Canad. J. Math. Vol. 70 (4), 2018 pp. 925-942 http://dx.doi.org/10.4153/CJM-2017-028-0 © Canadian Mathematical Society 2017



# Edge-Maximal Graphs on Surfaces

## Colin McDiarmid and David R. Wood

Abstract. We prove that for every surface  $\Sigma$  of Euler genus g, every edge-maximal embedding of a graph in  $\Sigma$  is at most O(g) edges short of a triangulation of  $\Sigma$ . This provides the first answer to an open problem of Kainen (1974).

#### 1 Introduction

For a graph class  $\mathcal{G}$ , a graph  $G \in \mathcal{G}$  is *edge-maximal* if adding any non-edge to G produces a graph not in  $\mathcal{G}$ . We emphasise that "graph" here means a simple graph with no parallel edges and no loops. A graph class  $\mathcal{G}$  is *pure* if |E(G)| = |E(H)| for all edge-maximal graphs  $G, H \in \mathcal{G}$  with |V(G)| = |V(H)|. For example, each of the following graph classes is pure: forests, outerplanar graphs, planar graphs; and for each positive integer k, the k-degenerate graphs, the graphs of treewidth at most k, and the chordal graphs with clique number at most k + 1 (where the last two classes have the same edge-maximal members, the k-trees). On the other hand, toroidal graphs are not pure: [2] proved that  $K_8 - E(C_5)$  is an edge-maximal toroidal graph but is not a toroidal triangulation (see Figure 1).

Motivated by this example, Kainen [4] posed the following open problem: by how many edges can an edge-maximal graph embeddable in a given surface fail to be a triangulation? This paper addresses this natural question, which surprisingly has been ignored in the literature. We prove that for every surface  $\Sigma$  of Euler genus g, every edge-maximal graph embeddable in  $\Sigma$  is O(g) edges short of a triangulation (regardless of the number of vertices).

We formulate this result as follows. A graph class  $\mathcal{G}$  is *k*-impure if  $||E(G)| - |E(H)|| \le k$  for all edge-maximal graphs  $G, H \in \mathcal{G}$  with |V(G)| = |V(H)|. Thus,  $\mathcal{G}$  is 0-impure means that it is pure. We introduce this notation as a measure of how far  $\mathcal{G}$  is from being pure. For  $h \ge 0$ , let  $\mathbb{S}_h$  be the sphere with *h* handles. For  $c \ge 0$ , let  $\mathbb{N}_c$  be the sphere with *c* cross-caps. Every surface is homeomorphic to  $\mathbb{S}_h$  or  $\mathbb{N}_c$ . The *Euler genus* of  $\mathbb{S}_h$  is 2*h*. The *Euler genus* of  $\mathbb{N}_c$  is *c*. The *Euler genus* of a graph *G* is the minimum Euler genus of a surface in which *G* embeds. See [8] for definitions and background regarding graphs embedded in surfaces. The following is our main theorem; see Theorems 3.3 and 3.4 for fuller forms of this result.

**Theorem 1.1** The class of graphs embeddable in a surface  $\Sigma$  of Euler genus g is O(g)-impure.

Received by the editors May 9, 2017.

Published electronically October 2, 2017.

Research of Wood is supported by the Australian Research Council.

AMS subject classification: 05C10.

Keywords: graph, surface, embedding.



*Figure 1*: An embedding of  $K_8 - E(C_5)$  in the torus. Every such embedding has one 4-face, which induces  $K_4$ , so no non-edge can be added.

To add some perspective to this result, note that several interesting graph classes are not at all pure. Consider, for example, the  $K_5$ -minor-free graphs. The 8-vertex Mobius ladder W is  $K_5$ -minor-free with 12 edges. Pasting copies of W on edges (in an arbitrary way) produces a  $K_5$ -minor-free graph with  $n \equiv 2 \pmod{6}$  vertices and (11n - 16)/6 edges (see Figure 2). It is edge-maximal with no  $K_5$ -minor by Wagner's characterisation [16]. On the other hand, every *n*-vertex edge-maximal planar graph is edge-maximal with no  $K_5$ -minor, yet has 3(n - 2) edges for  $n \ge 3$ . Thus, the difference between the numbers of edges in these two classes of edge-maximal  $K_5$ -minorfree graphs grows with n, and indeed is  $\Omega(n)$ . In general,  $K_t$ -minor-free graphs can have as many as  $ct\sqrt{\log t} n$  edges [5,13,14], but there are edge-maximal  $K_t$ -minor-free graphs, namely (t - 2)-trees, with only  $(t - 2)n - {t-1 \choose 2}$  edges (for  $n \ge t - 1$ ).



Figure 2: Pasting copies of W.

