connected component. In particular, an infinite tree is one-ended if and only if it does not contain a simple bi-infinite path. Components of the WUSF are known to be one-ended a.s. in several other classes of graphs [2, 10, 26, 28, 37, 43], and are recurrent in any graph [41]. Recall that a plane graph is said to have **bounded codegree** if its dual has bounded degree.

THEOREM 1.2. *Every component of the WUSF is one-ended a.s. in any bounded codegree proper plane network with edge resistances bounded above.*

The uniform spanning trees of \mathbb{Z}^2 and other planar Euclidean lattices are very well understood due to the deep theory of conformally invariant scaling limits. The study of the UST on \mathbb{Z}^2 led Schramm, in his seminal paper [48], to introduce the Schramm-Loewner Evolution (SLE) processes, which he conjectured to describe the scaling limits of the loop-erased random walk and UST. This conjecture was subsequently proven in the celebrated work of Lawler *et al.* [36]. Overall, Schramm's introduction of SLE has revolutionized the understanding of statistical physics in two dimensions; see for example, [17, 34, 46] for guides to the extensive literature in this very active field.

Although our own setting is too general to apply this theory, we nevertheless keep conformal invariance in mind throughout this paper. Indeed, the key to our proofs is *circle packing*, a canonical method of drawing planar graphs that is closely related to conformal mapping (see for example, [23, 25, 45, 46, 52] and references therein). For many purposes, one can pretend that the random walk on the packing is a quasiconformal image of standard planar Brownian motion: Effective resistances, heat kernels, and harmonic measures on the graph can each be estimated in terms of the corresponding Brownian quantities [4, 16].

It is natural to ask whether Theorem 1.1 extends to *all* bounded degree planar graphs, rather than just those admitting proper embeddings into the plane. In the sequel to this paper, we provide an example, which was analysed in collaboration with Gady Kozma, to show that the theorem does not admit such an extension. However, we also show there that the conclusion of Theorem 1.1 does continue to hold on any bounded degree planar network that admits a proper embedding into a domain with countably many boundary components.

1.1. Universal USF exponents via circle packing. A **circle packing** P is a set of discs in the Riemann sphere $\mathbb{C} \cup \{\infty\}$ that have disjoint interiors (that is, do not overlap) but can be tangent. The **tangency graph** $G = G(P)$ of a circle packing P is the plane graph with the centres of the circles in P as its vertices and with edges given by straight lines between the centres of tangent circles. The Koebe–Andreev–Thurston Circle Packing theorem [33, 39, 53] states that every *finite*, simple planar graph arises as the tangency graph of a circle packing, and that if the graph is a triangulation then its circle packing is unique up to Möbius transformations and reflections (see [14] for a combinatorial proof). The Circle Packing theorem was extended to infinite plane triangulations by He and Schramm [23, 24], who proved that every *infinite, proper*, simple plane triangulation admits a locally finite circle packing in either the Euclidean plane or the hyperbolic plane (identified with the interior of the unit disc), but not both. (Recall that in the Poincaré disc model, Euclidean and hyperbolic circles in the disc coincide as sets but may have different centres and radii.) We call an infinite, simple, proper plane triangulation **CP parabolic** if it admits a circle packing in the plane and **CP hyperbolic** otherwise.

He and Schramm [24] also initiated the use of circle packing to study probabilistic questions on plane graphs. In particular, they showed that a bounded degree, simple, proper plane triangulation is CP parabolic if and only if simple random walk on the triangulation is recurrent (that is, visits every vertex infinitely often a.s.). Circle packing has since proven instrumental in the study of planar graphs, and random walks on planar graphs in particular. Most relevantly to us, circle packing was used by Benjamini and Schramm [11] to prove that every transient, bounded degree planar graph admits nonconstant harmonic Dirichlet functions; BLPS [10] later applied this result to deduce that the free and wired uniform spanning forest of a bounded degree plane graph coincide if and only if the graph is recurrent. We refer the reader to [52] and [46] for background on circle packing, and to [4, 5, 7, 11, 12, 16, 18, 19, 24, 27, 29, 30, 42] for further probabilistic applications.

A guiding principle of the works mentioned above is that circle packing endows a triangulation with a geometry that, for many purposes, is better than the usual graph metric. The results described in this section provide a compelling instance of this principle in action: we find that the critical exponents governing the geometry of the USFs are universal over all transient, bounded degree, proper plane triangulations, provided that measurements are made using the hyperbolic geometry of their circle packings rather than the usual combinatorial geometry of the graphs. It is crucial here that we use the circle packing to take our measurements: The exponents in the graph distance are not universal and need

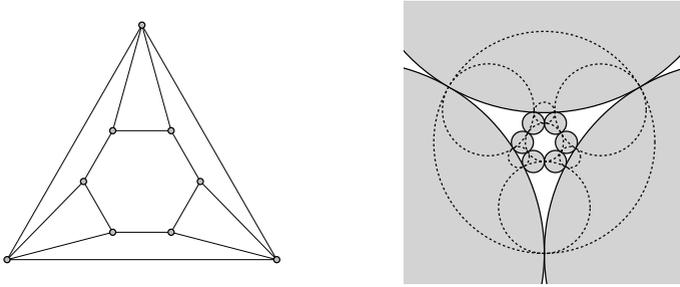


Figure 1. A finite polyhedral plane graph (left) and its double circle packing (right). Primal circles are filled and have solid boundaries, dual circles have dashed boundaries.

not even exist (see Figure 2). In Remark 5.2 we give an example to show that no such universal exponents hold in the CP parabolic case at this level of generality.

In order to state the quantitative versions of Theorems 1.1 and 1.2 in their full generality, we first introduce **double circle packing**. Let G be a plane graph with vertex set V and face set F . A double circle packing of G is a pair of circle packings $P = \{P(v) : v \in V\}$ and $P^\dagger = \{P^\dagger(f) : f \in F\}$ in $\mathbb{C} \cup \{\infty\}$ satisfying the following conditions (see Figure 1):

- (1) (**G is the tangency graph of P .**) For each pair of vertices u and v of G , the discs $P(u)$ and $P(v)$ are tangent if and only if u and v are adjacent in G .
- (2) (**G^\dagger is the tangency graph of P^\dagger .**) For each pair of faces f and g of G , the discs $P^\dagger(f)$ and $P^\dagger(g)$ are tangent if and only if f and g are adjacent in G^\dagger .
- (3) (**Primal and dual circles are perpendicular.**) For each vertex v and face f of G , the discs $P^\dagger(f)$ and $P(v)$ have nonempty intersection if and only if f is incident to v , and in this case the boundary circles of $P^\dagger(f)$ and $P(v)$ intersect at right angles.

It is easily seen that any finite plane graph admitting a double circle packing must be **polyhedral**, meaning that it is both simple (that is, not containing any loops or multiple edges) and 3-connected (meaning that the subgraph induced by $V \setminus \{u, v\}$ is connected for each $u, v \in V$). Conversely, Thurston's interpretation of Andreev's theorem [39, 53], or the Brightwell–Scheinerman theorem [14], implies that every finite, polyhedral plane graph admits a double circle packing. The corresponding infinite theory was developed by He [22], who proved that every infinite, polyhedral, proper plane graph G with locally finite dual admits a

double circle packing in either the Euclidean plane or the hyperbolic plane, but not both, and that this packing is unique up to Möbius transformations and reflections. (He phrased his results more generally than this, see Section 2.6 for an explanation of how the statement given here follows from the one in [22].) In particular, in the hyperbolic case, the packing is unique up to isometries of the hyperbolic plane. As before, we say that G is CP parabolic or CP hyperbolic as appropriate. As for triangulations [24], CP hyperbolicity is equivalent to transience for graphs with bounded degrees and codegrees [22].

Let G be CP hyperbolic and let (P, P^\dagger) be a double circle packing of G in the hyperbolic plane. We write $r_{\mathbb{H}}(v)$ for the hyperbolic radius of the circle $P(v)$. For each subset $A \subset V(G)$, we define $\text{diam}_{\mathbb{H}}(A)$ to be the hyperbolic diameter of the set of hyperbolic centres of the circles in P corresponding to vertices in A , and define $\text{area}_{\mathbb{H}}(A)$ to be the hyperbolic area of the union of the circles in P corresponding to vertices in A . Since (P, P^\dagger) is unique up to isometries of the hyperbolic plane, $r_{\mathbb{H}}(v)$, $\text{diam}_{\mathbb{H}}(A)$, and $\text{area}_{\mathbb{H}}(A)$ do not depend on the choice of (P, P^\dagger) .

We say that a network G has **bounded local geometry** if there exists a constant M such that $\deg(v) \leq M$ for every vertex v of G and $M^{-1} \leq c(e) \leq M$ for every edge e of G . Given a plane network G , we let F be the set of faces of G , and define

$$\mathbf{M} = \mathbf{M}_G = \max \left\{ \sup_{v \in V} \deg(v), \sup_{f \in F} \deg(f), \sup_{e \in E} c(e), \sup_{e \in E} c(e)^{-1} \right\}.$$

Given a spanning tree \mathfrak{F} and two vertices x and y in G , we write $\Gamma_{\mathfrak{F}}(x, y)$ for the unique path in \mathfrak{F} connecting x and y . In the following theorem, the graph G satisfies the conditions of Theorem 1.1 and so the path $\Gamma_{\mathfrak{F}}$ is well defined a.s.

THEOREM 1.3 (Free diameter exponent). *Let G be a transient, polyhedral, proper plane network with $\mathbf{M}_G < \infty$, let \mathfrak{F} be the FUSF of G , and let $e = (x, y)$ be an edge of G . Then there exist positive constants*

$$k_1 = k_1(\mathbf{M}, r_{\mathbb{H}}(x)),$$

increasing in $r_{\mathbb{H}}(x)$, and $k_2 = k_2(\mathbf{M})$ such that

$$k_1 R^{-1} \leq \mathbb{P}(\text{diam}_{\mathbb{H}}(\Gamma_{\mathfrak{F}}(x, y)) \geq R) \leq k_2 R^{-1}$$

for every $R \geq 1$.

Given a spanning forest \mathfrak{F} of G in which every component is a one-ended tree and an edge e of G , the **past** of e in \mathfrak{F} , denoted $\text{past}_{\mathfrak{F}}(e)$, is defined to be the

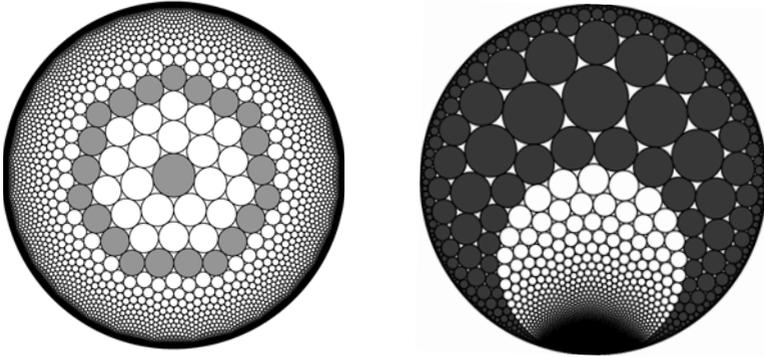


Figure 2. Two bounded degree, simple, proper plane triangulations for which the graph distance is not comparable to the hyperbolic distance. Similar examples are given in [52, Figure 17.7]. Left: In this example, rings of degree seven vertices (grey) are separated by growing bands of degree six vertices (white), causing the hyperbolic radii of circles to decay. The bands of degree six vertices can grow surprisingly quickly without the triangulation becoming recurrent [50]. Right: In this example, half-spaces of the 8-regular (grey) and 6-regular (white) triangulations have been glued together along their boundaries; the circles corresponding to the 6-regular half-space are contained inside a horodisc and have decaying hyperbolic radii.

unique finite connected component of $\mathfrak{F} \setminus \{e\}$ if $e \in \mathfrak{F}$ and to be the empty set otherwise. The following theorem is equivalent to Theorem 1.3 by duality (see Sections 2.5.1 and 4.6).

THEOREM 1.4 (Wired diameter exponent). *Let G be a transient, polyhedral, proper plane network with $\mathbf{M}_G < \infty$, let \mathfrak{F} be the WUSF of G , and let $e = (x, y)$ be an edge of G . Then there exist positive constants $k_1 = k_1(\mathbf{M}, r_{\mathbb{H}}(x))$, increasing in $r_{\mathbb{H}}(x)$, and $k_2 = k_2(\mathbf{M})$ such that*

$$k_1 R^{-1} \leq \mathbb{P}(\text{diam}_{\mathbb{H}}(\text{past}_{\mathfrak{F}}(e)) \geq R) \leq k_2 R^{-1}$$

for every $R \geq 1$.

By similar methods, we are also able to obtain a universal exponent of $1/2$ for the tail of the area of the past of an edge in the WUSF.

THEOREM 1.5 (Wired area exponent). *Let G be a transient, polyhedral, proper plane network with $\mathbf{M}_G < \infty$, let \mathfrak{F} be the WUSF of G , and let $e = (x, y)$ be an edge of G . Then there exist positive constants*

$$k_1 = k_1(\mathbf{M}, r_{\mathbb{H}}(x)),$$

increasing in $r_{\mathbb{H}}(x)$, and $k_2 = k_2(\mathbf{M})$ such that

$$k_1 R^{-1/2} \leq \mathbb{P}(\text{area}_{\mathbb{H}}(\text{past}_{\mathfrak{F}}(e)) \geq R) \leq k_2 R^{-1/2}$$

for every $R \geq 1$.

The exponents 1 and 1/2 occurring in Theorems 1.4 and 1.5 should, respectively, be compared with the analogous exponents for the survival time and total progeny of a critical branching process whose offspring distribution has finite variance (see for example, [38, §5,12]).

For *uniformly transient* proper plane graphs (that is, proper plane graphs in which the escape probabilities of the random walk are bounded uniformly away from zero), the hyperbolic radii of circles in (P, P^\dagger) are bounded away from zero uniformly (Proposition 4.16). This implies that the hyperbolic metric and the graph metric are rough-isometric (Corollary 4.17). Consequently, Theorems 1.3–1.5 hold with the graph distance and counting measure as appropriate (Corollary 4.18). This yields the following appealing corollary, which applies, for example, to planar Cayley graphs of cocompact Fuchsian groups.

COROLLARY 1.6 (Free length exponent). *Let G be a uniformly transient, polyhedral, proper plane network with $\mathbf{M}_G < \infty$ and let \mathfrak{F} be the FUSF of G . Let $\mathbf{p} > 0$ be a uniform lower bound on the escape probabilities of G . Then there exist positive constants $k_1 = k_1(\mathbf{M}, \mathbf{p})$ and $k_2 = k_2(\mathbf{M}, \mathbf{p})$ such that*

$$k_1 R^{-1/2} \leq \mathbb{P}(|\Gamma_{\mathfrak{F}}(x, y)| \geq R) \leq k_2 R^{-1/2}$$

for every edge $e = (x, y)$ of G and every $R \geq 1$.

See [9, 13, 31, 37, 40, 49] and the survey [8] for related results on Euclidean lattices.

1.2. Rough outline of the proof. The proof of Theorem 1.2 goes by showing that for every edge $e = (x, y)$ of the graph, the conditional probability that e is in the WUSF conditioned on having two long disjoint WUSF paths, one starting

at x and the other starting at y , is small as the length of the paths grows. Taking a limit one gets (7) from which we obtain that a.s. no bi-infinite paths in the WUSF exist. To prove this we take the rightmost path from x and the left most path from y and use the circle packing embedding to infer that no matter what these paths are, a large quantity of electric current can flow from one path to the other (with voltage difference 1) within the space that has not been explored by the WUSF. This current is also allowed to flow ‘via infinity’ since the boundary is wired, which handles the case where the space between the two paths is large.

It turns out that this approach already yields the correct order of the upper bound in Theorem 1.4. The lower bound, proved in Section 4, is significantly more complicated. The intuition is as follows: If we consider the intersection of the past of the origin in the WUSF with the hyperbolic shell $B(0, r + 1) \setminus B(0, r)$, the hyperbolic area of this intersection behaves similarly to the number of particles at generation n in a critical (finite variance) branching process; See (25) and below. The bulk of the work in Section 4 is to bound the first and second moments of the hyperbolic areas of these intersections. The bounds we obtain are analogous to the fact that the n th generation of a critical (finite variance) branching process has first moment 1 and second moment of order n .

Organization. Section 2 contains definitions and background on those notions that will be used throughout the paper. Experienced readers are advised that this section also includes proofs of a few simple folklore-type lemmas and propositions. Section 3 contains the proofs of Theorems 1.1 and 1.2, as well as the upper bound of Theorem 1.4 in the case that G is a triangulation. Section 4 completes the proofs of Theorems 1.3–1.5 and Corollary 1.6; the most substantial component of this section is the proof of the lower bound of Theorem 1.4. Section 4 also includes the statement and proof of an extension of the Ring lemma to double circle packings. We conclude with a remark and an open problem in Section 5.

2. Background and definitions

2.1. Notation. We write E^\rightarrow for the set of oriented edges of a network $G = (V, E)$. An oriented edge $e \in E^\rightarrow$ is oriented from its tail e^- to its head e^+ , and has reversal $-e$.

For each $r', r > 0$, and $z \in \mathbb{C}$, we define the open ball and closed annulus

$$B_z(r) = \{z' \in \mathbb{C} : |z - z'| < r\} \quad \text{and} \quad A_z(r, r') = \{z' \in \mathbb{C} : r \leq |z' - z| \leq r'\}.$$

We write $d_{\mathbb{C}}$ for the Euclidean metric on \mathbb{C} and write $d_{\mathbb{H}}$ for the hyperbolic metric on \mathbb{D} .

2.2. Hyperbolic geometry. Let \mathbb{D} be the open unit disc. The hyperbolic metric $d_{\mathbb{H}}$ on \mathbb{D} is defined by

$$d_{\mathbb{H}}(z_1, z_2) = 2 \tanh^{-1} \left| \frac{z_1 - z_2}{1 - z_1 \bar{z}_2} \right|.$$

Some basic facts we will make frequent use of include:

- (1) Möbius transformations and reflections fixing \mathbb{D} are isometries of the hyperbolic plane.
- (2) Euclidean circles in \mathbb{D} are hyperbolic circles with different centres and radii.
- (3) For every $a < \infty$ there exists A such that if $C \subseteq \mathbb{D}$ is a circle with hyperbolic radius $r_{\mathbb{H}} \leq a$ and Euclidean centre z and radius r , then

$$\frac{r}{A(1 - |z|)} \leq r_{\mathbb{H}} = 2 \tanh^{-1} \left| \frac{2r}{1 + r^2 - |z|^2} \right| \leq \frac{Ar}{1 - |z|}.$$

See for example, [3] for further background on hyperbolic geometry.

2.3. Uniform Spanning Forests. We begin with a succinct review of some basic facts about USFs, referring the reader to [10] and [38, Chs 4 and 10] for a comprehensive overview. For each finite, connected graph G , we define UST_G to be the uniform measure on the set of spanning trees of G (that is, connected subgraphs of G containing every vertex and no cycles). More generally, for a finite network $G = (G, w)$, we define UST_G to be the probability measure on spanning trees of G for which the measure of each tree is proportional to the product of the conductances of the edges in the tree.

An **exhaustion** of an infinite network G is a sequence $\langle V_n \rangle_{n \geq 0}$ of finite, connected subsets of V such that $V_n \subseteq V_{n+1}$ for all $n \geq 0$ and $\bigcup_n V_n = V$. Given such an exhaustion, let the network G_n be the subgraph of G induced by V_n together with the conductances inherited from G . The **free uniform spanning forest** measure FUSF_G is defined to be the weak limit of the sequence $\langle \text{UST}_{G_n} \rangle_{n \geq 1}$, so that

$$\text{FUSF}_G(S \subset \mathfrak{F}) = \lim_{n \rightarrow \infty} \text{UST}_{G_n}(S \subset T)$$

for each finite set $S \subset E$, where \mathfrak{F} is a sample of FUSF_G and T is a sample of UST_{G_n} . For each n , we also construct a network G_n^* from G by gluing (wiring) every vertex of $G \setminus G_n$ into a single vertex, denoted ∂_n , and deleting all the self-loops that are created. We identify the set of edges of G_n^* with the set of edges

of G that have at least one endpoint in V_n . The **wired uniform spanning forest** measure WUSF_G is defined to be the weak limit of the sequence $\langle \text{UST}_{G_n^*} \rangle_{n \geq 1}$, so that

$$\text{WUSF}_G(S \subset \mathfrak{F}) = \lim_{n \rightarrow \infty} \text{UST}_{G_n^*}(S \subset T)$$

for each finite set $S \subset E$, where \mathfrak{F} is a sample of WUSF_G and T is a sample of $\text{UST}_{G_n^*}$. These weak limits were both implicitly proven to exist by Pemantle [43], although the WUSF was not considered explicitly until the work of Häggström [20]. Both measures are easily seen to be concentrated on the set of **essential spanning forests** of G , that is, cycle-free subgraphs of G including every vertex such that every connected component is infinite. The measure FUSF_G stochastically dominates WUSF_G for every infinite network G .