Let  $\mathcal{G}_H$  denote the class of graphs not containing H as a minor. McDiarmid and Przykucki [6] proved that (ignoring  $K_1$ ) the only connected graphs H such that  $\mathcal{G}_H$ is pure are  $K_2$ ,  $K_3$ ,  $K_4$ , and  $P_3$  (the 3-vertex path). Furthermore, for each connected graph *H*, either  $\mathcal{G}_H$  is *k*-impure for some *k*, or there are *n*-vertex graphs  $G_n$  and  $G'_n$  in  $\mathcal{G}_H$  such that  $|E(G_n)| - |E(G'_n)|$  is  $\Omega(n)$ .

# 2 Main Proof

An embedding of a graph *G* in a surface is *edge-maximal* if for every non-edge *e* of *G*, it is not possible to add *e* to the embedding (without changing the embedding of *G* or creating parallel edges or loops). Observe that an embedding of a graph *G* in a surface is edge-maximal if and only if for each face *F*, the set of vertices on *F* induce a clique in *G*. Also note that a graph *G* is edge-maximal embeddable in a surface  $\Sigma$  if and only if every embedding of *G* in  $\Sigma$  is edge-maximal. We mentioned above that Theorems 3.3 and 3.4 give fuller forms of Theorem 1.1; in fact, they concern edge-maximal embeddings (as well as giving explicit constants). The distinction between edge-maximal embeddings and edge-maximal embeddable graphs is exemplified by the following fact. An embedding is *cellular* (or *2-cell*) if each face is homeomorphic to an open disc.

**Proposition 2.1** For each surface  $\Sigma$ , there are infinitely many planar graphs, each with an edge-maximal cellular embedding in  $\Sigma$ .

**Proof** First suppose that  $\Sigma = \mathbb{N}_g$ . Let  $G_0$  be a triangulation of the sphere with at least g faces. Say  $F_1, \ldots, F_g$  are distinct faces of  $G_0$ . Note that  $K_4$  has a cellular embedding in the projective plane with two triangular faces and one face of length 6 (see Figure 3). Let  $Q_1, \ldots, Q_g$  be g copies of this embedding of  $K_4$ . For  $i \in [1, g]$ , identify  $F_i$  with a triangular face of  $Q_i$ . We obtain a graph G embedded in  $\mathbb{N}_g$ , in which each face induces a clique. Thus, this embedding of G is edge-maximal. Note that G is a planar triangulation, since it is obtained from  $G_0$  by simply adding a degree-3 vertex inside g faces of  $G_0$ . An analogous proof works for  $\Sigma = \mathbb{S}_h$ , since  $K_4$  has a cellular embedding in the torus with one triangular face and one face of length 9 (see Figure 3).



*Figure 3*: Embeddings of  $K_4$  in  $\mathbb{N}_1$  and  $\mathbb{S}_1$ .

A *pseudograph* is a graph possibly with parallel edges and loops. A (*pseudograph*) *triangulation* is a cellular embedded (pseudo)graph in which each face has length exactly 3. Euler's formula implies that every pseudograph with  $n \ge 3$  vertices that embeds in a surface of Euler genus g such that each face has length at least 3 has at most 3(n + g - 2) edges, with equality if and only if the embedding is a pseudograph triangulation. Of course, every face in an embedding of a graph (with at least three vertices) has length at least 3. Thus, every graph with  $n \ge 3$  vertices that embeds in a surface of Euler genus g has at most 3(n + g - 2) edges, with equality if and only if the embedding is a pseudograph triangulation. Of course, every face in an embedding of a graph (with at least three vertices) has length at least 3. Thus, every graph with  $n \ge 3$  vertices that embeds in a surface of Euler genus g has at most 3(n + g - 2) edges, with equality if and only if the embedding is a triangulation. Also note that Euler's formula implies that every bipartite graph with  $n \ge 3$  vertices that embeds in a surface of Euler genus g has at most 2(n + g - 2) edges. We implicitly use these facts throughout the paper.

Given an embedding of an *n*-vertex graph in a surface  $\Sigma$  of Euler genus *g* (where  $n+g \ge 3$ ), we can add edges (if necessary) to obtain a pseudograph triangulation with exactly 3(n+g-2) edges. Moreover, note as an aside that if  $n \ge 3$ , then we do not need loops, as shown in Proposition 4.1. When we say that an edge-maximal embedding in  $\Sigma$  or an edge-maximal graph embeddable in  $\Sigma$  is "*k* edges short of a triangulation", we mean that it has exactly 3(n+g-2) - k edges.

We need the following lemmas about edge-maximal embeddings. The first says that we can restrict our attention to cellular embeddings.

**Lemma 2.2** Let  $c \ge 3$ , and assume that for every g and every surface  $\Sigma$  of Euler genus g, every edge-maximal cellular embedding in  $\Sigma$  is at most cg edges short of a triangulation of  $\Sigma$ . Then for every g and every surface  $\Sigma$  of Euler genus g, every edge-maximal embedding in  $\Sigma$  is at most cg edges short of a triangulation of  $\Sigma$ .

**Proof** Consider an edge-maximal embedding of a graph *G* in some surface  $\Sigma$  of Euler genus *g*. This embedding defines a combinatorial embedding of *G*, which corresponds to a cellular embedding in some surface  $\Sigma'$  of Euler genus  $g' \leq g$ . If a non-edge of *G* can be added to this embedding in  $\Sigma'$ , then the same non-edge can be added to the original embedding in  $\Sigma$ . Since the embedding in  $\Sigma$  is edge-maximal, so too is the embedding in  $\Sigma'$ . By assumption, *G* is at most cg' edges short of a triangulation in  $\Sigma'$ . That is,

$$|E(G)| \ge 3(|V(G)| + g' - 2) - cg' = 3(|V(G)| - 2) - (c - 3)g'$$
  
$$\ge 3(|V(G)| - 2) - (c - 3)g = 3(|V(G)| + g - 2) - cg.$$

That is, *G* is at most *cg* edges short of a triangulation in  $\Sigma$ .

For a vertex v in a graph G, let  $N(v) := \{w \in V(G) : vw \in E(G)\}$  be the open neighbourhood of v and let  $N[v] := N(v) \cup \{v\}$  be the closed neighbourhood of v.

*Lemma 2.3* Every graph G with  $n \ge 4$  vertices that has an edge-maximal cellular embedding in some surface is 3-connected.

**Proof** The graph *G* is connected, since the embedding is edge-maximal and Euler genus is additive on components and blocks [8]. If *G* contains a vertex *v* of degree 1 and *vw* is the edge incident to *v*, then *w* has a distinct neighbour, so the facial walk starting with *vw* is followed by *wx* for some  $x \notin \{v, w\}$ , and the edge *vx* can be added

to *G*, contradicting the edge-maximality of the embedding of *G*. Thus, *G* has minimum degree at least 2. Let  $\pi_v$  denote the cyclic ordering of edges incident to each vertex v in an embedding of *G* in  $\Sigma$ .

Suppose *G* contains a vertex *v* of degree 2. Let *u* and *w* be the neighbours of *v*. We can assume that the edges *uv* and *vw* have signature +1. For clarity, observe that the edge *uw* must be in *G*, since if not we could add it. Since *G* is connected and  $n \ge 4$ , at least one of *u* and *w*, say *w*, has a neighbour not in  $\{u, v, w\}$ . Consider the cyclic order  $\pi_w$ : if *wu* follows *wv*, then let *wx* be the edge preceding *wv*, otherwise, let *wx* be the edge following *wv*. Note that *x* is not in  $\{u, v, w\}$ . We can add the edge *vx*, with signature +1, as follows. Insert *vx* in  $\pi_v$  after *vw* and insert *xv* in  $\pi_x$  before *xw*. The original facial walk *W* starting *xwvu*... is replaced by two facial walks  $W_1 = xwvx$  and  $W_2 = xvu$ ..., where  $W_2$  is obtained from *W* by replacing the two-edge path *xwv* by the single edge *xv*. By maximality, *G* has minimum degree at least 3.

We now prove that *G* is *locally Hamiltonian*; that is, for each vertex *v* the subgraph induced on N(v) has a Hamilton cycle. Without loss of generality, the edges incident to *v* have signature +1. Let  $(vv_1, vv_2, ..., vv_d)$  be the cyclic ordering of the edges incident to *v*, where  $d \ge 3$ . We claim that  $v_1v_2...v_d$  is a cycle. For suppose that say  $v_1$ and  $v_2$  are not adjacent. If *F* is the face with facial walk starting  $(v_1v, vv_2, ...)$ , then we can add the edge  $v_1v_2$  across *F*, which is a contradiction. Thus *G* is locally Hamiltonian. Finally, every connected locally Hamiltonian graph is 3-connected. This was shown in [15, Proposition 7.1], but we give a proof here for completeness since it is short. Clearly *G* cannot have a separating vertex. Suppose *G* has a separating pair of vertices u, v. Thus,  $V(G) \setminus \{u, v\}$  can be partitioned into two non-empty parts *U* and *W* such that there are no U-W edges. Then *v* must have a neighbour  $a \in U$  and  $b \in W$  (otherwise *u* is a separating vertex), and there are two internally disjoint *ab*paths in G - v (around a Hamilton cycle in N(v)). But both paths must go through u, a contradiction. Hence *G* is 3-connected.