2.3.1. The spatial Markov property. Let $G = (V, E)$ be a finite or infinite network and let A and B be subsets of E . We write $(G - B)/A$ for the (possibly disconnected) network formed from G by deleting every edge in B and contracting (that is, identifying the two endpoints of) every edge in A . We identify the edges of $(G - B)/A$ with $E \setminus B$. Suppose that G is finite, and that

$$\text{UST}_G(A \subseteq T, B \cap T = \emptyset) > 0.$$

Then, given the event that T contains every edge in A and none of the edges in B , the conditional distribution of T is equal to the union of A and the UST of $(G - B)/A$. That is, for every event $\mathcal{A} \subseteq \{0, 1\}^E$,

$$\text{UST}_G(T \in \mathcal{A} \mid A \subset T, B \cap T = \emptyset) = \text{UST}_{(G-B)/A}(T \cup A \in \mathcal{A}).$$

This is the **spatial Markov property** of the UST. Taking limits over exhaustions, we obtain a corresponding spatial Markov property for the USFs: if $G = (V, E)$ is an infinite network and A and B are subsets of E such that

$$\text{WUSF}_G(A \subseteq \mathfrak{F}, B \cap \mathfrak{F} = \emptyset) > 0$$

and the network $(G - B)/A$ is locally finite, then

$$\text{WUSF}_G(\mathfrak{F} \in \mathcal{A} \mid A \subset \mathfrak{F}, B \cap \mathfrak{F} = \emptyset) = \text{WUSF}_{(G-B)/A}(\mathfrak{F} \cup A \in \mathcal{A}). \quad (1)$$

(If $(G - B)/A$ is not connected, we define $\text{WUSF}_{(G-B)/A}$ to be the product of the WUSF measures on the connected components of $(G - B)/A$.) A similar spatial Markov property holds for the FUSF.

The UST and USFs also enjoy a *strong* form of the spatial Markov property. Let G be a finite or infinite network, and suppose that H is a random element of

$\{0, 1\}^E$, which we think of as a random subgraph of G . We say that a random element K of $\{0, 1\}^E$, defined on the same probability space as H , is a **local set for H** if for every set $W \subseteq E$, the event $\{K \subseteq W\}$ is, up to a null set, measurable with respect to the σ -algebra generated by the collection of random variables $\{H(e) : e \in W\}$.

For example, suppose that \mathfrak{F} is a random spanning forest of a network G , and that u and v are vertices of G . Let $K = K(\mathfrak{F})$ be the random set of edges that is equal to the path connecting u and v in \mathfrak{F} if such a path exists, and otherwise equal to all of E . Then K is a local set for \mathfrak{F} , since, for every proper subset W of E , the event $\{K \subseteq W\}$ is equal (up to the null event in which \mathfrak{F} is not a forest) to the event that u and v are in the same component of \mathfrak{F} and that the unique path connecting u and v in \mathfrak{F} is contained in W , which is clearly measurable with respect to the restriction of \mathfrak{F} to W .

An example of a random set K that is not a local set for \mathfrak{F} is, for instance, the edges of \mathfrak{F} incident to a given vertex v . If v has degree at least 2 and e is some edge incident to v , then the event $\{K \subset \{e\}\}$ cannot be determined just by knowing $\mathfrak{F}(e)$. Note also that Proposition 2.1 does not hold for this K , since if we know which edges are included in K then we also know that $\mathfrak{F}(e) = 0$ for all edges $e \notin K$ incident to v .

Given a random subgraph H of a network G and a local set K for H , we write \mathcal{F}_K to denote the σ -algebra generated by K and the random collection of random variables $\{H(e) : e \in K\}$. We also write $K_o = \{e \in E : e \in K, H(e) = 1\}$ and $K_c = K \setminus K_o$ for the sets of edges in K that are included (open) in H and not included (closed) in H , respectively.

PROPOSITION 2.1 (Strong Spatial Markov Property for the WUSF). *Let $G = (V, E)$ be an infinite network, let \mathfrak{F} be a sample of the WUSF of G , and let K be a local set for \mathfrak{F} that is either finite or equal to all of E a.s. Conditional on \mathcal{F}_K and the event that K is finite, let $\hat{\mathfrak{F}}$ be a sample of the WUSF of $(G - K_c)/K_o$. Let $\mathfrak{F}' = K_o \cup \hat{\mathfrak{F}}$ if K_o is finite and $\mathfrak{F}' = \mathfrak{F}$ if $K_o = E$. Then \mathfrak{F} and \mathfrak{F}' have the same distribution.*

Proof. Expand the conditional probability

$$\begin{aligned} & \text{WUSF}_G(\mathfrak{F} \in \mathcal{A} \mid \mathcal{F}_K) \mathbb{1}(K \neq E) \\ &= \sum_{W_1, W_2 \subseteq E \text{ finite}} \text{WUSF}_G(\mathfrak{F} \in \mathcal{A} \mid K_o = W_1, K_c = W_2) \mathbb{1}(K_o = W_1, K_c = W_2). \end{aligned}$$

Since K is a local set for \mathfrak{F} , the right-hand side is equal to

$$\sum_{W_1, W_2 \subseteq E \text{ finite}} \text{WUSF}_G(\mathfrak{F} \in \mathcal{A} \mid W_1 \subseteq \mathfrak{F}, W_2 \cap \mathfrak{F} = \emptyset) \mathbb{1}(K_o = W_1, K_c = W_2).$$

Applying the spatial Markov property gives

$$\begin{aligned} & \text{WUSF}_G(\mathfrak{F} \in \mathcal{A} \mid \mathcal{F}_K)\mathbb{1}(K \neq E) \\ &= \sum_{W_1, W_2 \subset E \text{ finite}} \text{WUSF}_{(G-W_2)/W_1}(\hat{\mathfrak{F}} \cup W_1 \in \mathcal{A})\mathbb{1}(K_o = W_1, K_c = W_2) \\ &= \text{WUSF}_{(G-K_c)/K_o}(\hat{\mathfrak{F}} \cup K_o \in \mathcal{A})\mathbb{1}(K \neq E), \end{aligned}$$

from which the claim follows immediately. \square

Similar strong spatial Markov properties hold for the FUSF and UST, and admit very similar proofs.

2.4. Random walk, effective resistances. Given a network G and a vertex u of G , we write \mathbf{P}_u^G for the law of the simple random walk on G started at u , and will often write simply \mathbf{P}_u if the choice of network is unambiguous. For each set of vertices A , we let τ_A be the first time the random walk visits A , letting $\tau_A = \infty$ if the walk never visits A . Similarly, τ_A^+ is defined to be the first positive time that the random walk visits A . The conductance $c(u)$ of a vertex u is defined to be the sum of the conductances of the edges emanating from u .

Let A and B be sets of vertices in a finite network G . The **effective conductance** between A and B in G is defined to be

$$\mathcal{C}_{\text{eff}}(A \leftrightarrow B; G) = \sum_{v \in A} c(v)\mathbf{P}_v(\tau_B < \tau_A^+),$$

while the effective resistance $\mathcal{R}_{\text{eff}}(A \leftrightarrow B; G)$ is defined to be the reciprocal of the effective conductance, $\mathcal{R}_{\text{eff}}(A \leftrightarrow B; G) = \mathcal{C}_{\text{eff}}(A \leftrightarrow B; G)^{-1}$. Now suppose that G is an infinite network with exhaustion $\langle V_n \rangle_{n \geq 0}$ and let A and B be finite subsets of V . Let $\langle G_n \rangle_{n \geq 0}$ and $\langle G_n^* \rangle_{n \geq 0}$ be defined as in Section 2.3. The **free effective resistance between A and B** is defined to be the limit

$$\mathcal{R}_{\text{eff}}^{\text{F}}(A \leftrightarrow B; G) = \lim_{n \rightarrow \infty} \mathcal{R}_{\text{eff}}(A \leftrightarrow B; G_n),$$

while the **wired effective resistance between A and B** is defined to be

$$\mathcal{R}_{\text{eff}}^{\text{W}}(A \leftrightarrow B; G) = \lim_{n \rightarrow \infty} \mathcal{R}_{\text{eff}}(A \leftrightarrow B; G_n^*).$$

Free and wired effective conductances are defined by taking reciprocals. The free effective resistance between two, possibly infinite, sets A and B is defined to be the limit of the free effective resistances between $A \cap V_n$ and $B \cap V_n$, which are decreasing in n . We also define

$$\mathcal{R}_{\text{eff}}(A \leftrightarrow B \cup \{\infty\}) = \lim_{n \rightarrow \infty} \mathcal{R}_{\text{eff}}(A \leftrightarrow B \cup \{\partial_n\}; G_n^*).$$

See for example, [38] for further background on electrical networks.

We will make frequent use of **Rayleigh monotonicity principle** [38, Ch. 2.4], which states that the effective conductance between any two sets in a network is an increasing function of the edge conductances. In particular, it follows that the effective conductance between two sets decreases when edges are deleted from the network (which corresponds to taking the conductance of those edges to zero), and increases when edges are contracted (which corresponds to taking the conductance of those edges to infinity).

The proof of Theorem 1.2 will require the following simple lemma.

LEMMA 2.2. *Let A and B be sets of vertices in an infinite network G . Then*

$$\mathcal{R}_{\text{eff}}^{\text{W}}(A \leftrightarrow B; G) \leq 3 \max\{\mathcal{R}_{\text{eff}}(A \leftrightarrow B \cup \{\infty\}; G), \mathcal{R}_{\text{eff}}(B \leftrightarrow A \cup \{\infty\}; G)\}.$$

Proof. Let $M = \max\{\mathcal{R}_{\text{eff}}(A \leftrightarrow B \cup \{\infty\}; G), \mathcal{R}_{\text{eff}}(B \leftrightarrow A \cup \{\infty\}; G)\}$. Recall that for any three sets A, B and C in G_n^* (or any other finite network) [38, Exercise 2.33],

$$\mathcal{R}_{\text{eff}}(A \leftrightarrow B \cup C; G_n^*)^{-1} \leq \mathcal{R}_{\text{eff}}(A \leftrightarrow B; G_n^*)^{-1} + \mathcal{R}_{\text{eff}}(A \leftrightarrow C; G_n^*)^{-1}.$$

Letting $C = \{\partial_n\}$ and taking the limit as $n \rightarrow \infty$, we obtain that

$$M^{-1} \leq \mathcal{R}_{\text{eff}}(A \leftrightarrow B \cup \{\infty\}; G)^{-1} \leq \mathcal{R}_{\text{eff}}^{\text{W}}(A \leftrightarrow B; G)^{-1} + \mathcal{R}_{\text{eff}}(A \leftrightarrow \infty; G)^{-1}.$$

Rearranging, we have that

$$\mathcal{R}_{\text{eff}}(A \leftrightarrow \infty; G) \leq \frac{M \mathcal{R}_{\text{eff}}^{\text{W}}(A \leftrightarrow B; G)}{\mathcal{R}_{\text{eff}}^{\text{W}}(A \leftrightarrow B; G) - M}.$$

By symmetry, the inequality continues to hold when we exchange the roles of A and B . Combining both of these inequalities with the triangle inequality for effective resistances [38, Exercise 9.29] yields that

$$\begin{aligned} \mathcal{R}_{\text{eff}}^{\text{W}}(A \leftrightarrow B; G) &\leq \mathcal{R}_{\text{eff}}(A \leftrightarrow \infty; G) + \mathcal{R}_{\text{eff}}(B \leftrightarrow \infty; G) \\ &\leq 2 \frac{M \mathcal{R}_{\text{eff}}^{\text{W}}(A \leftrightarrow B; G)}{\mathcal{R}_{\text{eff}}^{\text{W}}(A \leftrightarrow B; G) - M}, \end{aligned}$$

which rearranges to give the claimed inequality. □

2.4.1. *Kirchhoff’s effective resistance formula.* The connection between effective resistances and spanning trees was first discovered by Kirchhoff [32] (see also [15]).

THEOREM 2.3 (Kirchhoff's Effective Resistance Formula). *Let G be a finite network. Then for every $e \in E^\rightarrow$*

$$\text{UST}_G(e \in T) = c(e)\mathcal{R}_{\text{eff}}(e^- \leftrightarrow e^+; G).$$

Taking limits over exhaustions, we also have the following extension of Kirchhoff's formula:

$$\text{WUSF}_G(e \in \mathfrak{F}) = c(e)\mathcal{R}_{\text{eff}}^{\text{W}}(e^- \leftrightarrow e^+; G). \quad (2)$$

A similar equality holds for the FUSF.

2.4.2. The method of random paths. We say that a path Γ in a network G is **simple** if it does not visit any vertex more than once. Given a probability measure ν on simple paths Γ in a network G , we define the **energy** of ν to be

$$\mathcal{E}(\nu) = \frac{1}{2} \sum_{e \in E^\rightarrow} \frac{1}{c(e)} (\nu(e \in \Gamma) - \nu(-e \in \Gamma))^2.$$

Effective resistances can be bounded from above by energies of random paths: In particular, if G is an infinite network and A and B are two finite sets of vertices in G , then

$$\mathcal{R}_{\text{eff}}(A \leftrightarrow B \cup \{\infty\}; G) = \min \left\{ \begin{array}{l} \nu \text{ a probability measure on simple} \\ \mathcal{E}(\nu) : \text{paths in } G \text{ starting in } A \text{ that are} \\ \text{either infinite or finite and end in } B. \end{array} \right\}$$

See for example, [38, Section 3] and [44] for more detail. Obtaining resistance bounds by defining flows using random paths in this manner is referred to as the **method of random paths**.

It will be convenient to use the following weakening of the method of random paths. Given the law μ of a random subset $W \subset V(G)$, define

$$\mathcal{E}(\mu) = \sum_{v \in V} \mu(v \in W)^2. \quad (3)$$

LEMMA 2.4 (Method of random sets). *Let A and B be two finite sets of vertices in an infinite network G , and let μ be a measure on subsets $W \subset V(G)$ such that the subgraph of G induced by V a.s. contains a path starting at A that is either infinite or finite and ends at B . Then*

$$\mathcal{R}_{\text{eff}}(A \leftrightarrow B \cup \{\infty\}; G) \leq \sup_{e \in E} c(e)^{-1} \mathcal{E}(\mu). \quad (4)$$

Proof. Given W , let Γ be a simple path connecting A to B that is contained in W . Then, letting ν be the law of Γ ,

$$\mathcal{E}(\nu) \leq \sup_e \frac{1}{c(e)} \sum_{e \in E \rightarrow} \nu(e \in \Gamma)^2.$$

Letting Γ' be an independent random path with the same law as Γ , the sum above is exactly the expected number of oriented edges that are used by both Γ and Γ' . Since these paths are simple, they each contain at most one oriented edge emanating from v for each vertex $v \in V$. It follows that the number of oriented edges included in both paths is at most the number of vertices included in both paths. This yields that

$$\mathcal{E}(\nu) \leq \sup_{e \in E} \frac{1}{c(e)} \sum_{v \in V} \nu(v \in \Gamma)^2 \leq \sup_{e \in E} \frac{1}{c(e)} \sum_{v \in V} \mu(v \in W)^2 = \sup_{e \in E} \frac{1}{c(e)} \mathcal{E}(\mu). \quad \square$$

2.5. Plane graphs and their USFs. Given a graph $G = (V, E)$, let \mathbb{G} be the metric space defined as follows. For each edge e of G , choose an orientation of e arbitrarily and let $\{I(e) : e \in E\}$ be a set of disjoint isometric copies of the interval $[0, 1]$. The metric space \mathbb{G} is defined as a quotient of the union $\bigcup_e I(e) \cup V$, where we identify the endpoints of $I(e)$ with the vertices e^- and e^+ , respectively, and is equipped with the path metric.

Let S be an orientable surface without boundary, which in this paper will always be a domain $D \subseteq \mathbb{C} \cup \{\infty\}$. A **proper embedding** of a graph G into S is a continuous, injective map $z : \mathbb{G} \rightarrow S$ satisfying the following conditions:

- (1) **(Every face is a topological disc.)** Every connected component of the complement $S \setminus z(\mathbb{G})$, called a **face** of (G, z) , is homeomorphic to the disc. Moreover, for each connected component U of $S \setminus z(\mathbb{G})$, the set of oriented edges of G that have their left-hand side incident to U forms either a cycle or a bi-infinite path in G .
- (2) **(z is locally finite.)** Every compact subset of S intersects at most finitely many edges of $z(\mathbb{G})$. Equivalently, the preimage $z^{-1}(K)$ of every compact set $K \subseteq S$ is compact in \mathbb{G} .

A locally finite, connected graph is planar if and only if it admits a proper embedding into some domain $D \subseteq \mathbb{C} \cup \{\infty\}$. A **plane graph** $G = (G, z)$ is a planar graph G together with a specified embedding $z : \mathbb{G} \rightarrow D \subseteq \mathbb{C} \cup \{\infty\}$. A **plane network** $G = (G, z, c)$ is a planar graph together with a specified embedding and an assignment of positive conductances $c : E \rightarrow (0, \infty)$.

Given a pair $G = (G, z)$ of a graph together with a proper embedding z of G into a domain D , the **dual** G^\dagger of G is the graph that has the faces of G as vertices, and has an edge drawn between two faces of G for each edge incident to both of the faces in G . By drawing each vertex of G^\dagger in the interior of the corresponding face of G and each edge of G^\dagger so that it crosses the corresponding edge of G but no others, we obtain an embedding z^\dagger of G^\dagger in D . The edge sets of G and G^\dagger are in natural correspondence, and we write e^\dagger for the edge of G^\dagger corresponding to e . If e is oriented, we let e^\dagger be oriented so that it crosses e from right to left as viewed from the orientation of e . If $G = (G, z, c)$ is a plane network, we assign the conductances $c^\dagger(e^\dagger) = c(e)^{-1}$ to the edges of G^\dagger .

2.5.1. USF duality. Let G be a plane network with dual G^\dagger . For each set $W \subseteq E$, let $W^\dagger := \{e^\dagger : e \notin W\}$. If G is finite and t is a spanning tree of G , then t^\dagger is a spanning tree of G^\dagger : the subgraph t^\dagger is connected because t has no cycles, and has no cycles because t is connected. Moreover, the ratio

$$\frac{\prod_{e \in t} c(e)}{\prod_{e^\dagger \in t^\dagger} c^\dagger(e^\dagger)} = \prod_{e \in E} c(e)$$

does not depend on t . It follows that if T is a random spanning tree of G with law UST_G , then T^\dagger is a random spanning tree of G^\dagger with law UST_{G^\dagger} . This duality was extended to infinite proper plane networks by BLPS.

THEOREM 2.5 [10, Theorem 12.2 and Proposition 12.5]. *Suppose that G is an infinite proper plane network with locally finite dual G^\dagger . Then if \mathfrak{F} is a sample of FUSF_G , the subgraph \mathfrak{F}^\dagger is an essential spanning forest of G^\dagger with law WUSF_{G^\dagger} . In particular, \mathfrak{F} is connected a.s. if and only if every component of \mathfrak{F}^\dagger is one-ended a.s.*

2.6. Circle packing. We now give some background on circle packing. The **carrier** of a circle packing P , denoted $\text{carr}(P)$, is the union of all the discs of P and of all the faces of $G(P)$, so that the embedding z of $G(P)$ defined by drawing straight lines between the centres of tangent circles is a proper embedding of $G(P)$ into $\text{carr}(P)$. Similarly, the carrier of a double circle packing (P, P^\dagger) is defined to be the union of all the discs in $P \cup P^\dagger$. Given a domain $D \subset \mathbb{C} \cup \{\infty\}$, a circle packing P (or double circle packing (P, P^\dagger)) is said to be **in** D if its carrier is D . In particular, a (double) circle packing P is said to be **in the plane** if its carrier is the plane \mathbb{C} and **in the disc** if its carrier is the open unit disc \mathbb{D} . The following theorems, which we stated in the introduction, are the cornerstones of the theory for infinite proper plane graphs.