**Lemma 2.4** Let G be a graph with at least four vertices that has an edge-maximal cellular embedding in a surface. Then every non-triangular face contains four distinct vertices that are consecutive on the facial walk. Furthermore, for each string of six vertices that are consecutive on the facial walk, at least one of the three substrings of length 4 consists of distinct vertices.

**Proof** If *a*, *b*, *c* are consecutive vertices on a face *F*, then *a*, *b*, *c* are distinct, as otherwise deg(*b*) = 1, which would contradict Lemma 2.3. Thus, if *F* has length 4 or 5, then all the vertices on *F* are distinct, and we are done. Now assume that *F* has length at least 6. Let  $v_1, \ldots, v_6$  be consecutive vertices on *F*. If  $v_1 = v_4$  and  $v_2 = v_5$  and  $v_3 = v_6$ , then the sequence is  $v_1, v_2, v_3, v_1, v_2, v_3$ , and the graph is  $K_3$  (embedded in  $\mathbb{N}_1$ ). Without loss of generality,  $v_1 \neq v_4$ , implying that  $v_1, v_2, v_3, v_4$  are distinct.

We noted earlier that Euler genus is additive on components and blocks. The main tool used in our proof is the following, more general, additivity theorem, proved independently by several authors. **Theorem 2.5** ([1,7,10]) If graphs  $G_1$  and  $G_2$  have at most two vertices in common, then the Euler genus of  $G_1 \cup G_2$  is at least the Euler genus of  $G_1$  plus the Euler genus of  $G_2$ .

Say a sequence  $v_1, \ldots, v_s$  of vertices in a graph *G* is *ordered* if for each  $i \in [2, s]$ ,

$$\left|N[v_i] \cap \left(\bigcup_{j=1}^{i-1} N[v_j]\right)\right| \leq 2.$$

Theorem 2.5 implies the following result.

**Corollary 2.6** ([3,9]) If  $v_1, ..., v_s$  is an ordered sequence of vertices in a graph G, and each  $N[v_i]$  is a clique on at least five vertices, then the Euler genus of G is at least s.

We prove in (2.2) that given integers  $g \ge 0$  and  $s \ge 1$ , there is an integer b such that for every bipartite graph G with Euler genus at most g, if (A, B) is a bipartition of G such that |B| > b and every vertex in B has degree at most 4, then B contains an ordered sequence of s vertices. Let  $f_g(s)$  be the least such integer b. We now give some illustrative examples. Since one vertex forms an ordered sequence,  $f_g(1) = 0$  for each  $g \ge 0$ . The planar bipartite graph Q shown in Figure 4 has a colour class B with three vertices, each pair of which has three common neighbours. Thus, B contains no ordered sequence of length 2. Thus,  $f_0(2) \ge 3$ . It is easily seen that  $f_0(2) \le 3$ (using a straightforward adaptation of the proof of Lemma 3.1). Thus,  $f_0(2) = 3$ . Now consider general  $g \ge 0$ . Ringel [11, 12] proved that the Euler genus of  $K_{3,2g+2}$  equals g. If *B* is the colour class of degree-3 vertices in  $K_{3,2g+2}$ , then each pair of vertices in *B* has three common neighbours. Thus, B contains no ordered sequence of length 2, and  $f_g(2) \ge 2g+2$ . Lemma 3.1 proves this inequality is tight for  $g \ge 1$ . These constructions can be combined as follows. Fix  $g \ge 0$  and  $s \ge 2$ . Let G be the graph obtained from  $K_{3,2g+2}$  by adding s - 2 disjoint copies of Q. Then G has Euler genus g, and G has a bipartition (A, B) where |B| = 2g+2+3(s-2), and every ordered sequence in B has at most one vertex from each of the s - 1 components of G. Thus B contains no ordered sequence of length s, and

$$(2.1) f_g(s) \ge 2g + 3s - 4.$$

The next lemma motivates the definition of  $f_g(s)$ .

**Lemma 2.7** Every edge-maximal embedding of a graph G in a surface  $\Sigma$  of Euler genus  $g \ge 1$  is at most  $5f_g(g+1) - 1$  edges short of a triangulation of  $\Sigma$ .

**Proof** Note that  $f_g(g+1) \ge 5g-1$  by (2.1), which implies that  $5f_g(g+1)-1 \ge 3g$ . Thus, we can assume this embedding is cellular by Lemma 2.2. Let n := |V(G)|. If  $n \le 7g$ , then the number of edges in a triangulation, 3(n+g-2), is at most  $24g-6 < 5(5g-1)-1 \le 5f_g(g+1)-1$  by (2.1), and the result holds. Now assume that  $n \ge 7g+1 \ge 8$ .

By Lemma 2.3, *G* has minimum degree at least 3. We can assume the embedding of *G* is not a triangulation. Let *G'* be the embedded pseudograph obtained from *G* as follows. Consider in turn each face *F* in *G* with length  $t \ge 4$ . We shall add edges to *G* across *F* so that each of the resulting faces in *G'* still contains at least four distinct vertices. By Lemma 2.4, *F* contains four distinct consecutive vertices. Let



*Figure 4*: The graph *Q*.

 $(v_0, v_1, v_2, ..., v_{t-1})$  be a facial walk of *F*, where  $v_0, v_1, v_2, v_3$  are distinct. Add the edge  $v_0v_i$  to *G* whenever  $i \equiv 3 \pmod{5}$  and  $3 \le i \le t-5$ ; this divides *F* into  $\lfloor \frac{t+2}{5} \rfloor$  faces in *G'*, each containing at least four distinct vertices (since  $v_0, v_1, v_2, v_3$  are distinct, and every other face contains six consecutive vertices in *F*, and thus, by Lemma 2.4, has at least four distinct vertices).

For each non-triangular face *F* of *G'*, add a vertex inside *F* adjacent to four distinct vertices of *F*. Let *B* be the set of these added vertices, and let *G''* be the resulting embedded graph. Since the embedding of *G* is edge-maximal, each face of *G* induces a clique. Thus,  $N_{G''}[v]$  induces  $K_5$  for each  $v \in B$ .

Consider a non-triangular face *F* of length *t* in *G*. Then *B* contains exactly  $\lfloor \frac{t+2}{5} \rfloor$  vertices corresponding to *F*. Note that  $t-3 \leq 5 \lfloor \frac{t+2}{5} \rfloor - 1$  edges are sufficient (and necessary) to triangulate *F*. Thus, the embedding of *G* can be extended to a triangulation by adding at most 5|B| - 1 edges.

Let G''' be the induced bipartite subgraph of G'' with bipartition  $\{B, \cup_{v \in B} N_{G''}(v)\}$ . By construction, G''' embeds in  $\Sigma$  and every vertex in B has degree 4.

Suppose for a contradiction that  $|B| > f_g(g + 1)$ . Thus, *B* contains an ordered sequence  $v_1, \ldots, v_{g+1}$  in *G''*. Since  $N_{G''}[v_i]$  induces  $K_5$ , by Corollary 2.6, the Euler genus of *G''* is at least g + 1, which is a contradiction. Thus,  $|B| \leq f_g(g + 1)$ . Hence *G* is at most  $5f_g(g + 1) - 1$  edges short of a triangulation.

It remains to show how to find ordered sequences. The next lemma is useful.

**Lemma 2.8** Fix an integer  $c \ge 7$ . Let G be a bipartite graph with bipartition A, B and with Euler genus at most g. If B is non-empty and  $|B| > \frac{2c}{c-6}(g-2)$ , then some vertex in B has at most two neighbours with degree at least c.

**Proof** Let *A'* be the set of vertices in *A* with degree at least *c*. Suppose for a contradiction that every vertex in *B* has at least three neighbours in *A'*. Double-counting the edges with endpoints in *A'* and *B* gives  $c|A'| \leq 2(|A'| + |B| + g - 2)$  and  $3|B| \leq 2(|A'| + |B| + g - 2)$ . Adding 2 times the first inequality plus c - 2 times the second inequality gives  $|B| \leq \frac{2c}{c-6}(g-2)$ , which is the desired contradiction.

We have the following recursive upper bound for  $f_g(s)$ .

*Lemma 2.9* For integers  $g \ge 1$  and  $s \ge 2$  and  $c \ge 7$ ,

$$f_g(s) \leq \max\left\{\frac{2c}{c-6}(g-2), 2c-3+f_g(s-1)\right\}.$$

**Proof** Let *G* be a bipartite graph with Euler genus at most *g*, where (*A*, *B*) is a bipartition of *G* such that  $|B| > \max \{ \frac{2c}{c-6}(g-2), 2c-3 + f_g(s-1) \}$  and every vertex in *B* has degree at most 4. Our goal is to show that *B* contains an ordered sequence of *s* vertices. Since *B* is non-empty and  $|B| > \frac{2c}{c-6}(g-2)$ , by Lemma 2.8, some vertex  $v_s$  in *B* has at most two neighbours with degree at least *c*. If deg( $v_s$ )  $\leq 2$  then let  $X := \emptyset$ . Otherwise, let *X* be the set of neighbours of  $v_s$  other than two of highest degree. Thus  $|X| \leq 2$  and each vertex  $u \in X$  has degree at most c - 1. Let *G'* be obtained from *G* by deleting N[u] for each  $u \in X$ . Let *A'*, *B'* be the bipartition of *G'* inherited from *G*. Note that

$$|B'| \ge |B| - (2c - 3) > \max\left\{\frac{2c}{c-6}(g - 2), 2c - 3 + f_g(s - 1)\right\} - (2c - 3) \ge f_g(s - 1).$$

Thus, B' contains an ordered sequence  $v_1, \ldots, v_{s-1}$  in G'. By construction,  $v_s$  has at most two neighbours in G'. Thus  $v_1, \ldots, v_s$  is an ordered sequence in G.