THEOREM 2.6 (He–Schramm existence and uniqueness theorem [22–24, 47]). *Let G be an infinite, polyhedral, proper plane graph with locally finite dual. Then G admits a double circle packing either in the plane or the disc, but not both, and this packing is unique up to Möbius transformations and reflections of the plane or the disc as appropriate.*

THEOREM 2.7 (He–Schramm recurrence theorem [22, 24]). *Let G be an infinite, polyhedral, proper plane graph with bounded degrees and locally finite dual. Then G is CP parabolic if and only if it is recurrent for simple random walk.*

Let us now explain how these theorems follow from the more general statements given in [22]. In that paper, He considers **disc patterns** of simple proper plane triangulations, which are like circle packings except that circles corresponding to adjacent vertices intersect at a specified angle, between 0 and π , rather than being tangent. Given a proper plane network $G = (V, E)$ with locally finite dual and face set F , we can form a proper plane triangulation T with vertex set $V \cup F$ by adding a vertex inside each face of G and connecting this vertex to each vertex in the boundary of the face. It is easily verified that T is simple if and only if G is polyhedral. Moreover, a double circle packing of T , either in the plane or the disc, can now be obtained using [22, Theorem 1.3] by requiring that the angle between a pair of circles that are adjacent in T is 0 if the pair corresponds to an edge of G and $\pi/2$ otherwise. It is straightforward to check that conditions (C1) and (C2) of [22, Theorem 1.3] hold, and that the packing obtained this way is indeed the double circle packing of G . Thus, the existence statement of Theorem 2.6 follows. The uniqueness of this packing, as formulated in Theorem 2.6, is the content of [22, Theorems 1.1 and 1.2]. Finally, Theorem 2.7 is a direct consequence of [22, Theorem 1.3] together with [24, Theorems 2.6 and 8.1].

Given a double circle packing (P, P^\dagger) of a plane graph G in a domain $D \subseteq \mathbb{C}$, we write $\text{diam}_{\mathbb{C}}(A)$ for the Euclidean diameter of the set $\{z(v) : v \in A\}$, and write $d_{\mathbb{C}}(A, B)$ for the Euclidean distance between the sets $\{z(v) : v \in A\}$ and $\{z(v) : v \in B\}$. We write $r(v)$ and $r(f)$ for the Euclidean radii of the circles $P(v)$ and $P^\dagger(f)$. If (P, P^\dagger) has carrier \mathbb{D} , we write $\sigma(v)$ for $1 - |z(v)|$, which is the distance between $z(v)$ and the boundary of \mathbb{D} , and write $r_{\mathbb{H}}(v)$ for the hyperbolic radius of $P(v)$. (Recall that Euclidean circles in \mathbb{D} are also hyperbolic circles with different centres and radii.)

We will also use the Ring lemma of Rodin and Sullivan [45]; see [21] and [1] for quantitative versions. In Section 4.1 we formulate and prove a version of the Ring lemma for double circle packings, Theorem 4.1.

THEOREM 2.8 (The Ring Lemma). *There exists a sequence of positive constants $\langle k_m : m \geq 3 \rangle$ such that for every circle packing P of every simple triangulation T in a domain $D \subseteq \mathbb{C} \cup \{\infty\}$, and every pair of adjacent vertices u and v of T such that $P(v)$ does not contain ∞ , the ratio of Euclidean radii $r(v)/r(u)$ is at most $k_{\deg(v)}$.*

An immediate corollary of the Ring lemma is that, whenever P is a circle packing in \mathbb{D} of a CP hyperbolic proper plane triangulation T and v is a vertex of T , the hyperbolic radius of $P(v)$ is bounded above by a constant $C = C(\deg(v))$, see the first paragraph of Section 4.3.

3. Connectivity of the FUSF

In this section we prove Theorem 1.2 and show it easily implies Theorem 1.1. We will write \leq , \geq and \asymp for inequalities that hold up to a positive multiplicative constant depending only on $\sup_{f \in F} \deg(f)$ and $\sup_{e \in E} c(e)^{-1}$.

Disclaimer: The proof of Theorem 1.2 will also essentially yield a proof of the upper bound of Theorems 1.3 and 1.4. In particular, the inequality (10) is a direct analogue of this bound, see the inequality on the right-hand side of (31). Unfortunately, since Theorems 1.1 and 1.2 require weaker assumptions on the graph G than Theorems 1.3 and 1.4 do, it will be technically convenient for the proofs of this section to use the circle packing of the triangulation obtained by adding a star inside each face of G (sometimes referred to as the *ball-bearing packing*), rather than the double circle packing of G : For Theorems 1.1 and 1.2 it is better to use the ball-bearing packing than the double circle packing due to the way the Ring lemma works, while the fact that the double circle packing defines a good embedding of G in the sense of [4] makes it better to use in Theorems 1.3 and 1.4. Nevertheless, we will want to use the bound analogous to (10) for double circle packing in Section 4, which holds for any polyhedral proper plane graph with bounded local geometry. The proof of this bound follows by a superficial modification of the proof of (10), the details of which are omitted. The reader may wish to keep this in mind as they read this section, and observe that the proofs are easily generalized to the alternative setting once the Ring lemma for double circle packings (Theorem 4.1) is established.

3.1. Preliminaries. We begin with some preliminary estimates that will be used in the proof. We first reduce the statement of Theorem 1.2 to the case where the graph is simple and has no *peninsulas*.

Let G be a bounded codegree proper plane network with locally finite dual G^\dagger . Recall that a **peninsula** of G is a finite connected component of $G \setminus \{v\}$ for some vertex $v \in V$. In order to apply the theory of circle packing, we first reduce to the case in which G is simple and does not contain a peninsula; this will ensure that the triangulation $T = T(G)$ formed by drawing a star inside each face of G is simple. We will then use the circle packing of T to analyse the WUSF of G . We first deal with the case that every vertex of G is in a peninsula.

LEMMA 3.1. *Let G be a network such that every vertex v of G is contained in a peninsula of G . Then every essential spanning forest of G is connected and one-ended. In particular, it follows that the WUSF of G is connected and one-ended a.s.*

Proof. Let v_0 be an arbitrary vertex of G , and for each $i \geq 1$ let v_i be a vertex of G such that v_{i-1} is contained in a finite connected component of $G \setminus \{v_i\}$. Let u be another vertex of G and let γ be a path from v_0 to u in G . If u is contained in an infinite connected component of $G \setminus \{v_i\}$, then γ must pass through the vertex v_i . Since γ is finite, we deduce that there exists an integer I such that u is contained in a finite connected component of $G \setminus \{v_i\}$ for all $i \geq I$. Thus, any infinite simple path from u in G must visit v_i for all $i \geq I$. We deduce that every essential spanning forest of G is connected and one-ended as claimed. \square

Thus, to prove Theorem 1.2, it suffices to consider the case that there is some vertex of G that is not contained in a peninsula. In this case, let G' be the simple plane network formed from G by first splitting every edge e of G into a path of length 2 with both edges given the weight $c(e)$, and then deleting every peninsula of the resulting network. The assumption that G has bounded codegrees and edge resistances bounded above ensures that G' does also (indeed, the maximum codegree of G' is at most twice that of G).

LEMMA 3.2. *Let G be a plane network with at least one vertex that is not contained in a peninsula, and let G' be as above. Then every component of the WUSF of G is one-ended a.s. if and only if every component of the WUSF of G' is one-ended a.s.*

Proof. Let \mathfrak{F} be a sample of WUSF_G , and let \mathfrak{F}' be the essential spanning forest of G' defined as follows. For each edge $e \in \mathfrak{F}$ such that neither endpoint of e is contained in a peninsula of G , let both of the edges of the path of length two corresponding to e in G' be included in \mathfrak{F}' . For each edge $e \notin \mathfrak{F}$ such that neither endpoint of e is contained in a peninsula of G , choose uniformly

and independently exactly one of the two edges of the path of length two corresponding to e in G' to be included in \mathfrak{F}' . Since \mathfrak{F} has no cycles, and since an infinite simple path in \mathfrak{F} cannot enter a peninsula that it does not start inside, it is easily deduced that \mathfrak{F} is one-ended if and only if \mathfrak{F}' is one-ended. Thus, to prove the claim it suffices to prove that \mathfrak{F}' is distributed according to $\text{WUSF}_{G'}$.

To prove this, first notice that if v is a vertex in a finite connected network H such that v 's removal disconnects H into components H_1, \dots, H_k , then the UST of H is the union of k independent USTs on $H_i \cup \{v\}$ for $1 \leq i \leq k$. Applying this argument on finite subgraphs of G and taking appropriate limits, it follows that the forest obtained by deleting every edge of \mathfrak{F} that has at least one of its endpoints in a peninsula of G is distributed as the WUSF of the network obtained from G by erasing all its peninsulas.

Now suppose that H is a finite network on n vertices and m edges and that H' is obtained from H by splitting every edge of H into a path of length 2 with both edges given conductance $c(e)$ so that H' has $n + m$ vertices and $2m$ edges. For each spanning tree T' of H' , we obtain a spanning tree $\phi(T')$ of H by erasing each leaf of T' that is one of the newly added vertices and replacing each path of length 2 in which the middle vertex is newly added by a single edge. Each spanning tree T of H has exactly 2^{m-n+1} preimages under the map, and every preimage $T' \in \phi^{-1}(T)$ has

$$\prod_{e' \in T'} c(e') = \left[\prod_{e \in E} c(e) \right] \left[\prod_{e \in T} c(e) \right].$$

Indeed, if $e \in T$ then both of the corresponding edges in H' must be included in any tree in $\phi^{-1}(T)$, while if $e \notin T$ then each tree in $\phi^{-1}(T)$ must include exactly one of the corresponding edges of H' , and any combination is possible. Thus, we deduce that if \mathfrak{T}' is distributed as a uniform spanning tree of H' then $\phi(\mathfrak{T}')$ is distributed as a uniform spanning tree of H . Applying this argument on finite subgraphs of G and taking appropriate limits, we deduce by this and the previous paragraph that the law of \mathfrak{F}' is $\text{WUSF}_{G'}$ as claimed. \square

Thus it suffices to consider the case that G is simple and has no peninsulas. This assumption allows us to circle pack G as follows. Let T be the triangulation obtained by adding a vertex inside each face of G and drawing an edge between this vertex and each vertex of the face it corresponds to. The assumption that G is simple and does not contain a peninsula ensures that T is a simple triangulation. We identify the vertices of T with $V(G) \cup F(G)$, where $F(G)$ is the set of faces of G . Let $P = \{P(v) : v \in V(G)\} \cup \{P(f) : f \in F(G)\}$ be a circle packing of T in either the plane or the unit disc. (Note that this is not the double circle packing

of G , which is less useful for us at this stage since we are not assuming that G has bounded degrees.)

We proceed with two geometric lemmas. For each $r' > r > 0$ and z in the carrier of P let $V_z(r, r')$ be the set of vertices v of G such that either the intersection of the circle $P(v)$ with the annulus $A_z(r, r')$ is nonempty or the circle $P(v)$ is contained in the ball $B_z(r)$ and there is a face f incident to v such that the intersection $P(f) \cap A_z(r, r')$ is nonempty. We define the set $W_z(r) := V_z(r, r)$ similarly by using a circle instead of an annulus.

LEMMA 3.3 (Existence of a uniformly large number of disjoint annuli around a point close to the boundary). *Suppose that T is CP hyperbolic so that the carrier of P is \mathbb{D} . Then the following hold:*

- (1) *There exists a decreasing sequence $\langle r_n \rangle_{n \geq 0}$ with $r_n \in (0, 1/4)$ such that for every $z \in \mathbb{D}$ with $|z| \geq 1 - r_n$, the sets*

$$V_z(r_i, 2r_i) \quad 1 \leq i \leq n$$

are disjoint.

- (2) *If G has bounded degrees, then there exists a constant $C = C(\mathbf{M}_G)$ such that we may take $r_n = C^{-n}$ in the previous statement.*

Proof. We first prove item (1). We construct the sequence recursively, letting $r_0 = 1/8$. Suppose that $\langle r_i \rangle_{i=0}^n$ satisfying the conclusion of the lemma have already been chosen, and consider the set of vertices

$$K_n = \left\{ v \in V(G) : r(v) \geq \frac{r_n}{4} \text{ or } r(f) \geq \frac{r_n}{4} \text{ for some face } f \text{ incident to } v \right\}.$$

Since the carrier of P has finite area K_n is a finite set. We define r_{n+1} to be

$$r_{n+1} = \sup\{r \leq r_n/4 : P(v) \subseteq \overline{B_0(1 - 3r)} \text{ for all } v \in K_n\}$$

which is positive since K_n is finite. We claim that $\langle r_i \rangle_{i=0}^{n+1}$ continues to satisfy the conclusion of the lemma. That is, we claim that $V_z(r_{n+1}, 2r_{n+1}) \cap V_z(r_i, 2r_i) = \emptyset$ for every $1 \leq i \leq n$ and every $z \in \mathbb{D}$ with $|z| \geq 1 - r_{n+1}$.

Indeed, let z be such that $|z| \geq 1 - r_{n+1}$ and let $v \in V_z(r_{n+1}, 2r_{n+1})$. By definition of $V_z(r_{n+1}, 2r_{n+1})$, the intersection $P(v) \cap B_z(2r_{n+1})$ is nonempty, and we deduce that $P(v)$ is not contained in the closure of $B_0(1 - 3r_{n+1})$. By definition of r_{n+1} , it follows that $v \notin K_n$ and hence that $r(v) \leq r_n/4$ and $r(f) \leq r_n/4$ for every face f incident to v . Thus, since $r_{n+1} \leq r_n/4$,

$$P(v) \subseteq B_z(2r_{n+1} + 2r_n/4) \subseteq B_z(r_n) \quad \text{and} \quad P(f) \subseteq B_z(2r_{n+1} + 2r_n/4) \subseteq B_z(r_n)$$

for every face f incident to v . It follows that $V_z(r_{n+1}, 2r_{n+1}) \cap V_z(r_i, 2r_i) = \emptyset$ for every $1 \leq i \leq n$ as claimed.

To prove (2), observe that, by the Ring lemma (Theorem 2.8), there exists a constant $k = k(\mathbf{M})$ such that for every $C > 1$, every z with $|z| \geq 1 - C^{-m}$, and every $0 \leq n \leq m$, every circle in P that either has centre in $A_z(C^{-n}, 2C^{-n})$ or is tangent to some circle with centre in $A_z(C^{-n}, 2C^{-n})$ has radius at most kC^{-n} . Thus, this set of circles is contained in the ball $B_z((2 + 4k)C^{-n})$. It follows that taking $C = 1/(4 + 8k)$ suffices. \square

Finally, we estimate the energy of a random set of vertices that we will frequently use.

LEMMA 3.4. *Let z be a point in the carrier of P (which may be either \mathbb{C} or \mathbb{D}), let U be a uniform random variable on the interval $[1, 2]$ and, for each $r > 0$, let μ_r be the law of the random set of vertices $W_z(Ur) = V_z(Ur, Ur)$. Then*

$$\mathcal{E}(\mu_r) \leq 1$$

uniformly in $r > 0$, where $\mathcal{E}(\mu_r)$ was defined in (3).

Proof. For a vertex v of G to be included in $W_z(Ur)$, the circle

$$\{z' \in \mathbb{C} : |z' - z| = Ur\}$$

must either intersect the circle $P(v)$ or intersect $P(f)$ for some face f incident to v . The union of $P(v)$ and all the $P(f)$ incident to $P(v)$ is contained in the ball of radius $r(v) + 2 \max_{f \sim v} r(f)$ around $z(v)$. Since the codegrees of G are bounded, the Ring lemma implies that $r(f) \leq r(v)$ for all incident $v \in V$ and $f \in F$, and so

$$\mu_r(v \in W_z(Ur)) \leq \frac{1}{r} \min\left(2r(v) + 4 \max_{f \sim v} r(f), r\right) \leq \frac{1}{r} \min\{r(v), r\}. \quad (5)$$

We claim that

$$\sum_{v \in V_z(r, 2r)} \min\{r(v), r\}^2 \leq 16r^2. \quad (6)$$

To see this, replace each circle of a vertex in $V_z(r, 2r)$ that has radius larger than r with a circle of radius r that is contained in the original circle and intersects $B_z(2r)$: The circles in this new set still have disjoint interiors, are contained in the ball $B_z(4r)$, and have total area $\pi \sum_{v \in V_z(r, 2r)} \min\{r(v), r\}^2$, yielding (6). The claim follows from (5) and (6) by definition of the energy $\mathcal{E}(\mu_r)$. \square

We are now ready to prove Theorem 1.2.

3.2. Proof of Theorem 1.2. Let G be a simple proper plane network with bounded codegrees and bounded edge resistances. By Lemma 3.2 we can assume that G does not contain a peninsula. Let \mathfrak{F} be a sample of WUSF_G and given an edge $e = (x, y)$ let \mathcal{A}^e be the event that x and y are in distinct infinite connected components of $\mathfrak{F} \setminus \{e\}$. It is clear that every component of \mathfrak{F} is one-ended a.s. if and only if

$$\text{WUSF}_G(e \in \mathfrak{F}, \mathcal{A}^e) = 0 \tag{7}$$

for every edge e of G .

Consider the triangulation T obtained from G and its circle packing $P = \{P(v) : v \in V(G)\} \cup \{P(f) : f \in F(G)\}$, as described in the previous subsection. By applying a Möbius transformation, we normalize P by setting the centres $z(x)$ and $z(y)$ to be on the negative and positive real axes, respectively, setting the circles $P(x)$ and $P(y)$ to have the origin as their tangency point and, in the parabolic case, fixing the scale by setting $z(y) - z(x) = 1$.

Let $\varepsilon > 0$ be arbitrarily small. If T is CP hyperbolic, let V_ε be the set of vertices of G with $|z(v)| \leq 1 - \varepsilon$. Otherwise, T is CP parabolic and we define V_ε to be the set of vertices of G with $|z(v)| \leq \varepsilon^{-1}$. We also denote by E_ε the set of edges that have both endpoints in V_ε .

Let $\mathcal{B}_\varepsilon^e$ be the event that every component of $\mathfrak{F} \setminus \{e\}$ intersects $V \setminus V_\varepsilon$. On the event $\mathcal{B}_\varepsilon^e$, we define η^x to be the rightmost path in $\mathfrak{F} \setminus \{e\}$ from x to $V \setminus V_\varepsilon$ when looking at x from y , and η^y to be the leftmost path in $\mathfrak{F} \setminus \{e\}$ from y to $V \setminus V_\varepsilon$ when looking at y from x . Note that the paths η_x and η_y are not necessarily disjoint. Nonetheless, concatenating the reversal of η^x with e and η^y separates $E_\varepsilon \setminus (\eta_x \cup \eta_y)$ into two sets of edges, \mathcal{L} and \mathcal{R} , which are to the left and right of e (when viewed from x to y), respectively. See Figure 3 for an illustration of the case when η_x and η_y are disjoint (when they are not, \mathcal{R} is a ‘bubble’ separated from $V \setminus V_\varepsilon$).

On the event $\mathcal{B}_\varepsilon^e$, let K be the set of edges that are either in \mathcal{L} or belong to $\eta_x \cup \eta_y$, setting $K = E$ off of this event. The condition that η^x and η^y are the rightmost and leftmost paths to $V \setminus V_\varepsilon$ from x and y is equivalent to the condition that K does not contain any open path from x to $V \setminus V_\varepsilon$ other than η^x , and does not contain any open path from y to $V \setminus V_\varepsilon$ other than η^y . It follows from this characterization that K is a local set for \mathfrak{F} . (Indeed, we note that K can be explored algorithmically, without querying the status of any edge in $E \setminus K$, by performing a right-directed depth-first search of x ’s component in \mathfrak{F} and a left-directed depth-first search of y ’s component in \mathfrak{F} , stopping each search when it first leaves V_ε .)