Since  $f_g(1) = 0$ , Lemma 2.9 implies that for all integers  $c \ge 7$  and  $s \ge 1$ ,

(2.2) 
$$f_g(s) \leq (2c-3)(s-2) + \max\left\{\frac{2c}{c-6}(g-2), 2c-3\right\}$$

With any choice of  $c \ge 7$ , this implies that  $f_g(g+1)$  is O(g). Lemma 2.7 then implies that every edge-maximal embedding in a surface  $\Sigma$  of Euler genus g is O(g) edges short of a triangulation of  $\Sigma$ . Therefore the graphs embeddable in  $\Sigma$  are O(g)-impure, which is the main result of this paper (Theorem 1.1). For example, with c = 8 and  $g \ge 4$ ,

$$f_g(g+1) \leq 13(g-1) + \max\left\{8(g-2), 13\right\} = 21g - 29,$$

and by Lemma 2.7 every edge-maximal graph embeddable in a surface of Euler genus  $g \ge 4$  is at most 105g - 146 edges short of a triangulation.

# 3 Improving the Constants

Our proof of Theorem 1.1 shows that the class of graphs embeddable in a given surface of Euler genus  $g \ge 4$  is (105g - 146)-impure. In this section we shall improve this bound; see Theorems 3.3 and 3.4. We first give a precise result for ordered sequences of length 2, improving on the bound in (2.2) with s = 2.

*Lemma* 3.1  $f_g(2) = 2g + 2$  for  $g \ge 1$ .

**Proof** We proved above that  $K_{3,2g+2}$  shows that  $f_g(2) \ge 2g + 2$  for  $g \ge 1$ . We now prove the corresponding upper bound.

Let *G* be a bipartite graph *G* with Euler genus at most *g*. Assume that (A, B) is a bipartition of *G* such that every vertex in *B* has degree at most 4 and  $|B| \ge 2g + 3$ . We claim that *B* contains an ordered sequence of two vertices. That is, *B* contains two vertices with at most two common neighbours. Suppose for a contradiction that each pair of vertices in *B* has at least three common neighbours.

By adding degree-1 vertices in *A*, we can assume that every vertex in *B* has degree exactly 4. Without loss of generality,  $A = \bigcup_{b \in B} N(b)$ . We have  $4|B| \le |E(G)| \le$ 2(|A| + |B| + g - 2) implying  $|B| \le |A| + g - 2$  and  $|A| \ge (2g + 3) - (g - 2) = g + 5 \ge 6$ . Let  $a, b \in B$  have  $N(a) \ne N(b)$ . Let  $X = N(a) \cap N(b)$  and  $Y = N(a) \cup N(b)$ . Then |X| = 3 and |Y| = 5. Let  $N(a) = X \cup \{a'\}$  and  $N(b) = X \cup \{b'\}$ . Since  $|A| \ge 6$ , there is a vertex  $c \in B$  with N(c) not contained in *Y*. If  $a' \in N(c)$ , then  $|N(c) \cap N(b)| \le 2$ , so  $a' \notin N(c)$ ; and similarly  $b' \notin N(c)$ . Hence  $N(c) \cap Y = X$ , so we can write N(c) = $X \cup \{c'\}$ . Note that a', b', c' are distinct and not in *X*, so we have symmetry between (a, a'), (b, b'), and (c, c').

Now for any  $v \in B$ , N(v) cannot contain  $\{a', b', c'\}$  (since then for example  $|N(v) \cap N(a)| \leq 2$ ); so assume without loss of generality that  $c' \notin N(v)$ . But then we must have  $N(v) \cap N(c) = X$ , and so N(v) contains X. We have shown that N(v) contains X for each  $v \in B$ . But now the induced bipartite graph with parts X and B is complete. Hence  $3|B| \leq 2(3 + |B| + g - 2)$ , implying  $|B| \leq 2g + 2 < 2g + 3$ . This contradiction completes the proof.

Lemmas 2.9 and 3.1 imply that for  $g \ge 1$  and  $s \ge 2$ ,

(3.1) 
$$f_g(s) \leq \begin{cases} 2g+2 & \text{if } s=2, \\ \min\left\{\max\left\{\frac{2c_s}{c_s-6}(g-2), 2c_s-3+f_g(s-1)\right\}: c_s \ge 7 \right\} & \text{if } s \ge 3. \end{cases}$$

For non-orientable surfaces, Table 1 shows the optimal choice of  $c_3, \ldots, c_{g+1}$  in (3.1) for each value of  $g \leq 20$ , along with the corresponding lower bound on the number of edges in an edge-maximal graph.

The next lemma shows a method for choosing the constants  $c_s$  in (3.1). All logarithms are natural.

*Lemma 3.2* Let  $\lambda = 25 - 11(\frac{48332}{114345} + \frac{16}{33}\log 2) \approx 16.6533\cdots$  to four decimal places. Then for  $g \ge 2$ ,

$$f_g(g+1) \leq \lambda(g-2) + 2\left[\sqrt{\frac{3}{2}(g-2)}\right] + 33.$$

**Proof** For  $i \ge 7$ , let

$$\alpha_i := \sum_{j=i+1}^{\infty} \frac{12}{(j-7)(j-6)(2j-3)}$$

Then

$$0.758757\cdots = \alpha_7 > \alpha_8 > \alpha_9 > \cdots$$

These numbers  $\alpha_i$  are used below to calculate the values  $c_s$  in (3.1). For example,  $\alpha_7 \approx 0.76$  means that  $c_s = 7$  roughly for  $0.76g \le s \le g$ , and  $\alpha_8 \approx 0.30$  means that  $c_s = 8$ roughly for  $0.30g \le s \le 0.76g$ . This behaviour is evident in the lower rows of Table 1. The definition of  $\alpha_i$  is designed to minimise the "max" operation in (3.1).

We now upper bound  $\alpha_k$ . Since  $(j-6)(2j-3) \ge 2(j-7)^2$  for  $j \ge 7$ ,

$$\alpha_k = \sum_{j=k+1}^{\infty} \frac{12}{(j-7)(j-6)(2j-3)} \leq \sum_{j=k+1}^{\infty} \frac{6}{(j-7)^3} \leq \int_{k+1}^{\infty} \frac{6}{(j-8)^3} dj = \frac{3}{(k-7)^2}.$$

g	surface	$c_3,\ldots,c_{g+1}$	impurity ≼	$ E(G)  \ge$
1	$\mathbb{N}_1$		19	3n - 22
2	$\mathbb{N}_2$	7	84	3n - 84
3	$\mathbb{N}_3$	7,7	149	3n - 146
4	$\mathbb{N}_4$	8,7,7	224	3n - 218
5	$\mathbb{N}_5$	8, 8, 7, 7	299	3n - 290
6	$\mathbb{N}_6$	9, 8, 8, 7, 7	384	3n - 372
7	$\mathbb{N}_7$	9, 8, 8, 7, 7, 7	459	3n - 444
8	$\mathbb{N}_8$	10, 8, 8, 8, 7, 7, 7	534	3n - 516
9	$\mathbb{N}_{9}$	10, 9, 8, 8, 8, 7, 7, 7	619	3n - 598
10	$\mathbb{N}_{10}$	10, 9, 8, 8, 8, 8, 7, 7, 7	699	3n - 675
11	$\mathbb{N}_{11}$	11, 9, 8, 8, 8, 8, 8, 7, 7, 7	784	3n - 757
12	$\mathbb{N}_{12}$	11, 9, 9, 8, 8, 8, 8, 7, 7, 7, 7	864	3n - 834
13	$\mathbb{N}_{13}$	11, 10, 9, 8, 8, 8, 8, 8, 8, 7, 7, 7, 7	944	3n - 911
14	$\mathbb{N}_{14}$	12, 10, 9, 8, 8, 8, 8, 8, 8, 7, 7, 7, 7	1024	3n - 988
15	$\mathbb{N}_{15}$	12, 10, 9, 9, 8, 8, 8, 8, 8, 8, 8, 7, 7, 7, 7	1109	3n-1070
16	$\mathbb{N}_{16}$	12, 10, 9, 9, 8, 8, 8, 8, 8, 8, 8, 8, 7, 7, 7, 7	1189	3n - 1147
17	$\mathbb{N}_{17}$	13, 10, 9, 9, 8, 8, 8, 8, 8, 8, 8, 8, 7, 7, 7, 7, 7	1269	3n-1224
18	$\mathbb{N}_{18}$	13, 10, 9, 9, 9, 8, 8, 8, 8, 8, 8, 8, 8, 7, 7, 7, 7, 7	1359	3n - 1311
19	$\mathbb{N}_{19}$	13, 11, 10, 9, 9, 8, 8, 8, 8, 8, 8, 8, 8, 8, 7, 7, 7, 7, 7	1439	3n - 1388
20	$\mathbb{N}_{20}$	13, 11, 10, 9, 9, 8, 8, 8, 8, 8, 8, 8, 8, 8, 8, 7, 7, 7, 7, 7	1519	3n - 1465

Table 1: Number of edges in an edge-maximal graph embeddable in a non-orientable surface.