Let $\mathcal{A}_\varepsilon^e$ denote the event that $\mathcal{B}_\varepsilon^e$ occurs and that K does not contain an open path from x to y , or, equivalently, that η_x and η_y are disjoint (since η_x and η_y were chosen to be rightmost and leftmost, respectively). Note that $\mathcal{A}_\varepsilon^e$ is measurable

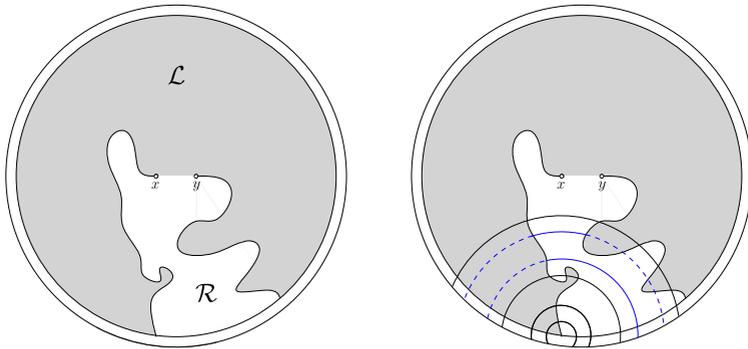


Figure 3. Illustration of the proof of Theorem 1.2 in the case that T is CP hyperbolic. Left: On the event $\mathcal{A}_\varepsilon^e$, the paths η^x and η^y split $E_\varepsilon \setminus (\eta_x \cup \eta_y)$ into two pieces, \mathcal{L} and \mathcal{R} . Right: We define a random set containing a path (solid blue) from η^x to $\eta^y \cup \{\infty\}$ in $G \setminus K_c$ using a random circle (dashed blue). Here we see two examples, one in which the path ends at η^y , and the other in which the path ends at the boundary (that is, at infinity).

with respect to \mathcal{F}_K (as defined in Section 2.3.1), and that $\mathcal{A}^e = \bigcap_{\varepsilon > 0} \mathcal{A}_\varepsilon^e$. Thus,

$$\text{WUSF}_G(e \in \mathfrak{F}, \mathcal{A}^e) \leq \text{WUSF}_G(e \in \mathfrak{F} \mid \mathcal{A}_\varepsilon^e) = \mathbb{E}[\text{WUSF}_G(e \in \mathfrak{F} \mid \mathcal{F}_K, \mathcal{A}_\varepsilon^e)].$$

By the strong spatial Markov property (Proposition 2.1), conditioned on \mathcal{F}_K and the event $\mathcal{A}_\varepsilon^e$, the law of \mathfrak{F} is equal to the union of K_o with a sample of the WUSF of the network $(G - K_c)/K_o$ obtained from G by deleting all the revealed closed edges and contracting all the revealed open edges. In particular, by Kirchhoff's Effective Resistance Formula (Theorem 2.3),

$$\text{WUSF}_G(e \in \mathfrak{F} \mid \mathcal{F}_K, \mathcal{A}_\varepsilon^e) \leq c(e) \mathcal{R}_{\text{eff}}^{\text{W}}(\eta^x \leftrightarrow \eta^y; G - K_c). \quad (8)$$

Since the edge e was arbitrary, to prove (7) (and hence Theorem 1.2) it suffices to prove that there is an upper bound on the effective resistance appearing in (8) that tends to zero as $\varepsilon \rightarrow 0$ uniformly in \mathcal{F}_K . We perform this analysis now according to whether T is CP hyperbolic or parabolic.

Proof of Theorem 1.2, hyperbolic case. Suppose that T is CP hyperbolic, let v^x be the endpoint of the path η^x and let $z_0 = z(v^x)$. On the event $\mathcal{A}_\varepsilon^e$, for each $1 - |z_0| \leq r \leq 1/4$, we claim that the set $W_{z_0}(r)$, as defined just before Lemma 3.3, contains a path in G from η^x to $\eta^y \cup \{\infty\}$ that is contained in $\mathcal{R} \cup \eta^x \cup \eta^y$, and is therefore a path in $G \setminus K_c$.

Indeed, consider the arc $\mathcal{A}'(r) = \{z \in \overline{\mathbb{D}} : |z - z_0| = r\}$, parameterized in the clockwise direction. Let $\mathcal{A}(r)$ be the subarc of $\mathcal{A}'(r)$ beginning at the last time that $\mathcal{A}'(r)$ intersects a circle corresponding to a vertex in the trace of η^x , and ending at the first time after this time that $\mathcal{A}'(r)$ intersects either $\partial\mathbb{D}$ or a circle corresponding to a vertex in the trace of η^y (see Figure 3). Thus, on the event $\mathcal{A}_\varepsilon^e$, the set of vertices of T whose corresponding circles in P are intersected by $\mathcal{A}(r)$ contains a path in T from η^x to $\eta^y \cup \{\infty\}$, for every $1 - |z_0| \leq r \leq 1/4$. (Indeed, for Lebesgue a.e. $1 - |z_0| \leq r \leq 1/4$, the arc $\mathcal{A}(r)$ is not tangent to any circle in P , and in this case the set is precisely the trace of a simple path in T .) To obtain a path in G rather than T , we divert the path counterclockwise around each face of G . That is, whenever the path passes from a vertex u of G to a face f of G and then to a vertex v of G , we replace this section of the path with the list of vertices of G incident to f that are between u and v in the counterclockwise order. This construction shows that the subgraph of $G \setminus K_c$ induced by the set $W_{z_0}(r)$ contains a path from η^x to $\eta^y \cup \{\infty\}$, as claimed.

Let $\langle r_n \rangle_{n \geq 0}$ be as in Lemma 3.3 and let $n(\varepsilon)$ be the maximum n such that $\varepsilon < r_n$. By Lemma 3.3 the measures μ_{r_i} are supported on sets that are contained in the disjoint sets $V_z(r_i, 2r_i)$. Thus, by Lemma 2.4, our assumption that edge resistances are bounded above, and Lemma 3.4, we have

$$\mathcal{R}_{\text{eff}}^W(\eta^x \leftrightarrow \eta^y \cup \{\infty\}; G \setminus K_c) \leq \mathcal{E}\left(\frac{1}{n(\varepsilon)} \sum_{i=1}^{n(\varepsilon)} \mu_{r_i}\right) = \frac{1}{n(\varepsilon)^2} \sum_{i=1}^{n(\varepsilon)} \mathcal{E}(\mu_{r_i}) \leq \frac{1}{n(\varepsilon)}$$

and hence, by symmetry,

$$\mathcal{R}_{\text{eff}}^W(\eta^y \leftrightarrow \eta^x \cup \{\infty\}; G \setminus K_c) \leq \frac{1}{n(\varepsilon)}.$$

Applying Lemma 2.2 and (8), we have

$$\text{WUSF}(e \in \mathfrak{F} \mid \mathcal{F}_K, \mathcal{A}_\varepsilon^e) \leq \frac{c(e)}{n(\varepsilon)}, \tag{9}$$

which by Lemma 3.3 converges to zero as $\varepsilon \rightarrow 0$, completing the proof of Theorem 1.2 in the case that T is CP hyperbolic. If G has bounded degrees, then combining (9) with Lemma 3.3 implies that there exists a positive constant $C = C(\mathbf{M})$ such that

$$\text{WUSF}(e \in \mathfrak{F} \mid \mathcal{F}_K, \mathcal{A}_\varepsilon^e) \leq C \frac{c(e)}{\log(1/\varepsilon)} \tag{10}$$

for all $\varepsilon \leq 1/2$. This inequality is a direct analogue of the upper bound of Theorem 1.4, see the inequality on the right of (31). \square

Proof of Theorem 1.2, parabolic case. Suppose that T is CP parabolic. Let $r_1 = 1$ and define $\langle r_n \rangle_{n \geq 1}$ recursively by

$$r_n = 1 + \inf\{r \geq r_{n-1} : V_0(r, 2r) \cap V_0(r_{n-1}, 2r_{n-1}) = \emptyset\}.$$

Since $V_0(r_{n-1}, 2r_{n-1})$ is finite, this infimum is finite. By definition, the sets $V_0(r_i, 2r_i)$ are disjoint. A similar analysis to the hyperbolic case shows that, on the event $\mathcal{A}_\varepsilon^e$, for each $r \geq 1$, the set $W_0(r)$ contains a path in G from η^x to η^y that is contained in $\mathcal{R} \cup \eta^x \cup \eta^y$, and is therefore a path in $G \setminus K_c$. For each $\varepsilon > 0$, let $n(\varepsilon)$ be the maximal n such that $r_n \leq \varepsilon^{-1}$. Then on the event $\mathcal{A}_\varepsilon^e$, by Lemma 2.4 and Lemma 3.4,

$$\mathcal{R}_{\text{eff}}^F(\eta^x \leftrightarrow \eta^y; G \setminus K_c) \leq \mathcal{E}\left(\frac{1}{n(\varepsilon)} \sum_{i=1}^{n(\varepsilon)} \mu_{r_i}\right) \leq \frac{1}{n(\varepsilon)^2} \sum_{i=1}^{n(\varepsilon)} \mathcal{E}(\mu_{r_i}) \leq \frac{1}{n(\varepsilon)}.$$

Thus, by (8),

$$\text{WUSF}_G(e \in \mathfrak{F} \mid \mathcal{F}_K, \mathcal{B}_\varepsilon^e) \leq \frac{c(e)}{n(\varepsilon)}. \quad (11)$$

The right-hand side converges to zero as $\varepsilon \rightarrow 0$, completing the proof of Theorem 1.2. \square

REMARK 3.5. Since the random sets used in the CP parabolic case above are always finite, they can also be used to bound free effective resistances. Therefore, by repeating the proof above with the FUSF in place of the WUSF, we deduce that every component of the FUSF of G is one-ended a.s. if T is CP parabolic. Since the FUSF stochastically dominates the WUSF, and an essential spanning forest with one-ended components does not contain any strict subgraphs that are also essential spanning forests, we deduce that if T is CP parabolic, then the FUSF and WUSF of G coincide. In particular, using [10, Theorem 7.3], we obtain the following result: If T is a CP parabolic proper plane triangulation then T does not admit nonconstant harmonic functions of finite Dirichlet energy.

This result can alternatively be obtained by noticing that the usual proof of the He–Schramm theorem establishes that every bounded degree CP parabolic polyhedral proper plane map is recurrent (that is, bounded codegree is not needed). Thus, if T is a CP parabolic proper plane triangulation then the *dual* of T is recurrent, so that in particular its free and wired spanning forests coincide, and the claim in the previous paragraph follows by duality.

3.3. The FUSF is connected.

Proof of Theorem 1.1. First suppose that G^\dagger is locally finite. Since G has bounded degrees and bounded conductances, G^\dagger has bounded codegrees and bounded resistances. Thus, Theorem 1.2 implies that every component of the WUSF on G^\dagger is one-ended a.s. and consequently, by Theorem 2.5, that the FUSF of G is connected a.s.

Now suppose that G does not have locally finite dual. In this case, we form a plane network (G', c') from G by adding edges to triangulate the infinite faces of G while keeping the degrees bounded. We enumerate these additional edges $\{e_i\}_{i \geq 1}$ and define conductances

$$c'(e_i) = 2^{-i-1} \mathcal{R}_{\text{eff}}^F(e_i^- \leftrightarrow e_i^+; G)^{-1}.$$

Since G' has bounded degrees, bounded conductances and locally finite dual, its FUSF is connected a.s. By Kirchhoff’s Effective Resistance Formula (Theorem 2.3), Rayleigh monotonicity and the union bound, the probability that the FUSF of G' contains any of the additional edges e_i is at most

$$\sum_{i \geq 0} c'(e_i) \mathcal{R}_{\text{eff}}^F(e_i^- \leftrightarrow e_i^+; G') \leq \sum_{i \geq 0} c'(e_i) \mathcal{R}_{\text{eff}}^F(e_i^- \leftrightarrow e_i^+; G) \leq 1/2.$$

In particular, there is a positive probability that none of the additional edges are contained in the FUSF of G' . The conditional distribution of the FUSF of G' on this event is FUSF_G by the spatial Markov property, and it follows that the FUSF of G is connected a.s. □

4. Critical exponents

4.1. The Ring lemma for double circle packings. In this section we extend the Ring lemma to double circle packings. While unsurprising, we were unable to find such an extension in the literature.

THEOREM 4.1 (Ring Lemma). *There exists a family of positive constants $\{k_{n,m} : n \geq 3, m \geq 3\}$ such that if (P, P^\dagger) is a double circle packing in $\mathbb{C} \cup \{\infty\}$ of a polyhedral plane graph G and v is a vertex of G , then for every $f \in F$ incident to v such that $P(v)$ does not contain ∞ , then*

$$r(v)/r(f) \leq k_{\deg(v), \max_{g \perp v} \deg(g)}$$

where $g \perp v$ means that the face g is incident to the vertex v .

As for triangulations, the Ring lemma immediately implies that whenever a polyhedral, CP hyperbolic proper plane network G with bounded degrees and codegrees is circle packed in \mathbb{D} , the hyperbolic radii of the discs in $P \cup P^\dagger$ are bounded above by a constant C depending only on the maximal degree and codegree.

Proof of Theorem 4.1. Let n be the degree of v and let m be the maximum degree of the faces incident to v in G . We may assume that $r(v) = 1$. Note that for each two distinct discs $P^\dagger(f), P^\dagger(f') \in P^\dagger$ that are not tangent, there is at most one disc $P(u) \in P$ that intersects both $P^\dagger(f)$ and $P^\dagger(f')$, while if $P^\dagger(f)$ and $P^\dagger(f')$ are tangent, there exist exactly two discs in P that intersect both $P^\dagger(f)$ and $P^\dagger(f')$. For each two faces f and f' incident to v , the complement $\partial P(v) \setminus (P^\dagger(f_1) \cup P^\dagger(f_2))$ is either a single arc (if $P^\dagger(f)$ and $P^\dagger(f')$ are tangent), or is equal to the union of two arcs (if $P^\dagger(f)$ and $P^\dagger(f')$ are not tangent).

We claim that there exists a function $\psi_m(\cdot, \cdot) : (0, \infty)^2 \rightarrow (0, 2\pi]$, increasing in both coordinates, such that if f and f' are two distinct faces of G incident to v , then each of the (one or two) arcs forming the complement $\partial P(v) \setminus (P^\dagger(f_1) \cup P^\dagger(f_2))$ have length at least $\psi_m(r(f), r(f'))$. Indeed, let $r(f)$ and $r(f')$ be fixed, and suppose that one of the arcs forming the complement $\partial P(v) \setminus (P^\dagger(f) \cup P^\dagger(f'))$ is extremely small, with length ε . Let the primal circles incident to f be enumerated $v_1, \dots, v_{\deg(f)}$, where $v_1 = v$, $P(v_2)$ is the primal circle that is tangent to $P(v)$ and intersects $P^\dagger(f)$ on the same side as the small arc, $P(v_3)$ is the next primal circle that is tangent to $P(v_2)$ and intersects $P^\dagger(f)$, and so on. Since ε is small, $P(v_2)$ must also be very small, as it does not intersect $P^\dagger(f')$. Similarly, if ε is sufficiently small, $P(v_3)$ must also be small, since it also does not intersect $P^\dagger(f')$. See Figure 4. Applying this argument recursively, we see that, if ε is sufficiently small, then the circles $P(v_2), \dots, P(v_{\deg(f)})$ are collectively too small to contain $\partial P^\dagger(f) \setminus P(v)$ in their union, a contradiction. We write $\psi_m(r(f), r(f'))$ for the minimal ε that is not ruled out as impossible by this argument.

Let the faces incident to the vertex v be indexed in clockwise order f_1, \dots, f_n , where $n = \deg(v)$ and f_1 has maximal radius among the faces incident to v . For each face f incident to v , the arc $\partial P(v) \cap P^\dagger(f)$ has length $2 \tan^{-1}(r(f))$. Since $\partial P(v) = \bigcup_{i=1}^n (\partial P(v) \cap P^\dagger(f_i))$, we deduce that $r(f_1)$ is bounded below by $\tan(\pi/n)$. By definition of ψ_m , we have that $r(f_k)$ satisfies

$$\psi_m(\tan(\pi/n), r(f_k)) \leq \psi_m(r(f_1), r(f_k)) \leq \sum_{i=2}^{k-1} 2 \tan^{-1}(r(f_i)) \quad (12)$$

for all $3 \leq k \leq n$. For each such k , (12) yields an implicit upper bound on $r(f_k)$ which converges to zero as $r(f_2)$ converges to zero. This in turn yields a uniform lower bound on $r(f_2)$: if $r(f_2)$ were sufficiently small, the bound (12) would

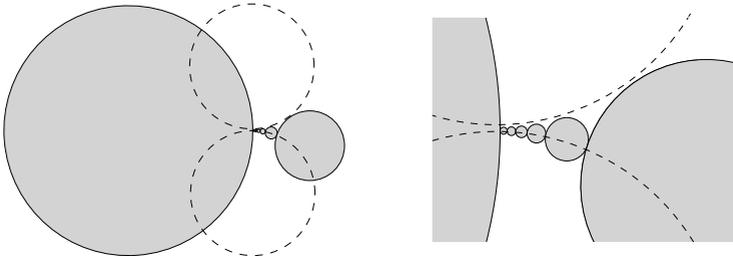


Figure 4. Proof of the Double Ring lemma. If two dual circles are close but do not touch, there must be many primal circles contained in the crevasse between them. This forces the two dual circles to each have large degree. The right-hand figure is a magnification of the left-hand figure.

imply that $\sum_{k=2}^n 2 \tan^{-1}(r(f_k))$ would be less than π , a contradiction (since we have trivially that $2 \tan^{-1}(r(f_1)) \leq \pi$). We obtain uniform lower bounds on $r(f_k)$ for each $3 \leq k \leq n$ by repeating the above argument inductively. \square

4.2. Good embeddings of planar graphs. If G is a plane graph and (P, P^\dagger) is a double circle packing of G in a domain $D \subseteq \mathbb{C}$, then drawing straight lines between the centres of the circles in P yields a proper embedding of G in D in which every edge is a straight line. We call such an embedding a **proper straight-line embedding of G** in D . Following [4], a proper straight-line embedding of a graph G in a domain $D \subseteq \mathbb{C}$ is said to be **η -good** if the following conditions are satisfied:

- (1) **(No near-flat, flat, or reflexive angles.)** All internal angles of every face in the drawing are at most $\pi - \eta$. In particular, every face is convex.
- (2) **(Adjacent edges have comparable lengths.)** For every pair of edges e, e' sharing a common endpoint, the ratio of the lengths of the straight lines corresponding to e and e' in the drawing is at most η^{-1} .

Let G be a plane graph with bounded degrees and codegrees and let $e = (u, v)$ and $e' = (v, w)$ be two edges of a face f meeting at a vertex v and consider its double circle packing (P, P^\dagger) . By the Ring lemma (Theorem 4.1) $P(u), P(v)$ and $P(w)$ all have comparable sizes. On the other hand, if the angle between the straight lines of e and e' was very close to π , this would imply that $P(f)$ is much larger than $P(v)$, contradicting the previous assertion. Thus we have the following corollary.

COROLLARY 4.2. *The proper straight-line embedding given by any double circle packing of a plane graph G with bounded degrees and bounded codegrees in a domain $D \subseteq \mathbb{C}$ is η -good for some positive $\eta = \eta(\mathbf{M}_G)$.*

We remark that, in contrast, the embedding of G obtained by circle packing the triangulation $T(G)$ formed by drawing a star inside each face of G (and then erasing these added vertices), as done in Section 3, does not necessarily yield a good embedding of G .

For the remainder of this section and in Sections 4.3, 4.5 and 4.6, G will be a fixed transient, polyhedral proper plane network with bounded codegrees and bounded local geometry, and (P, P^\dagger) will be a double circle packing of G in \mathbb{D} . We will write \leq , \geq and \asymp to denote inequalities or equalities that hold up to positive multiplicative constants depending only upon \mathbf{M} . We will also fix an edge $e = (x, y)$ of G and, by applying a Möbius transformation if necessary, normalize (P, P^\dagger) by setting the centres $z(x)$ and $z(y)$ to be on the negative real axis and positive real axis, respectively, and setting the circles $P(x)$ and $P(y)$ to have the origin as their tangency point.

In [4] and [16], several estimates are established that allow one to compare the random walk on a good embedding of a planar graph with Brownian motion. The following estimates, proven for general good embeddings of proper plane graphs in [4, Theorems 1.4 and 1.5], are of central importance to the proofs of the lower bounds in Theorems 1.4 and 1.5. Recall that $\sigma(v)$ is defined to be $1 - |z(v)|$.

THEOREM 4.3 (Diffusive time estimate [4, Theorem 1.5]). *There exists a constant $C_1 = C_1(\mathbf{M}) \geq 1$ such that the following holds. For each vertex v of G , let $\langle X_n \rangle_{n \geq 0}$ be a random walk on G started at v , let $r(v) \leq r \leq C_1^{-1} \sigma(v)$ and let T_r be the first time n that $|z(X_n) - z(v)| \geq r$. Then*

$$\mathbf{E}_v \sum_{n=0}^{T_r} r(X_n)^2 \asymp r^2.$$

Murugan [42] has shown that the constant C_1 above can in fact be taken to be 1.

THEOREM 4.4 (Cone estimate [4, Theorem 1.4]). *Let $\eta = \eta(\mathbf{M}) > 0$ be the constant from Corollary 4.2. There exists a positive constant $q_1 = q_1(\mathbf{M})$ such that the following holds. For each vertex v of G , let $\langle X_n \rangle_{n \geq 0}$ be a random walk on G started at v , let $0 \leq r \leq \sigma(v)$, and let T_r be the first time n that $|z(X_n) - z(v)| \geq r$. Then for any interval $I \subset \mathbb{R}/(2\pi\mathbb{Z})$ of length at least $\pi - \eta$ we have*

$$\mathbf{P}_v(\arg(z(X_{T_r}) - z(v)) \in I) \geq q_1.$$

We also apply the following result of Benjamini and Schramm [11, Lemma 5.3]; their proof was given for simple triangulations but extends immediately to our setting.

THEOREM 4.5 (Convergence estimate [11]). *For every vertex v of G , we have that*

$$\mathbf{P}_v(|z(X_n) - z(v)| \geq t\sigma(v) \text{ for some } n \geq 0) \leq \frac{1}{\log t}.$$

We remark that the logarithmic decay in Theorem 4.5 is not sharp. Sharp polynomial estimates of the same quantity have been obtained by Chelkak [16, Corollary 7.9].

4.3. Preliminary estimates. For each vertex v of G , let $a_{\mathbb{H}}(v)$ denote the hyperbolic area of $P(v)$ (recall that by the uniqueness of Theorem 2.6 this quantity does not depend on the choice of packing). Consider a Möbius transformation of the unit disc taking $z(v)$ to the origin. Then by the Ring lemma (Theorem 4.1), the Euclidean radius of the image of $P(v)$ under this Möbius transformation must be bounded by a constant $c = c(\mathbf{M}) < 1$ and we deduce that the hyperbolic radii of the discs in P are bounded from above. Thus,

$$a_{\mathbb{H}}(v) \asymp r_{\mathbb{H}}(v)^2 \asymp \sigma(v)^{-2}r(v)^2 \tag{13}$$

for every vertex v of G .

Recall that given a set of vertices B we write τ_B for the first time the random walk visits B , letting $\tau_B = \infty$ if no such time exists. In the following section, we will wish to estimate sums of the form

$$\sum_{u \in A} a_{\mathbb{H}}(u) \mathbf{P}_u(\tau_B < \infty) \tag{14}$$

where A and B are subsets of V . We begin with a useful preliminary estimate. For $r \in (0, 1]$ and a vertex v let $\mathfrak{H}(v, r)$ be the half-plane containing $z(v)$ whose boundary is the unique straight line with distance $r\sigma(v)$ to $z(v)$ that is orthogonal to the line connecting $z(v)$ to the origin.

LEMMA 4.6 (Remain in half-plane estimate). *For any $r \in (0, 1]$ there exists a positive constant $q_2 = q_2(\mathbf{M}, r)$, increasing in r , such that for any vertex v the probability of the random walk starting at v to remain in $\mathfrak{H}(v, r)$ forever is at least q_2 . That is,*

$$\mathbf{P}_v(z(X_n) \in \mathfrak{H}(v, r) \text{ for all } n) \geq q_2. \tag{15}$$

Proof. This estimate is similar to those used in the proof of [4, Theorem 1.4]. We will use Theorem 4.4 to ‘push’ the walker deep inside the half-plane $\mathfrak{H}(v, r)$ (the number of times we apply Theorem 4.4 will only depend on r) and once it is close enough to the boundary Theorem 4.5 will ensure that it stays close forever with positive probability. We now make this precise. For each vertex u , let $\mathfrak{C}_v(u)$ be the cone

$$\mathfrak{C}_v(u) = \{z \in \mathbb{C} : |\arg(z - z(u)) - \arg z(v)| \leq \pi/2 - \eta/2\}.$$

Define numbers $\langle \rho_i \rangle_{i \geq 0}$ and stopping times $\langle T_i \rangle_{i \geq 0}$ by putting $T_0 = 0$ and $\rho_0 = r\sigma(v)$ and recursively setting for $i \geq 1$

$$T_i = \min\{n \geq 0 : |z(X_n) - z(X_{T_{i-1}})| \geq \rho_{i-1}/2\},$$

and setting ρ_i to be the distance between $z(X_{T_i})$ and $\partial\mathfrak{H}(v, r) \cup \partial\mathbb{D}$. Since the closed ball of radius $\rho_{i-1}/2$ around $z(X_{T_{i-1}})$ is contained in \mathbb{D} , it has only finitely many vertices in it and therefore $T_i < \infty$ a.s. for all $i \geq 1$.

Denote by \mathcal{A}_i the event that $z(X_{T_i}) \in \mathfrak{C}_v(X_{T_{i-1}})$. That is, \mathcal{A}_i is the event that the random walk starting from $X_{T_{i-1}}$ exits the ball of radius $\rho_{i-1}/2$ around $z(X_{T_{i-1}})$ inside the cone that is parallel to $\mathfrak{C}_v(0)$ and is centred at $z(X_{T_{i-1}})$. By basic trigonometry this event implies that

$$\sigma(X_{T_i}) \leq \sigma(X_{T_{i-1}}) - \rho_{i-1} \sin(\eta/2)/2, \quad (16)$$

and

$$d_{\mathbb{C}}(X_{T_i}, \partial\mathfrak{H}(v, r)) \geq d_{\mathbb{C}}(X_{T_{i-1}}, \partial\mathfrak{H}(v, r)) + \rho_{i-1} \sin(\eta/2)/2. \quad (17)$$

In particular, on the event $\bigcap_{i=0}^n \mathcal{A}_i$ the distance between $z(X_{T_i})$ and $\partial\mathbb{D}$ is decreasing in i for $0 \leq i \leq n$ and the distance between $z(X_{T_i})$ and $\partial\mathfrak{H}(v, r)$ is increasing in i for $0 \leq i \leq n$.

By Theorem 4.5 there exists $t_0 = t_0(\mathbf{M}) \geq 1$ such that we have

$$\mathbf{P}_u(|z(X_n) - z(u)| \geq t_0\sigma(u) \text{ for some } n \geq 0) \leq \frac{1}{2},$$

for every vertex u . Set $\delta = \min\{r(2t_0)^{-1}, \sin(\eta/2)\}$ and put $n_0 = 2\lceil 1/\delta^2 \rceil$. We first note that by (17) the event $\bigcap_{i=0}^{n_0} \mathcal{A}_i$ implies that $z(X_n) \in \mathfrak{H}(v, r)$ for $0 \leq n \leq T_{n_0}$. Next we have that if X_{T_i} has distance at least $\delta\sigma(v)$ from $\partial\mathfrak{H}(v, r) \cup \partial\mathbb{D}$, then $\rho_i \sin(\eta/2) \geq \delta^2\sigma(v)$. Therefore, if the event $\bigcap_{i=0}^{n_0} \mathcal{A}_i$ occurs, then $X_{T_{n_0}}$ has distance at most $\delta\sigma(v)$ from $\partial\mathbb{D}$ (otherwise the contradictory assertion that $\sigma(X_{T_{n_0}}) \leq 0$ follows by (16)) and has distance at least $r\sigma(v)$ from $\partial\mathfrak{H}(v, r)$ by (17). That is,

$$\sigma(X_{T_{n_0}}) \leq r(2t_0)^{-1}\sigma(v) \quad \text{and} \quad d_{\mathbb{C}}(X_{T_{n_0}}, \partial\mathfrak{H}(v, r)) \geq r\sigma(v).$$

Therefore, by the definition of t_0 and the strong Markov property, conditioned on the event $\bigcap_{i=0}^{n_0} \mathcal{A}_i$ we have that $|z(X_n) - z(X_{T_{n_0}})| \leq r\sigma(v)/2$ for all $n \geq T_{n_0}$ with probability at least $1/2$. Thus, we deduce that

$$\mathbf{P}_v \left(z(X_n) \in \mathfrak{H}(v, r) \text{ for all } n \geq T_{n_0} \mid \bigcap_{i=0}^{n_0} \mathcal{A}_i \right) \geq \frac{1}{2}.$$

Theorem 4.4 and the strong Markov property imply that $\mathbf{P}_v(\mathcal{A}_i \mid X_{T_{i-1}}) \geq q_1$, and hence that

$$\mathbf{P}_v \left(\bigcap_{i=0}^{n_0} \mathcal{A}_i \right) \geq q_1^{n_0},$$

which concludes our proof by taking $q_2 = q_1^{n_0}/2$. \square

Let $p, r \in (0, 1]$. We say that a set $A \subset V$ is (p, r) -**escapable** if for every vertex $v \in A$, the random walk started at v has probability at least p of not returning to A after first leaving the set of vertices whose corresponding discs in P have centres contained in the Euclidean ball of radius $r\sigma(v)$ about $z(v)$. To avoid trivialities, we also declare the empty set to be (p, r) -escapable for all p and r . Lemma 4.6 leads to the following corollary.

COROLLARY 4.7. *Let $r \in (0, 1]$. There exist positive constants $\delta_2 = \delta_2(\mathbf{M}, r)$ and $p_1 = p_1(\mathbf{M}, r)$, both increasing in r , such that the set*

$$\{v \in V : (1 - \delta_2)\varepsilon \leq \sigma(v) \leq \varepsilon\}$$

is (p_1, r) -escapable for every $\varepsilon > 0$.

Proof. Let $\eta = \eta(\mathbf{M})$ be the constant appearing in Corollary 4.2, and set

$$\delta_2 = \min\{r \sin(\eta/2)/2, 1/4\}.$$

Let $\varepsilon > 0$ be arbitrary, let $v \in \{v \in V : (1 - \delta_2)\varepsilon \leq \sigma(v) \leq \varepsilon\}$, and let T be the stopping time

$$T = \min\{n \geq 0 : |z(X_n) - z(v)| \geq r\sigma(v)\}.$$

We apply Theorem 4.4 with $r\sigma(v)$ in place of r to get

$$\mathbf{P}_v(\sigma(X_T) \leq \sigma(v) - r \sin(\eta/2)\sigma(v)) \geq q_1(\mathbf{M}) > 0. \quad (18)$$

By definition of δ_2 , on the event appearing in the left-hand side of (18), we have $\sigma(X_T) \leq \varepsilon(1 - 2\delta_2)$. Therefore, the half-plane $\mathfrak{H}(X_T, \delta_2)$, defined above Lemma 4.6, is disjoint from the ball $\{z \in \mathbb{C} : |z| \leq 1 - (1 - \delta_2)\varepsilon\}$. Thus, the claim follows from (18) and Lemma 4.6 by taking $p = q_1(\mathbf{M})q_2(\mathbf{M}, \delta_2)$, where q_2 from Lemma 4.6. \square

LEMMA 4.8. Let $A \subset V$ be (p, r) -escapable for some $p \in (0, 1)$ and $r \in (0, C_1^{-1}/2]$, where $C_1 = C_1(\mathbf{M})$ is the constant appearing in Theorem 4.3. Then for every vertex $u \in V$,

$$\mathbf{E}_u \left[\sum_{n \geq 0} a_{\mathbb{H}}(X_n) \mathbb{1}(X_n \in A) \right] \leq \frac{r^2}{p}. \quad (19)$$

Proof. Define the sequences of stopping times $\langle T_i^- \rangle_{i \geq 0}$ and $\langle T_i^+ \rangle_{i \geq 0}$ by letting $T_0^- = \tau_A$ and recursively letting

$$T_i^+ = \min\{n \geq T_i^- : |z(X_n) - z(X_{T_i^-})| \geq r\sigma(X_{T_i^-})\}$$

and $T_i^- = \min\{n \geq T_{i-1}^+ : X_n \in A\}$. Then

$$\mathbf{E}_u \left[\sum_{n \geq 0} a_{\mathbb{H}}(X_n) \mathbb{1}(X_n \in A) \right] \leq \sum_{i \geq 0} \mathbf{E}_u \left[\mathbb{1}(T_i^- < \infty) \sum_{n=T_i^-}^{T_i^+-1} a_{\mathbb{H}}(X_n) \right].$$

Since $\sigma(X_n) \geq (1-r)\sigma(X_{T_i^-})$ for all $T_i^- \leq n < T_i^+$, the Diffusive Time Estimate (Theorem 4.3), the strong Markov property, and (13) imply that

$$\begin{aligned} & \mathbf{E}_u \left[\sum_{n=T_i^-}^{T_i^+-1} a_{\mathbb{H}}(X_n) \mid T_i^- < \infty, X_{T_i^-} \right] \\ & \leq (1-r)^{-2} \sigma(X_{T_i^-})^{-2} \mathbf{E}_u \left[\sum_{n=T_i^-}^{T_i^+-1} r(X_n)^2 \mid T_i^- < \infty, X_{T_i^-} \right] \leq r^2. \end{aligned}$$

Meanwhile, since A is (p, r) -escapable, we have that $\mathbf{P}_u(T_i^- < \infty) \leq (1-p)^i$. Combining these estimates yields the desired inequality (19). \square

We say that a set $A \subset V$ is **C -short-lived** if

$$\mathbf{E}_u \left[\sum_{n \geq 0} a_{\mathbb{H}}(X_n) \mathbb{1}(X_n \in A) \right] \leq C \quad (20)$$

for every $u \in V$. Note that if A is a C -short-lived set of vertices, then every subset of A is also C -short-lived. Also note that Lemma 4.8 states that escapable sets are short-lived.

LEMMA 4.9. *Let A and B be two sets of vertices in G , and suppose that A is finite and C -short-lived for some $C > 0$. Then*

$$\sum_{v \in A} a_{\mathbb{H}}(v) \mathbf{P}_v(\tau_B < \infty) \leq C \mathcal{C}_{\text{eff}}^F(A \leftrightarrow B).$$

Proof. Let the stopping times τ_i be defined recursively by setting $\tau_0 = \tau_A$ and $\tau_{i+1} = \min\{t > \tau_i : X_t \in A\}$, so that τ_i is the i th time the random walk $\langle X_n \rangle_{n \geq 0}$ visits A . As usual if A is not hit by the random walk we set the corresponding stopping time to ∞ . Then

$$\begin{aligned} \sum_{v \in A} a_{\mathbb{H}}(v) \mathbf{P}_v(\tau_B < \infty) &\leq \sum_{v \in A} \sum_{i \geq 0} a_{\mathbb{H}}(v) \mathbf{P}_v(B \text{ hit between time } \tau_i \text{ and } \tau_{i+1}) \\ &= \sum_{v \in A} \sum_{u \in A} \sum_{i \geq 0} a_{\mathbb{H}}(v) \mathbf{P}_v(\tau_i < \infty, X_{\tau_i} = u) \mathbf{P}_u(\tau_B < \tau_A^+). \end{aligned}$$

Reversing time gives that $\mathbf{P}_v(\tau_i < \infty, X_{\tau_i} = u) = (c(u)/c(v)) \mathbf{P}_u(\tau_i < \infty, X_{\tau_i} = v)$ and since the degrees are bounded and conductances are bounded above and below we get that

$$\sum_{v \in A} a_{\mathbb{H}}(v) \mathbf{P}_v(\tau_B < \infty) \leq \sum_{v \in A} \sum_{u \in A} \sum_{i \geq 0} a_{\mathbb{H}}(v) \mathbf{P}_u(\tau_i < \infty, X_{\tau_i} = v) \mathbf{P}_u(\tau_B < \tau_A^+).$$

By exchanging the order of summation and using our assumption on A we obtain that

$$\begin{aligned} \sum_{v \in A} a_{\mathbb{H}}(v) \mathbf{P}_v(\tau_B < \infty) &\leq \sum_{u \in A} \mathbf{E}_u \left[\sum_{n \geq 0} a_{\mathbb{H}}(X_n) \mathbb{1}(X_n \in A) \right] \mathbf{P}_u(\tau_B < \tau_A^+) \\ &\leq C \sum_{u \in A} \mathbf{P}_u(\tau_B < \tau_A^+). \end{aligned}$$

To conclude, let $\langle V_j \rangle_{j \geq 1}$ be an exhaustion of V and let G_j be the subgraph of G induced by V_j , with conductances inherited from G , and observe that

$$\begin{aligned} \sum_{u \in A} c(u) \mathbf{P}_u^G(\tau_B < \tau_A^+) &= \lim_{j \rightarrow \infty} \sum_{u \in A} c(u) \mathbf{P}_u^{G_j}(\tau_B < \min\{\tau_A^+, \tau_{V \setminus V_j}\}) \\ &\leq \lim_{j \rightarrow \infty} \sum_{u \in A} c(u) \mathbf{P}_u^{G_j}(\tau_B < \tau_A^+) = \mathcal{C}_{\text{eff}}^F(A \leftrightarrow B), \end{aligned}$$

concluding the proof since $c(u)$ is bounded away from 0. □

Recall that $\text{diam}_{\mathbb{C}}(A)$ and $d_{\mathbb{C}}(A, B)$ denote the Euclidean diameter of $\{z(v) : v \in A\}$ and the Euclidean distance between $\{z(v) : v \in A\}$ and $\{z(v) : v \in B\}$, respectively.

LEMMA 4.10. *Let A and B be disjoint sets of vertices in G . Then*

$$\mathcal{C}_{\text{eff}}^{\text{F}}(A \leftrightarrow B) \leq \frac{\text{diam}_{\mathbb{C}}(A)^2}{\min\{\text{diam}_{\mathbb{C}}(A), d_{\mathbb{C}}(A, B)\}^2},$$

with the convention that the right-hand side is 1 if $\text{diam}_{\mathbb{C}}(A) = 0$.

Proof. Let $D = \text{diam}_{\mathbb{C}}(A)$. If $D = 0$ then A is a single vertex v and $\mathcal{C}_{\text{eff}}^{\text{F}}(A \leftrightarrow B) \leq c(v) \leq 1$, so assume not. Recall the **extremal length** characterization of the free effective conductance [38, Exercise 9.42]: For each function $\ell : E \rightarrow [0, \infty)$ assigning a nonnegative length to every edge e of G , let d_{ℓ} be the shortest path pseudometric on G induced by ℓ . Then

$$\mathcal{C}_{\text{eff}}^{\text{F}}(A \leftrightarrow B) = \inf \left\{ \frac{\sum_{e \in E} c(e)\ell(e)^2}{d_{\ell}(A, B)^2} : \ell : E \rightarrow [0, \infty), d_{\ell}(A, B) > 0 \right\}.$$

Let W be the set of vertices v of G whose corresponding circles intersect the D -neighbourhood of $z(A)$ in \mathbb{C} . Define lengths by setting $\ell(e)$ to be

$$\ell(e) = \begin{cases} \min\{|z(e^-) - z(e^+)|, D\} & \text{if } e \text{ has an endpoint in } W \\ 0 & \text{otherwise.} \end{cases}$$

Then

$$d_{\ell}(A, B) \geq \min\{D, d_{\mathbb{C}}(A, B)\} > 0 \quad (21)$$

while, since $|z(e^-) - z(e^+)| = r(e^-) + r(e^+)$,

$$\sum_e c(e)\ell(e)^2 \leq \sum_{v \in W} \sum_{v' \sim v} \min\{r(v) + r(v'), D\}^2 \leq \sum_{v \in W} \min\{r(v), D\}^2, \quad (22)$$

where the Ring lemma (Theorem 4.1) is used in the second inequality. As in the proof of Lemma 3.4, consider replacing each circle corresponding to a vertex in W that has radius larger than D with a circle of radius D that is contained in the original circle and intersects the D -neighbourhood of $z(A)$. This yields a set of circles contained in the $3D$ -neighbourhood of $z(A)$. Comparing the total area of this set of circles with that of the $3D$ -neighbourhood of $z(A)$ yields that

$$\sum_{v \in W} \min\{r(v), D\}^2 \leq D^2. \quad (23)$$

We conclude by combining (21), (22) and (23). □

4.4. Wilson’s algorithm. **Wilson’s algorithm rooted at infinity** [10, 54] is a powerful method of sampling the WUSF of an infinite, transient graph by joining together loop-erased random walks. We now give a very brief description of the algorithm. See [38] for a detailed exposition. Let G be a transient network. Let γ be a path in G that visits each vertex of G at most finitely many times. The **loop erasure** is formed by erasing cycles from γ chronologically as they are created. (The loop erasure of a random walk path is referred to as **loop-erased random walk** and was first studied by Lawler [35].) Let $\{v_j : j \in \mathbb{N}\}$ be an enumeration of the vertices of G . Let $\mathfrak{F}_0 = \emptyset$ and define a sequence of forests in G as follows:

- (1) Given \mathfrak{F}_i , start an independent random walk from v_{i+1} . Stop this random walk if it hits the set of vertices already included in \mathfrak{F}_i , running it forever otherwise.
- (2) Form the loop erasure of this random walk path and let \mathfrak{F}_{i+1} be the union of \mathfrak{F}_i with this loop-erased path.

Then the forest $\mathfrak{F} = \bigcup_{i \geq 0} \mathfrak{F}_i$ is a sample of the WUSF of G [10, Theorem 5.1].

4.5. Proof of Theorems 1.3, 1.4 and 1.5. Recall that $e = (x, y)$ is a fixed edge. We write \geq_e to denote a lower bound that holds up to a positive multiplicative constant that depends only on \mathbf{M} and $r_{\mathbb{H}}(x)$, and that is increasing in $r_{\mathbb{H}}(x)$.

Let (P, P^\dagger) be the double circle packing of G in \mathbb{D} normalized so that $z(x)$ and $z(y)$ lie on the negative and positive axes, respectively, and so that the tangency point of $P(x)$ and $P(y)$ is the origin. By the Ring lemma (Theorem 4.1) there exists a constant $s = s(\mathbf{M}) < 1$ such that for any $\varepsilon \in (0, 1)$ and any vertex v with $|z(v)| \geq 1 - \varepsilon$ which has a neighbour u with $|z(u)| < 1 - \varepsilon$ we must have that $|z(v)| \leq 1 - s\varepsilon$. It follows that in any simple path starting at a vertex u with $|z(u)| < 1 - \varepsilon$, the first time we visit a vertex v with $|z(v)| \geq 1 - \varepsilon$ we must have that $|z(v)| \geq 1 - s\varepsilon$. Thus, for every $\varepsilon > 0$, the set

$$\{v \in V : z(v) \in A_0(1 - \varepsilon, 1 - s\varepsilon)\}$$

disconnects e from ∞ in G . An equivalent formulation that we will use is that if (v, u) is an edge in the graph and $0 < \varepsilon_1 < \varepsilon_2$, then

$$\sigma(v) \in [\varepsilon_1, \varepsilon_2] \implies \sigma(u) \in [s\varepsilon_1, s^{-1}\varepsilon_2]. \quad (24)$$

For each $\varepsilon > 0$, we define

$$W_\varepsilon := \{v \in V : z(v) \in A_0(1 - \varepsilon, 1 - s^5\varepsilon)\}.$$

Recall the definition of a C -short-lived set from (20). We trivially have that unions of boundedly many short-lived sets are short-lived with a larger constant. In particular, letting $s = s(\mathbf{M})$ be as above and letting $\delta = \delta_2(\mathbf{M}, C_1^{-1}(\mathbf{M}))$, where C_1 is the constant appearing in Theorem 4.3 and δ_2 is the constant from Corollary 4.7, we have that

$$W_\varepsilon \subseteq \bigcup_{m=0}^{\lceil 5 \log_{1-\delta}(s) \rceil} \{v \in V : (1-\delta)^{m+1} \varepsilon \leq \sigma(v) \leq (1-\delta)^m \varepsilon\}$$

for every $\varepsilon > 0$. Applying Corollary 4.7 and Lemma 4.8 immediately yields the following.

COROLLARY 4.11. *There exists a constant $C_2 = C_2(\mathbf{M}) > 0$ such that the set W_ε is C_2 -short-lived for every $\varepsilon > 0$.*

Let \mathfrak{F} be the WUSF of G and let $\varepsilon \in (0, 1/2)$ be arbitrarily small. The proofs in this section will be based upon the study of the random variable

$$Z_\varepsilon = \sum_{v \in W_\varepsilon} a_{\mathbb{H}}(v) \mathbb{1}(v \in \text{past}_{\mathfrak{F}}(e)), \quad (25)$$

and its moments. Intuitively, we think of the process $\langle Z_{2^{-n}} \rangle_{n \geq 1}$ as behaving similarly to a critical branching process, so that in particular we expect the random variable Z_ε to be distributed similarly to the number of particles at generation $\log(1/\varepsilon)$ in a critical branching process. In particular, we expect its first moment to be of order 1 and its second moment to be of order $\log(1/\varepsilon)$. This is what we prove below. We begin by bounding its first moment.

LEMMA 4.12. $\mathbb{E}[Z_\varepsilon] \leq 1$ for all $\varepsilon > 0$.

Proof. If we generate \mathfrak{F} using Wilson's algorithm, starting with the vertex v , then v is in $\text{past}_{\mathfrak{F}}(e)$ if and only if the loop-erased random walk from v passes through $e = (x, y)$. In particular $\mathbb{P}(v \in \text{past}_{\mathfrak{F}}(e)) \leq \mathbf{P}_v(\tau_x < \infty)$. Thus, the claim follows immediately from Corollary 4.11 and Lemma 4.9, taking $A = W_\varepsilon \setminus \{x\}$ and $B = \{x\}$, since $\mathcal{C}_{\text{eff}}^{\text{F}}(x \leftrightarrow V \setminus \{x\}) \leq c(x) \leq 1$. \square

LEMMA 4.13. *There exists a positive constant $\delta_5 = \delta_5(\mathbf{M}, r_{\mathbb{H}}(x))$, increasing in $r_{\mathbb{H}}(x)$, such that $\mathbb{E}[Z_\varepsilon] \geq_e 1$ for all $\varepsilon \leq \delta_5$.*

Proof. As in the previous lemma, v is in $\text{past}_{\mathfrak{F}}(e)$ if and only if the loop-erased random walk from v passes through $e = (x, y)$. In particular, we obtain the lower

bound

$$\mathbb{P}(v \in \text{past}_{\mathfrak{F}}(e)) \geq \mathbf{P}_v(\tau_x < \infty, X_{\tau_x+1} = y, \langle X_n \rangle_{n \geq \tau_x+1} \text{ disjoint from } \langle X_n \rangle_{n=0}^{\tau_x})$$

and hence, decomposing according to the value of τ_x ,

$$\mathbb{E}[Z_\varepsilon] \geq \sum_{v \in W_\varepsilon} \sum_{m \geq 1} a_{\mathbb{H}}(v) \mathbf{P}_v(\tau_x = m, X_{m+1} = y, \langle X_n \rangle_{n \geq m} \text{ disjoint from } \langle X_n \rangle_{n=0}^{m-1}).$$

Letting $\langle Y_n \rangle_{n \geq 0}$ be a random walk started at x independent of $\langle X_n \rangle_{n \geq 0}$ and reversing time yields that

$$\mathbb{E}[Z_\varepsilon] \geq \sum_{v \in W_\varepsilon} \sum_{m \geq 1} a_{\mathbb{H}}(v) \mathbf{P}_x(X_m = v, \mathcal{C}_m) = \mathbf{E}_x \left[\sum_{m \geq 0} a_{\mathbb{H}}(X_m) \mathbb{1}(X_m \in W_\varepsilon, \mathcal{C}_m) \right],$$

where \mathcal{C}_m is the event

$$\mathcal{C}_m = \{Y_1 = y, \text{ and } \langle X_n \rangle_{n=1}^m \text{ disjoint from } \langle Y_n \rangle_{n \geq 0}\}.$$

(Note that on the event \mathcal{C}_m the walk $\langle X_n \rangle_{n=0}^m$ does not return to x after time 0.)

Let τ_1 be the first time that the random walk $\langle X_m \rangle_{m \geq 0}$ visits

$$\{v \in V : s^3 \varepsilon \leq \sigma(v) \leq s^2 \varepsilon\},$$

which is finite a.s. since this set separates x from ∞ by definition of s . Let τ_2 be the first time m after τ_1 that $|z(X_m) - z(X_{\tau_1})| \geq C_1^{-1} s^3 \varepsilon$, where $C_1 = C_1(\mathbf{M}) \geq 1$ is the constant from Theorem 4.3. By the triangle inequality

$$s^3 \varepsilon - C_1^{-1} s^3 \varepsilon \leq \sigma(X_{\tau_2-1}) \leq s^2 \varepsilon + C_1^{-1} s^3 \varepsilon,$$

and hence by (24) we get that

$$s^4 \varepsilon - C_1^{-1} s^4 \varepsilon \leq \sigma(X_{\tau_2}) \leq s \varepsilon + C_1^{-1} s^2 \varepsilon.$$

We may assume that s is sufficiently small that $X_m \in W_\varepsilon$ for all $\tau_1 \leq m \leq \tau_2$. Therefore,

$$\begin{aligned} \mathbb{E}[Z_\varepsilon] &\geq \mathbf{E}_x \left[\sum_{m=\tau_1}^{\tau_2} a_{\mathbb{H}}(X_m) \mathbb{1}(\mathcal{C}_m) \right] \geq \mathbf{E}_x \left[\sum_{m=\tau_1}^{\tau_2} a_{\mathbb{H}}(X_m) \mathbb{1}(\mathcal{C}_{\tau_2}) \right] \\ &\geq \mathbf{E}_x \left[\sum_{m=\tau_1}^{\tau_2} a_{\mathbb{H}}(X_m) \mathbb{1}(d_C(X_{\tau_1}, \{Y_n : n \geq 0\}) \geq 2\varepsilon/s \text{ and } \mathcal{C}_{\tau_2}) \right], \end{aligned}$$

where the second inequality follows since the events \mathcal{C}_m are decreasing and third inequality is trivial. We have that

$$|z(X_{\tau_2-1}) - z(X_{\tau_1})| \leq C_1^{-1} s^3 \varepsilon,$$

and since $X_{\tau_1-1} \in W_\varepsilon$ we have

$$|z(X_{\tau_2}) - z(X_{\tau_2-1})| \leq \varepsilon/s,$$

by the definition of s . By the triangle inequality $|z(X_{\tau_2}) - z(X_{\tau_1})| \leq 2\varepsilon/s$ since $s \leq 1/2$. We deduce that the events $\mathcal{C}_{\tau_1} \cap \{d_{\mathbb{C}}(X_{\tau_1}, \{Y_n : n \geq 0\}) \geq 2\varepsilon/s\}$ and $\mathcal{C}_{\tau_2} \cap \{d_{\mathbb{C}}(X_{\tau_1}, \{Y_n : n \geq 0\}) \geq 2\varepsilon/s\}$ are equal. Therefore,

$$\begin{aligned} \mathbb{E}[Z_\varepsilon] &\geq \mathbf{E}_x \left[\sum_{m=\tau_1}^{\tau_2} a_{\mathbb{H}}(X_m) \mathbb{1}(d_{\mathbb{C}}(X_{\tau_1}, \{Y_n : n \geq 0\}) \geq 2\varepsilon/s, \mathcal{C}_{\tau_1}) \right] \\ &= \mathbf{E}_x \left[\mathbf{E}_x \left[\sum_{m=\tau_1}^{\tau_2} a_{\mathbb{H}}(X_m) \mid X_{\tau_1} \right] \mathbb{1}(d_{\mathbb{C}}(X_{\tau_1}, \{Y_n : n \geq 0\}) \geq 2\varepsilon/s, \mathcal{C}_{\tau_1}) \right]. \end{aligned}$$

Applying (13) and the Diffusive Time Estimate (Theorem 4.3), we obtain that

$$\mathbf{E}_x \left[\sum_{m=\tau_1}^{\tau_2} a_{\mathbb{H}}(X_m) \mid X_{\tau_1} \right] \asymp \varepsilon^{-2} \mathbf{E}_x \left[\sum_{m=\tau_1}^{\tau_2} r(X_m)^2 \mid X_{\tau_1} \right] \asymp 1,$$

and hence

$$\mathbb{E}[Z_\varepsilon] \geq \mathbf{P}_x(d_{\mathbb{C}}(X_{\tau_1}, \{Y_n : n \geq 0\}) \geq 2\varepsilon/s, \mathcal{C}_{\tau_1}). \quad (26)$$

We now set

$$\delta_5 = \frac{s(r(x) + r(y))}{4},$$

and note that, as usual, the Ring lemma implies that $\delta_5 \asymp r_{\mathbb{H}}(x)$. Since the degrees, the edge conductances, and their inverses are bounded by \mathbf{M} , the probability that $Y_1 = y$ is at least \mathbf{M}^{-3} and we may assume this indeed occurs. We now apply Lemma 4.6 with $r_1 = r(x)/2\sigma(x)$ (note that by the Ring lemma $\sigma(x) \geq 1$ and hence $r_1 \geq r_{\mathbb{H}}(x)$) and establish that with probability at least q_2 the walker $\{X_n\}_{n \geq 0}$ is confined to the half-plane $\mathfrak{H}(x, r_1)$ (which is just $\{(x, y) : x \leq -r(x)/2\}$). Similarly, we put $r_2 = r(y)/2\sigma(y)$ and by Lemma 4.6 we get that with probability at least q_2 the walker $\{Y_n\}_{n \geq 1}$ starting from y is confined to the half-plane $\mathfrak{H}(y, r_2)$ (which is $\{(x, y) : x \geq r(y)/2\}$). If these two events occur, then the distance between $z(X_n)$ and $z(Y_m)$ is at least $(r(x) + r(y))/2$ for any $n \geq 0$ and $m \geq 1$. By our choice of δ_5 we deduce that for any $\varepsilon \leq \delta_5$

$$\mathbf{P}_x(d_{\mathbb{C}}(X_{\tau_1}, \{Y_n : n \geq 0\}) \geq 2\varepsilon/s, \mathcal{C}_{\tau_1}) \geq \mathbf{M}^{-1} q_2 (\mathbf{M}, r_{\mathbb{H}}(x))^2 \geq_e 1,$$

concluding the proof. \square

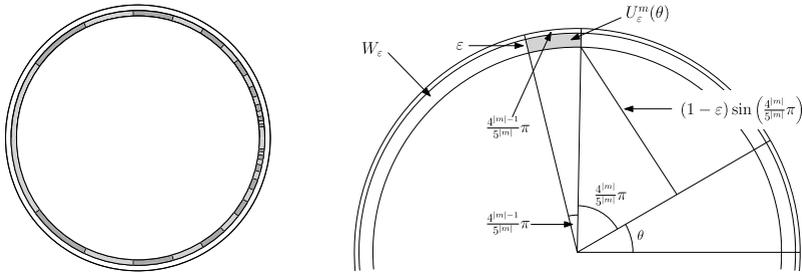


Figure 5. Left: The decomposition of W_ε that it used to bound the second moment of Z_ε . Each shaded region contributes at most a constant to the conditional first moment. Right: Illustration of the trigonometric calculations used to prove Lemma 4.14.

We now wish to estimate the second moment of Z_ε . Doing so amounts to estimating the conditional first moment of Z_ε given that some specified vertex of W_ε is contained in the past of e . We will show that this conditional first moment is $O(\log(1/\varepsilon))$ by decomposing the annulus W_ε into $O(\log(1/\varepsilon))$ many pieces, the sets U_ε^m below, and then proving that each such piece contributes at most a constant to the conditional first moment. Using the bounds established in Section 4.3 will reduce the problem of bounding the contribution to the conditional first moment of each piece to geometric estimates that can be proven by elementary trigonometry.

We now begin to carry out this programme. For each $\varepsilon > 0$ and $m \in \mathbb{Z} \setminus \{0\}$, we define

$$U_\varepsilon^m = \left\{ z \in \mathbb{D} : 1 - \varepsilon \leq |z| \leq 1 - s^3 \varepsilon \text{ and } \operatorname{sgn}(m) \frac{4^{|m|}}{5^{|m|}} \pi \leq \arg z \leq \operatorname{sgn}(m) \frac{4^{|m|-1}}{5^{|m|-1}} \pi \right\}$$

and define $U_\varepsilon^m(\theta)$ to be the rotated set $e^{i\theta} U_\varepsilon^m$.

LEMMA 4.14. *There exist universal constants $\delta_3, \delta_4 > 0$ and $k < \infty$ such that*

$$d_{\mathbb{C}}(U_\varepsilon^m(\theta), \{re^{i\theta} : r \geq 0\}) \geq (2 + \delta_3) \operatorname{diam}_{\mathbb{C}}(U_\varepsilon^m(\theta))$$

for all $\theta \in [-\pi, \pi]$, $\varepsilon \leq \delta_4$ and $|m| \leq \log_{5/4}(1/\varepsilon) - k$.

The constants here are more important than usual since we will later need to estimate the difference $\frac{1}{2} d_{\mathbb{C}}(U_\varepsilon^m(\theta), \{re^{i\theta} : r \geq 0\}) - \operatorname{diam}_{\mathbb{C}}(U_\varepsilon^m(\theta))$.

Proof. We begin by calculating, using elementary trigonometry, that

$$\text{diam}_{\mathbb{C}}(U_{\varepsilon}^m(\theta)) \leq \frac{4^{|m|-1}}{5^{|m|}}\pi + \varepsilon$$

and

$$d_{\mathbb{C}}(U_{\varepsilon}^m(\theta), \{re^{i\theta} : r \geq 0\}) = \begin{cases} 1 - \varepsilon & |m| \leq 3, \\ (1 - \varepsilon) \sin\left(\frac{4^{|m|}}{5^{|m|}}\pi\right) & |m| \geq 4. \end{cases}$$

See Figure 5 for an illustration. By concavity of the function $\sin(t)$ on $[0, \pi]$, we have that $\sin(t) \geq 2t/\pi$ for all $t \in [0, \pi/2]$, and in particular

$$\begin{aligned} & \frac{\text{diam}_{\mathbb{C}}(U_{\varepsilon}^m(\theta))}{d_{\mathbb{C}}(U_{\varepsilon}^m(\theta), \{re^{i\theta} : r \geq 0\})} \\ & \leq \frac{1}{8(1 - \varepsilon)}\pi + \frac{5^{|m|}\varepsilon}{2 \cdot 4^{|m|}(1 - \varepsilon)} \leq \frac{1}{8(1 - \delta_4)}\pi + \frac{4^k\delta_4}{2 \cdot 5^k(1 - \delta_4)} \end{aligned}$$

for all $\varepsilon \leq \delta_4$ and $|m| \leq \log_{5/4}(1/\varepsilon) - k$. This upper bound is less than $\pi/7 < 1/2$ when δ_4 is sufficiently small and k is sufficiently large. \square

LEMMA 4.15. $\mathbb{E}[Z_{\varepsilon}^2] \leq \mathbb{E}[Z_{\varepsilon}] \log(1/\varepsilon)$ for all $0 < \varepsilon \leq \delta_4$, where δ_4 is the constant from Lemma 4.14.

Again, we remark that $\log(1/\varepsilon)$ is of the same order as the hyperbolic distance between e and W_{ε} .

Proof. For each two vertices u and v in G , let

$$H(u, v) := \{w \in V : |z(w) - z(u)| \leq |z(w) - z(v)|\}$$

be the set of vertices closer to u than to v with respect to the Euclidean metric on the circle packing. Expand $\mathbb{E}[Z_{\varepsilon}^2]$ as the sum

$$\mathbb{E}[Z_{\varepsilon}^2] = \sum_{u, v \in W_{\varepsilon}} a_{\mathbb{H}}(u)a_{\mathbb{H}}(v)\mathbb{P}(u, v \in \text{past}_{\mathfrak{F}}(e)).$$

If u and v are both in the past of e in \mathfrak{F} , let $w(u, v)$ be the first vertex at which the unique simple paths in \mathfrak{F} from u to e and from v to e meet. Then

$$\mathbb{E}[Z_{\varepsilon}^2] \leq 2 \sum_{u, v \in W_{\varepsilon}} a_{\mathbb{H}}(u)a_{\mathbb{H}}(v)\mathbb{P}(u, v \in \text{past}_{\mathfrak{F}}(e) \text{ and } w(u, v) \in H(u, v)).$$

Consider generating \mathfrak{F} using Wilson’s algorithm rooted at infinity, starting first with u and then v . In order for u and v both to be in the past of e and for $w(u, v)$ to be in $H(u, v)$, we must have that the simple random walk from v hits $H(u, v)$, so that

$$\mathbb{E}[Z_\varepsilon^2] \leq 2 \sum_{u \in W_\varepsilon} a_{\mathbb{H}}(u) \mathbb{P}(u \in \text{past}_{\mathfrak{F}}(e)) \sum_{v \in W_\varepsilon} a_{\mathbb{H}}(v) \mathbf{P}_v(\tau_{H(u,v)} < \infty). \tag{27}$$

Thus, it suffices to prove that

$$\sum_{v \in W_\varepsilon} a_{\mathbb{H}}(v) \mathbf{P}_v(\tau_{H(u,v)} < \infty) \leq \log(1/\varepsilon) \tag{28}$$

for all $\varepsilon \leq \delta_4$ and $u \in W_\varepsilon$.

Recall the definition of U_ε^m from Lemma 4.14. Let k be the universal constant from Lemma 4.14, let $\ell(\varepsilon) = \lceil \log_{5/4}(\varepsilon) - k \rceil$, and let $\theta = \arg(u)$. For each $m \in \mathbb{Z}$ with $1 \leq |m| \leq \ell(\varepsilon)$, let $S_\varepsilon^m(\theta)$ be the set of vertices whose centres are contained in $U_\varepsilon^m(\theta)$, and let

$$S_\varepsilon^0 = S_\varepsilon^0(u) := \left\{ v \in W_\varepsilon : \left| \arg \frac{z(v)}{z(u)} \right| \leq \frac{5^k \pi \varepsilon}{4^k} \right\}.$$

Then

$$\begin{aligned} \sum_{v \in W_\varepsilon} a_{\mathbb{H}}(v) \mathbf{P}_v(\tau_{H(u,v)} < \infty) &= \sum_{m=-\ell(\varepsilon)}^{\ell(\varepsilon)} \sum_{v \in S_\varepsilon^m} a_{\mathbb{H}}(v) \mathbf{P}_v(\tau_{H(u,v)} < \infty) \\ &\leq \sum_{m=-\ell(\varepsilon)}^{\ell(\varepsilon)} \sum_{v \in S_\varepsilon^m} a_{\mathbb{H}}(v) \mathbf{P}_v(\tau_{H(u, S_\varepsilon^m)} < \infty), \end{aligned} \tag{29}$$

where for a set $A \subset V$ and a vertex u we denote $H(u, A) = \bigcup_{v \in A} H(u, v)$.

Since S_ε^m is contained in W_ε , it is C -short-lived for some $C = C(\mathbf{M})$ by Corollary 4.11. Thus, applying Lemmas 4.9 and 4.10 yields that

$$\sum_{v \in S_\varepsilon^m} a_{\mathbb{H}}(v) \mathbf{P}_v(\tau_{H(u, S_\varepsilon^m)} < \infty) \leq \frac{\text{diam}_{\mathbb{C}}(S_\varepsilon^m)^2}{\min(\text{diam}(S_\varepsilon^m), d_{\mathbb{C}}(S_\varepsilon^m, H(u, S_\varepsilon^m)))^2}. \tag{30}$$

Our goal is to show that the last ratio is ≤ 1 . If the minimum in the denominator is attained on the first term, then we are done. Otherwise, we bound

$$\begin{aligned} d_{\mathbb{C}}(S_{\varepsilon}^m, H(u, S_{\varepsilon}^m)) &= \inf_{v_1, v_2 \in S_{\varepsilon}^m} d_{\mathbb{C}}(v_1, H(u, v_2)) \\ &\geq \inf_{v_1 \in S_{\varepsilon}^m} d_{\mathbb{C}}(v_1, H(u, v_1)) - \sup_{v_1, v_2 \in S_{\varepsilon}^m} d_{\mathbb{C}}(v_1, v_2) \\ &\geq \frac{1}{2} d_{\mathbb{C}}(u, S_{\varepsilon}^m) - \text{diam}_{\mathbb{C}}(S_{\varepsilon}^m). \end{aligned}$$

Since $S_{\varepsilon}^m \subset U_{\varepsilon}^m$ have that $\text{diam}_{\mathbb{C}}(S_{\varepsilon}^m) \leq \text{diam}_{\mathbb{C}}(U_{\varepsilon}^m)$ and furthermore $d_{\mathbb{C}}(u, S_{\varepsilon}^m) \geq d_{\mathbb{C}}(U_{\varepsilon}^m(\theta), \{re^{i\theta} : r \geq 0\})$. Thus Lemma 4.14 implies that for $\varepsilon \leq \delta_4$ and $1 \leq |m| \leq \ell(\varepsilon)$ we have that

$$d_{\mathbb{C}}(S_{\varepsilon}^m, H(u, S_{\varepsilon}^m)) \geq \text{diam}_{\mathbb{C}}(U_{\varepsilon}^m).$$

We get that the right-hand side of (30) is ≤ 1 . We put this in (29) and use the fact that $\sum_{v \in S_{\varepsilon}^0} a_{\mathbb{H}}(v) \leq 1$ to bound the term $m = 0$ there. This gives (28) and concludes the proof. \square

We now have everything we need in place to prove Theorem 1.4.

Proof of Theorem 1.4. Denote by $\mathcal{R}_{\varepsilon}^e$ the event that $\text{past}_{\mathfrak{F}}^e(e)$ contains a vertex whose centre is within Euclidean distance ε of the unit circle, and let \mathcal{D}_R^e be the event that $\text{diam}_{\mathbb{H}}(\text{past}_{\mathfrak{F}}^e(e)) \geq R$. Then, letting $\varepsilon(R) = 1 - \tanh(R/2)$, we have that $\mathcal{R}_{\varepsilon(R)/2}^e \subseteq \mathcal{D}_R^e \subseteq \mathcal{R}_{\varepsilon(R)}^e$. Thus, to prove Theorem 1.4, it suffices to prove that

$$\log(1/\varepsilon)^{-1} \leq_e \mathbb{P}(\mathcal{R}_{\varepsilon}^e) \leq \log(1/\varepsilon)^{-1} \quad (31)$$

for all $\varepsilon \leq 1/2$.

As discussed at the beginning of Section 3, the proof of Theorem 1.2, and in particular of the estimate (10), adapts immediately to yield the desired upper bound when the analysis there is carried out using the double circle packing (P, P^{\dagger}) of G rather than the circle packing of $T(G)$. Indeed, we replace $P(f)$ with $P^{\dagger}(f)$ in the definition of $V_z(r, r')$ above Lemma 3.3 and observe that part (2) of Lemma 3.3 holds by the Ring lemma 2.8 since G has bounded degrees. Furthermore, the proof of Lemma 3.4 applies as written once we replace $P(f)$ with $P^{\dagger}(f)$. The proof of Theorem 1.2 now follows almost verbatim; we only note that the procedure of diverting the random path from T to G works seamlessly in the double circle packing setting and (10) is obtained.

The remainder of this proof is devoted to proving the lower bound. Towards that aim, note that by the definition of W_{ε} , the random variable Z_{ε} defined in (25) is positive if and only if $\mathcal{R}_{\varepsilon}^e$ occurs. Let δ_4 be the constant from Lemma 4.14 and

let $\delta_5 = \delta_5(\mathbf{M}, r_{\mathbb{H}}(x))$ be the constant from Lemma 4.13. The lower bound of Theorem 1.4 now follows from Lemmas 4.13 and 4.15 together with the Cauchy–Schwartz inequality, which imply that

$$\mathbb{P}(Z_\varepsilon > 0) \geq \frac{\mathbb{E}[Z_\varepsilon]^2}{\mathbb{E}[Z_\varepsilon^2]} \geq_e \frac{1}{\log(1/\varepsilon)}$$

for all $\varepsilon \leq \min\{\delta_4, \delta_5\}$. Since $\mathbb{P}(Z_\varepsilon > 0)$ is an increasing function of ε and $\min\{\delta_4, \delta_5\}$ is an increasing function of $r_{\mathbb{H}}(x)$, it follows that

$$\mathbb{P}(Z_\varepsilon > 0) \geq \mathbb{P}(Z_{\min\{\delta_4, \delta_5\}\varepsilon} > 0) \geq_e \frac{1}{\log(1/\varepsilon)}$$

for all $\varepsilon \leq 1/2$. □

Proof of Theorem 1.5. We continue to use the notation from the proof of Theorem 1.4. We first prove the upper bound. Let $R \geq 1$. Applying Markov’s inequality and Lemma 4.12 we obtain that

$$\mathbb{P}\left(\sum_{k=0}^{R^{1/2}} Z_{s^{-3k}} \geq R\right) \leq \frac{1}{R} \mathbb{E}\left[\sum_{k=0}^{R^{1/2}} Z_{s^{-3k}}\right] \leq R^{-1/2}, \tag{32}$$

and hence

$$\mathbb{P}(\text{area}_{\mathbb{H}}(\text{past}_{\mathfrak{F}}(e)) \geq R) \leq \mathbb{P}\left(\sum_{k=0}^{R^{1/2}} Z_{s^{-3k}} \geq R\right) + \mathbb{P}\left(\sum_{k=R^{1/2}}^{\infty} Z_{s^{-3k}} > 0\right) \leq R^{-1/2}, \tag{33}$$

where the second inequality follows from Theorem 1.4 and (32).

We now obtain a matching lower bound. Let δ_4 be the constant from Lemma 4.14 and let $\delta_5 = \delta_5(\mathbf{M}, r_{\mathbb{H}}(x))$ be the constant from Lemma 4.13, and let

$$R_0 = R_0(\mathbf{M}, r_{\mathbb{H}}(x)) = (4/9) \log_{1/s}^2(1/\min\{\delta_4, \delta_5\}),$$

so that $s^{-3R^{1/2}/2}$ is less than both δ_4 and δ_5 for all $R \geq R_0$. Let $R \geq R_0$ and let

$$W = \bigcup_{k=(1/2)R^{1/2}}^{R^{1/2}} W_{s^{-3k}} \quad \text{and} \quad \text{let } Z = \sum_{k=(1/2)R^{1/2}}^{R^{1/2}} Z_{s^{-3k}}.$$

The argument used to derive (27) is imitated and yields that

$$\mathbb{E}[Z^2] \leq 2 \sum_{u \in W} a_{\mathbb{H}}(u) \mathbb{P}(u \in \text{past}_{\mathfrak{F}}(e)) \sum_{v \in W} a_{\mathbb{H}}(v) \mathbf{P}_v(\tau_{H(u,v)} < \infty). \tag{34}$$

Let $u \in W$ and let $\theta = \arg(z(u))$. Let $\ell(s^{-3k})$ and let the sets $S_{s^{-3k}}^m(\theta)$ be defined as in the proof of Lemma 4.15. Then,

$$\sum_{v \in W} a_{\mathbb{H}}(v) \mathbf{P}_v(\tau_{H(u,v)} < \infty) \leq \sum_{k=(1/2)R^{1/2}}^{R^{1/2}} \sum_{m=-\ell(s^{-3k})}^{\ell(s^{-3k})} \left[\sum_{v \in S_{s^{-3k}}^m} a_{\mathbb{H}}(v) \mathbf{P}(\tau_{H(u, S_{s^{-3k}}^m)} < \infty) \right]. \tag{35}$$

The proof of Lemma 4.15 showed that the left-hand side of (30) is bounded by a constant depending only on \mathbf{M} for any ε , and it follows that this also holds for the sums appearing in the square brackets on the right-hand side of (35). Thus, we deduce that

$$\sum_{v \in W} a_{\mathbb{H}}(v) \mathbf{P}_v(\tau_{H(u,v)} < \infty) \leq R \tag{36}$$

for all $u \in W$. Putting (34) and (36) together yields that

$$\mathbb{E}[Z^2] \leq R\mathbb{E}[Z].$$

Next, Lemma 4.13 implies that $\mathbb{E}[Z] \geq_e R^{1/2}$, while Theorem 1.4 implies that $\mathbb{P}(Z > 0) \leq R^{-1/2}$. Thus, there exists a positive constant $C = C(\mathbf{M}, r_{\mathbb{H}}(x))$ such that $\mathbb{E}[Z \mid Z > 0] \geq CR$. Applying the second moment estimate above, the Paley–Zigmond Inequality implies that

$$\mathbb{P}\left(Z \geq \frac{C}{2}R \mid Z > 0\right) \geq \frac{\mathbb{E}[Z \mid Z > 0]^2}{4\mathbb{E}[Z^2 \mid Z > 0]} = \frac{\mathbb{E}[Z]^2}{4\mathbb{E}[Z^2]\mathbb{P}(Z > 0)} \geq_e 1. \tag{37}$$

Combining (37) with the lower bound of Theorem 1.4 yields that

$$\mathbb{P}\left(\text{area}_{\mathbb{H}}(\text{past}_{\mathfrak{F}}(e)) \geq \frac{C}{2}R\right) \geq \mathbb{P}\left(Z \geq \frac{C}{2}R\right) \geq_e R^{-1/2}$$

for all $R \geq R_0$. Since the probability on the left-hand side is decreasing in R , we conclude that

$$\mathbb{P}(\text{area}_{\mathbb{H}}(\text{past}_{\mathfrak{F}}(e)) \geq R) \geq_e R^{-1/2}$$

as claimed. □

Proof of Theorem 1.3. Let \mathfrak{F} be a spanning forest of G and let e^\dagger be the edge of G^\dagger dual to $e = (x, y)$. The past of e^\dagger in the dual forest \mathfrak{F}^\dagger is contained in the region of the plane bounded by e and $\Gamma_{\mathfrak{F}}(x, y)$, so that

$$\text{diam}_{\mathbb{H}}(\Gamma_{\mathfrak{F}}(x, y)) \geq \text{diam}_{\mathbb{H}}(\text{past}_{\mathfrak{F}^\dagger}(e^\dagger)).$$

Meanwhile, if $\text{past}_{\mathfrak{F}^\dagger}(e^\dagger)$ is nonempty, then every edge in the path $\Gamma_{\mathfrak{F}}(x, y)$ is incident to a face of G that is in $\text{past}_{\mathfrak{F}^\dagger}(e^\dagger)$. By the Ring lemma (Theorem 4.1),

the hyperbolic radii of circles in $P \cup P^\dagger$ are bounded above. We deduce that there exists a constant $C = C(\mathbf{M})$ such that

$$\text{diam}_{\mathbb{H}}(\Gamma_{\mathfrak{F}}(x, y)) \leq \text{diam}_{\mathbb{H}}(\text{past}_{\mathfrak{F}^\dagger}(e^\dagger)) + C.$$

We deduce Theorem 1.3 from Theorem 1.4 by applying these estimates when \mathfrak{F} is the FUSF of G and \mathfrak{F}^\dagger is the WUSF of G^\dagger . □

4.6. The uniformly transient case. Recall that a graph is said to be **uniformly transient** if $\mathbf{p} = \inf_{v \in V} \mathbf{P}_v(\tau_v^+ = \infty)$ is positive. If in addition the graph has bounded degrees, this is equivalent to the property that $\mathcal{C}_{\text{eff}}(v \leftrightarrow \infty)$ is bounded away from zero uniformly in v .

PROPOSITION 4.16. *Then there exists a constant $C = C(\mathbf{M})$ such that*

$$r_{\mathbb{H}}(v) \geq \frac{1}{C} \exp(-C \mathcal{R}_{\text{eff}}(v \leftrightarrow \infty)).$$

Proof. By applying a Möbius transformation if necessary, we may assume that the circle $P(v)$ is centred at the origin. By the Ring lemma (Theorem 4.1), $r_{\mathbb{H}}(v)$ is bounded above by a constant depending only on the maximum degree and codegree of G . Applying [18, Corollary 3.3] together with Theorem 4.1 yields that $\mathcal{R}_{\text{eff}}(v \leftrightarrow \infty) \geq c \log(1/r(v))$ for some constant $c = c(\mathbf{M})$. Since the hyperbolic radii are bounded from above, the Euclidean radius $r(v)$ is comparable to $r_{\mathbb{H}}(v)$. □

We are now ready to prove Corollary 1.6. In the rest of this subsection, we will use \leq, \geq and \asymp to denote inequalities or equalities that hold up to positive multiplicative constants depending only on \mathbf{M} and \mathbf{p} .

Proof of Corollary 1.6. Let $e = (x, y)$ be an edge of G and let \mathfrak{F} be sample of FUSF $_G$. By Theorem 2.5 the dual forest \mathfrak{F}^\dagger is distributed as WUSF $_G$ and so by Theorem 1.2 it is a.s. one-ended. Hence $\text{past}_{\mathfrak{F}^\dagger}(e^\dagger)$ (recall the definition above Theorem 1.4), where e^\dagger the edge of G^\dagger dual to e , is well defined.

The past of e^\dagger in the dual forest \mathfrak{F}^\dagger is contained in the region of the plane bounded by e and $\Gamma_{\mathfrak{F}}(e)$. The nonamenability of the hyperbolic plane implies that the perimeter of any set is at least a constant multiple of its area, and so

$$\sum_{v \in \Gamma_{\mathfrak{F}}(x, y)} r_{\mathbb{H}}(v) \geq \text{area}_{\mathbb{H}}(\text{past}_{\mathfrak{F}^\dagger}(e^\dagger)). \tag{38}$$

On the other hand, if $\text{past}_{\mathfrak{F}^\dagger}(e^\dagger)$ is nonempty, then every edge in the path $\Gamma_{\mathfrak{F}}(x, y)$ is incident to a face of G that is in $\text{past}_{\mathfrak{F}^\dagger}(e^\dagger)$. We deduce that if

$\text{past}_{\mathfrak{F}^\dagger}(e^\dagger)$ is nonempty then

$$\text{area}_{\mathbb{H}}(\text{past}_{\mathfrak{F}^\dagger}(e^\dagger)) \geq \sum_{v \in \Gamma_{\mathfrak{F}}(x, y)} a_{\mathbb{H}}(v). \quad (39)$$

Note that neither estimate (38) or (39) required uniform transience. Proposition 4.16 and the Ring lemma (Theorem 4.1) imply that

$$|\Gamma_{\mathfrak{F}}(x, y)| \asymp \sum_{v \in \Gamma_{\mathfrak{F}}(x, y)} r_{\mathbb{H}}(v) \asymp \sum_{v \in \Gamma_{\mathfrak{F}}(x, y)} a_{\mathbb{H}}(v).$$

Thus, we deduce Corollary 1.6 from Theorem 1.5. □

Let (X_1, d_1) and (X_2, d_2) be metric spaces and let α, β be positive. A (not necessarily continuous) function $\phi : X_1 \rightarrow X_2$ is said to be an (α, β) -**rough isometry** if the following hold.

- (1) (**ϕ roughly preserves distances.**) $\alpha^{-1}d_1(x, y) - \beta \leq d_2(\phi(x), \phi(y)) \leq \alpha d_1(x, y) + \beta$ for all $x, y \in X_1$.
- (2) (**ϕ is almost surjective.**) For every $x_2 \in X_2$, there exists $x_1 \in X_1$ such that $d_2(\phi(x_1), x_2) \leq \beta$.

See [38, §2.6] for further background on rough isometries. We write d_G for the graph distance on V .

COROLLARY 4.17. *Let G be a uniformly transient, polyhedral, proper plane network with bounded codegrees and bounded local geometry. Let d_G denote the graph distance on V , let (P, P^\dagger) be a double circle packing of G in \mathbb{D} , and let $z(v)$ be the centre of the disc in P corresponding to the vertex v . Then there exist positive constants $\alpha = \alpha(\mathbf{M}, \mathbf{p})$ and $\beta = \beta(\mathbf{M}, \mathbf{p})$ such that z is an (α, β) -rough isometry from (V, d_G) to $(\mathbb{D}, d_{\mathbb{H}})$.*

Proof. Proposition 4.16 implies that for every vertex v of G , $r_{\mathbb{H}}(v)$ is bounded both above and away from zero by positive constants. Almost surjectivity is immediate. For each two vertices u and v in G , the shortest graph distance path between them induces a curve in \mathbb{D} (by going along the hyperbolic geodesics between the centres of the circles in the path) whose hyperbolic length is $\geq d_G(u, v)$.

Conversely, let γ be the hyperbolic geodesic between $z(u)$ and $z(v)$, and consider the set W of vertices w of G such that either $P(w)$ intersects γ or $P^\dagger(f)$ intersects γ for some face f incident to w . Let d be the length of γ . Since all circles in (P, P^\dagger) have a uniform upper bound on their hyperbolic radii,

we deduce that all circles in W are contained in a hyperbolic neighbourhood of constant thickness about γ , and hence the total area of these circles is $\leq d$. Since the radii of the circles are also bounded away from zero, we deduce that the cardinality of W is also $\leq d$. Since W contains a path in G from u to v , we deduce that $d_G(u, v)$ is $\leq d$ as required. \square

We summarize the situation for uniformly transient graphs in the following corollary, which follows immediately by combining Corollary 4.17 with Theorems 1.3–1.5 and Corollary 1.6. We write diam_G for the graph distance diameter of a set of vertices in G .

COROLLARY 4.18 (Graph distance exponents). *Let G be a uniformly transient, polyhedral, proper plane network with bounded codegrees and bounded local geometry, and let $\mathbf{p} > 0$ be a uniform lower bound on the escape probabilities of G . Then there exist positive constants $k_1 = k_1(\mathbf{M}, \mathbf{p})$ and $k_2 = k_2(\mathbf{M}, \mathbf{p})$ such that*

$$\begin{aligned} k_1 R^{-1} &\leq \text{FUSF}_G(\text{diam}_G(\Gamma_{\mathfrak{F}}(x, y)) \geq R) \leq k_2 R^{-1}, \\ k_1 R^{-1} &\leq \text{WUSF}_G(\text{diam}_G(\text{past}_{\mathfrak{F}}(e)) \geq R) \leq k_2 R^{-1}, \\ k_1 R^{-1/2} &\leq \text{WUSF}_G(|\text{past}_{\mathfrak{F}}(e)| \geq R) \leq k_2 R^{-1/2}, \\ k_1 R^{-1/2} &\leq \text{FUSF}_G(|\Gamma_{\mathfrak{F}}(x, y)| \geq R) \leq k_2 R^{-1/2} \end{aligned}$$

for every edge $e = (x, y)$ of G and every $R \geq 1$, where \mathfrak{F} is a sample of either the free or wired uniform spanning forest of G as appropriate.

5. Two remarks and a problem

5.1. Two remarks.

REMARK 5.1. Theorem 1.1 fails without the assumption that the degrees and edge conductances are bounded. Indeed, take two vertex-disjoint infinite simple paths in \mathbb{Z}^2 and put conductance 2^{2k} on the k th edge in each of these paths. It is not too hard to show using Wilson's algorithm that the FUSF of this graph has 2 connected components. A similar affect can be achieved with a simple unweighted graph by replacing the edges with large numbers of paths of length 2.

One may also wonder if this kind of pathological behaviour cannot occur when the underlying graph G is itself random. This is indeed the case, and in fact much more is true. It was proved by the first author [28] that a.s. every component of the WUSF is one-ended on any transient unimodular random rooted graph (with no planarity assumption). Using USF duality, one may deduce that the FUSF is a.s. connected on any simply connected unimodular random map of finite expected

degree whose dual is locally finite. Details are given in [6, Theorem 5.13], where this result is also extended to the case that the dual is not locally finite.

REMARK 5.2 (Nonuniversality in the parabolic case). Unlike in the CP hyperbolic case, the exponents governing the behaviour of the USTs of CP parabolic, polyhedral proper plane graphs with bounded codegrees and bounded local geometry are not universal, and need not exist in general.

Indeed, consider the double circle packing of the proper plane quadrangulation with underlying graph $\mathbb{N} \times \mathbb{Z}_4$, pictured in Figure 6. Let the packing be normalized to be symmetric under rotation by $\pi/2$ about the origin and to have $r(0, 0) = 1$. It is possible to compute that $r(i, j) = (3 + 2\sqrt{2})^i$ and hence that $|z(i, j)|$ is comparable to $(3 + 2\sqrt{2})^i$ for every $(i, j) \in \mathbb{N} \times \mathbb{Z}_4$. Suppose that the edges connecting (i, j) to $(i \pm 1, j)$ are given weight 1 for every $(i, j) \in \mathbb{N} \times \mathbb{Z}_4$, while the edges connecting (i, j) to $(i, j \pm 1)$ are given weight c for each $(i, j) \in \mathbb{N} \times \mathbb{Z}_4$. It can be computed that the probability that a walk started at $(i, 0)$ hits $(0, 0)$ without ever changing its second coordinate is $a(c)^i := (1 + c - \sqrt{c^2 + 2c})^i$. Let $e = ((0, 0), (0, 1))$. By running Wilson's algorithm rooted at $(0, 0)$ starting from the vertices $(i, 0)$ and $(i, 1)$, we see that

$$\begin{aligned} \text{UST}(\text{past}_T(e) \cap \{i\} \times \mathbb{Z}_4 \neq \emptyset) &\geq \mathbf{P}_{(i,0)}(\tau_{(0,0)} < \tau_{\mathbb{N} \times \{1,2,3\}}) \mathbf{P}_{(i,1)}(\tau_{(0,1)} < \tau_{\mathbb{N} \times \{0,2,3\}}) \\ &\quad \cdot \mathbf{P}_{(0,1)}(X_1 = (0, 0)) \\ &= \frac{c}{2c + 1} a(c)^{2i}. \end{aligned}$$

The right-hand side is exactly the probability that the random walk from $(i, 0)$ hits $(0, 0)$ without ever changing its second coordinate, and that the random walk from $(i, 1)$ hits $(0, 1)$ without ever changing its second coordinate and then steps to $(0, 0)$.

Let $q(c)$ be the probability that a random walk started at (i, j) visits every vertex of $\{i\} \times \mathbb{Z}_4$ before changing its vertical coordinate, which tends to one as $c \rightarrow \infty$. Let Y be a loop-erased random walk from $(0, 0)$. It can be computed that the probability that a random walk started from (i, j) visits $\{0\} \times \mathbb{Z}_4$ before hitting the trace of Y is at most $b(c)^i := (1 - q(c))^i$. Thus, by Wilson's algorithm and a union bound,

$$\text{UST}(\text{past}_T(e) \cap \{i\} \times \mathbb{Z}_4 \neq \emptyset) \leq 4b(c)^i.$$

It follows that there exist positive constants $k(c)$, $\alpha(c)$ and $\beta(c)$ such that $\alpha(c) \rightarrow 0$ as $c \rightarrow 0$, $\beta(c) \rightarrow \infty$ as $c \rightarrow \infty$, and

$$k(c)^{-1} R^{-\alpha(c)} \leq \text{UST}(\text{diam}_{\mathbb{C}}(\text{past}_T(e)) \geq R) \leq k(c) R^{-\beta(c)}.$$

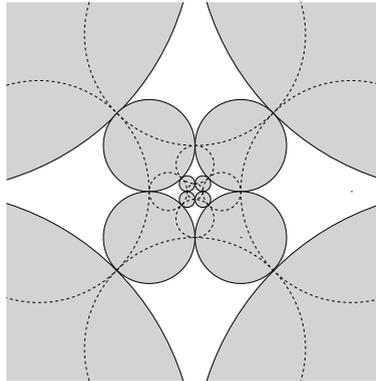


Figure 6. The double circle packing of $\mathbb{N} \times \mathbb{Z}_4$.

Thus, by varying c , we obtain CP parabolic proper plane networks with bounded codegrees and bounded local geometry with different exponents governing the diameter of the pasts of edges in their USTs: If c is large the diameter has a light tail, while if c is small the diameter has a heavy tail. Furthermore, by varying the weight of $((i, j), (i, j \pm 1))$ as a function of i in the above example (that is, making c small at some scales and large at others), it is possible to construct a polyhedral, CP parabolic proper plane network G with bounded codegrees and bounded local geometry such that

$$\frac{\log \text{UST}_G(\text{diam}_{\mathbb{C}}(\text{past}_T(e)) \geq R)}{\log(R)}$$

does not converge as $R \rightarrow \infty$ for some edge e of G . The details are left to the reader.

Similar constructions show that the behaviour of $\text{WUSF}_G(\text{diam}_{\mathbb{H}}(\text{past}_{\mathbb{F}}(e)) \geq R \cdot r_{\mathbb{H}}(x))$ is not universal over polyhedral, CP hyperbolic proper plane network G with bounded codegrees and bounded local geometry in the regime that $r_{\mathbb{H}}(x)$ is small.

5.2. A problem. It is natural to ask to what extent the assumption of planarity in Theorem 1.1 can be relaxed. Part (1) of the following question was suggested by R. Lyons.

QUESTION 5.3. Let G be a bounded degree proper plane graph.

- (1) Let H be a finite graph. Is the free uniform spanning forest of the product graph $G \times H$ connected almost surely?
- (2) Let G' be a bounded degree graph that is rough-isometric to G . Is the free uniform spanning forest of G connected almost surely?

Without the assumption of planarity, connectivity of the FUSF is not preserved by rough isometries; this can be seen from an analysis of the graphs appearing in [11, Theorem 3.5].

Acknowledgements

TH thanks Tel Aviv University and both authors thank the Issac Newton Institute, where part of this work was carried out, for their hospitality. We also thank Tyler Helmuth for finding several typos in an earlier version of this manuscript. Images of circle packings were created with Ken Stephenson's CirclePack software [51]. This project was supported by ERC starting grant RADNGEOM 676970.

References

- [1] D. Aharonov, 'The sharp constant in the ring lemma', *Complex Var. Theory Appl.* **33**(1–4) (1997), 27–31.
- [2] D. Aldous and R. Lyons, 'Processes on unimodular random networks', *Electron. J. Probab.* **12**(54) (2007), 1454–1508.
- [3] J. W. Anderson, *Hyperbolic Geometry*, 2nd edn, Springer Undergraduate Mathematics Series (Springer London, Ltd., London, 2005), xii+276.
- [4] O. Angel, M. T. Barlow, O. Gurel-Gurevich and A. Nachmias, 'Boundaries of planar graphs, via circle packings', *Ann. Probab.* **44**(3) (2016), 1956–1984.
- [5] O. Angel, T. Hutchcroft, A. Nachmias and G. Ray, 'Unimodular hyperbolic triangulations: circle packing and random walk', *Invent. Math.* **206**(1) (2016), 229–268.
- [6] O. Angel, T. Hutchcroft, A. Nachmias and G. Ray, 'Hyperbolic and parabolic unimodular random maps', *Geom. Funct. Anal.* **28**(4) (2018), 879–942.
- [7] O. Angel and G. Ray, 'The half plane UIPT is recurrent', *Probab. Theory Related Fields* **170**(3–4) (2018), 657–683.
- [8] M. T. Barlow, 'Loop erased walks and uniform spanning trees', in *Discrete Geometric Analysis*, MSJ Mem., 34 (Math. Soc. Japan, Tokyo, 2016), 1–32.
- [9] M. T. Barlow and R. Masson, 'Spectral dimension and random walks on the two dimensional uniform spanning tree', *Commun. Math. Phys.* **305**(1) (2011), 23–57.
- [10] I. Benjamini, R. Lyons, Y. Peres and O. Schramm, 'Uniform spanning forests', *Ann. Probab.* **29**(1) (2001), 1–65.
- [11] I. Benjamini and O. Schramm, 'Harmonic functions on planar and almost planar graphs and manifolds, via circle packings', *Invent. Math.* **126**(3) (1996), 565–587.

- [12] I. Benjamini and O. Schramm, ‘Recurrence of distributional limits of finite planar graphs’, *Electron. J. Probab.* **6**(23) (2001), 1–13.
- [13] S. Bhupatiraju, J. Hanson and A. A. Járai, ‘Inequalities for critical exponents in d -dimensional sandpiles’, *Electron. J. Probab.* **22**(85) (2017), 51.
- [14] G. R. Brightwell and E. R. Scheinerman, ‘Representations of planar graphs’, *SIAM J. Discrete Math.* **6**(2) (1993), 214–229.
- [15] R. Burton and R. Pemantle, ‘Local characteristics, entropy and limit theorems for spanning trees and domino tilings via transfer-impedances’, *Ann. Probab.* **21**(3) (1993), 1329–1371.
- [16] D. Chelkak *et al.*, ‘Robust discrete complex analysis: A toolbox’, *Ann. Probab.* **44**(1) (2016), 628–683.
- [17] C. Garban, ‘Quantum gravity and the KPZ formula [after Duplantier-Sheffield]’, *Astérisque* **352** (2013), 315–354.
- [18] O. Gurel-Gurevich and A. Nachmias, ‘Recurrence of planar graph limits’, *Ann. of Math. (2)* **177**(2) (2013), 761–781.
- [19] O. Gurel-Gurevich, A. Nachmias and J. Souto, ‘Recurrence of multiply-ended planar triangulations’, *Electron. Commun. Probab.* **22**(5) (2017), 6.
- [20] O. Häggström, ‘Random-cluster measures and uniform spanning trees’, *Stochastic Process. Appl.* **59**(2) (1995), 267–275.
- [21] L. J. Hansen, ‘On the Rodin and Sullivan ring lemma’, *Complex Var. Theory Appl.* **10**(1) (1988), 23–30.
- [22] Z.-X. He, ‘Rigidity of infinite disk patterns’, *Ann. of Math. (2)* **149** (1999), 1–33.
- [23] Z.-X. He and O. Schramm, ‘Fixed points, Koebe uniformization and circle packings’, *Ann. of Math. (2)* **137**(2) (1993), 369–406.
- [24] Z.-X. He and O. Schramm, ‘Hyperbolic and parabolic packings’, *Discrete Comput. Geom.* **14**(2) (1995), 123–149.
- [25] Z.-X. He and O. Schramm, ‘The C^∞ -convergence of hexagonal disk packings to the Riemann map’, *Acta Math.* **180**(2) (1998), 219–245.
- [26] T. Hutchcroft, ‘Wired cycle-breaking dynamics for uniform spanning forests’, *Ann. Probab.* **44**(6) (2016), 3879–3892.
- [27] T. Hutchcroft, ‘Harmonic Dirichlet functions on planar graphs’, *Discrete Comput. Geom.* **61**(3) 479–506.
- [28] T. Hutchcroft, ‘Interlacements and the wired uniform spanning forest’, *Ann. Probab.* **46**(2) (2018), 1170–1200.
- [29] T. Hutchcroft and Y. Peres, ‘Boundaries of planar graphs: a unified approach’, *Electron. J. Probab.* **22**(100) (2017), 20.
- [30] J. Jonasson and O. Schramm, ‘On the cover time of planar graphs’, *Electron. Comm. Probab.* **5** (2000), 85–90 (electronic).
- [31] R. Kenyon, ‘The asymptotic determinant of the discrete Laplacian’, *Acta Math.* **185**(2) (2000), 239–286.
- [32] G. Kirchhoff, ‘Ueber die auflösung der gleichungen, auf welche man bei der untersuchung der linearen vertheilung galvanischer ströme geführt wird’, *Ann. Phys.* **148**(12) (1847), 497–508.
- [33] P. Koebe, *Kontaktprobleme der konformen Abbildung*, (Hirzel, 1936).
- [34] G. F. Lawler, *Conformally Invariant Processes in the Plane*, Mathematical Surveys and Monographs, 114 (American Mathematical Society, Providence, RI, 2005), xii+242.
- [35] G. F. Lawler *et al.*, ‘A self-avoiding random walk’, *Duke Math. J.* **47**(3) (1980), 655–693.
- [36] G. F. Lawler, O. Schramm and W. Werner, ‘Conformal invariance of planar loop-erased random walks and uniform spanning trees’, *Ann. Probab.* **32**(1) (2004), 939–995.

- [37] R. Lyons, B. J. Morris and O. Schramm, 'Ends in uniform spanning forests', *Electron. J. Probab.* **13**(58) (2008), 1702–1725.
- [38] R. Lyons and Y. Peres, *Probability on Trees and Networks*, Cambridge Series in Statistical and Probabilistic Mathematics, 42 (Cambridge University Press, New York, 2016).
- [39] A. Marden and B. Rodin, 'On Thurston's formulation and proof of Andreev's theorem', in *Computational Methods and Function Theory (Valparaíso, 1989)*, Lecture Notes in Math., 1435 (Springer, Berlin, 1990), 103–115.
- [40] R. Masson, 'The growth exponent for planar loop-erased random walk', *Electron. J. Probab.* **14**(36) (2009), 1012–1073.
- [41] B. Morris, 'The components of the wired spanning forest are recurrent', *Probab. Theory Related Fields* **125**(2) (2003), 259–265.
- [42] M. Murugan, 'Quasisymmetric uniformization and heat kernel estimates', Preprint, 2018, [arXiv:1803.11296](https://arxiv.org/abs/1803.11296).
- [43] R. Pemantle, 'Choosing a spanning tree for the integer lattice uniformly', *Ann. Probab.* **19**(4) (1991), 1559–1574.
- [44] Y. Peres, 'Probability on trees: an introductory climb', in *Lectures on Probability Theory and Statistics (Saint-Flour, 1997)*, Lecture Notes in Math., 1717 (Springer, Berlin, 1999), 193–280.
- [45] B. Rodin and D. Sullivan, 'The convergence of circle packings to the Riemann mapping', *J. Differential Geom.* **26**(2) (1987), 349–360.
- [46] S. Rohde, 'Oded Schramm: from circle packing to SLE', *Ann. Probab.* **39** (2011), 1621–1667.
- [47] O. Schramm, 'Rigidity of infinite (circle) packings', *J. Amer. Math. Soc.* **4**(1) (1991), 127–149.
- [48] O. Schramm, 'Scaling limits of loop-erased random walks and uniform spanning trees', *Israel J. Math.* **118** (2000), 221–288.
- [49] D. Shiraishi, 'Growth exponent for loop-erased random walk in three dimensions', *Ann. Probab.* **46**(2) (2018), 687–774.
- [50] R. Siders, 'Layered circlepackings and the type problem', *Proc. Amer. Math. Soc.* **126**(10) (1998), 3071–3074.
- [51] K. Stephenson, 'Circle pack, Java 2.0', <http://www.math.utk.edu/~kens/CirclePack>.
- [52] K. Stephenson, *Introduction to Circle Packing: The Theory of Discrete Analytic Functions*, (Cambridge University Press, Cambridge, 2005).
- [53] W. Thurston, 'Hyperbolic geometry and 3-manifolds', in *Low-dimensional Topology (Bangor, 1979)*, London Math. Soc. Lecture Note Ser., 48 (Cambridge Univ. Press, Cambridge-New York, 1982), 9–25.
- [54] D. B. Wilson, 'Generating random spanning trees more quickly than the cover time', in *Proceedings of the Twenty-eighth Annual ACM Symposium on the Theory of Computing (Philadelphia, PA, 1996)* (ACM, New York, 1996), 296–303.