With  $k := |\sqrt{3/2(g-2)}| + 7$ , we have

$$(k-7)^2 \ge 3/2(g-2)$$
 and  $\alpha_k(g-2) \le 3/(k-7)^2(g-2) \le 2$ 

Let *k* be the minimum integer such that  $\alpha_k(g-2) \leq 2$ . Thus,  $k \leq |\sqrt{3/2(g-2)}| + 7$ . For  $i \in [7, k]$ , define

$$\beta_i \coloneqq [\alpha_i(g-2)]$$
 and  $\gamma_i \coloneqq \beta_i - \alpha_i(g-2)$ .

We claim that  $\beta_k = 2$ . If not, then  $\alpha_k(g - 2) \leq 1$ , implying

$$\frac{12}{(k-7)(k-6)(2k-3)} = \alpha_{k-1}(g-2) - \alpha_k(g-2) > 2 - 1 = 1,$$

which has no solution. Thus,  $\beta_k = 2$ . Define  $\beta_{2g+2} := 1$ . For  $i \in [7, 2g + 2]$ , define

$$L_i := \begin{cases} (\beta_i + 1, \beta_i + 2, \dots, \beta_{i-1}) & \text{if } i \in [8, 2g + 2], \\ (\beta_7 + 1, \beta_7 + 2, \dots, g + 1) & \text{if } i = 7. \end{cases}$$

Edge-Maximal Graphs on Surfaces

Then  $L_{2g+2}, \ldots, L_7$  is a partition of [2, g+1]. Define  $\ell_i := |L_i|$ . Then for  $i \in [8, 2g+2]$ ,

$$(3.2)\ell_i = (\alpha_{i-1} - \alpha_i)(g-2) + (\gamma_{i-1} - \gamma_i) = \frac{12(g-2)}{(i-7)(i-6)(2i-3)} + (\gamma_{i-1} - \gamma_i),$$
  
$$(3.3)\ell_7 = g+1 - \beta_7 = g+1 - \alpha_7(g-2) - \gamma_7 = (1 - \alpha_7)(g-2) - \gamma_7 + 3.$$

It may be that  $\ell_i = 0$  for some values of *i*. (For example, that there is no 12 in  $c_3, \ldots, c_{g+1}$  in the final row of Table 1 corresponds to  $\ell_{12} = 0$ .) If  $\ell_i > 0$  and i < 2g + 2, then let  $i^* := \min\{j > i : \ell_j > 0\}$ . Since  $\ell_{2g+2} > 0$ , this is well defined. Note that  $\ell_j = 0$  for  $j \in [i + 1, i^* - 1]$  and  $\beta_{i^*} + \ell_{i^*} = \beta_i$ . For  $s \in [2, g + 1]$ , there is a unique integer *i* such that  $\ell_i > 0$  and  $s \in L_i$ , in which case define  $c_s := i$ . Thus,  $c_s \ge 7$ . Note that *s* can be uniquely written  $s = \beta_i + z$  for some  $i \in [7, 2g + 2]$  with  $\ell_i > 0$  and  $z \in [1, \ell_i]$ . These definitions are summarised as follows:

$$\begin{split} & L_{2g+2} = (2 = \beta_{2g+2} + 1), \\ & L_{2g+1} = \emptyset, \\ & \vdots \\ & L_{k+1} = \emptyset, \\ & L_k = (3 = \beta_k + 1, \beta_k + 2, \dots, \beta_k + \ell_k = \beta_{k-1}), \\ & \vdots \\ & L_{i^*} = (\beta_{i^*} + 1, \beta_{i^*} + 2, \dots, \beta_{i^*} + \ell_{i^*} = \beta_i), \\ & L_{i^*-1} = \emptyset, \\ & \vdots \\ & L_{i+1} = \emptyset, \\ & L_i = (\beta_i + 1, \beta_i + 2, \dots, \beta_i + \ell_i = \beta_{i-1}), \\ & \vdots \\ & L_8 = (\beta_8 + 1, \beta_8 + 2, \dots, \beta_8 + \ell_8 = \beta_7), \\ & L_7 = (\beta_7 + 1, \beta_7 + 2, \dots, \beta_7 + \ell_7 = g + 1). \end{split}$$

Define

$$f'_g(s) := \begin{cases} 2g+2 & \text{if } s = 2, \\ \max\left\{\frac{2c_s}{c_s-6}(g-2), 2c_s - 3 + f'_g(s-1)\right\} & \text{if } s \ge 3. \end{cases}$$

It follows by induction on *s* that  $f_g(s) \leq f'_g(s)$ . Thus, to prove the desired upper bound on  $f_g(s)$  it suffices to prove the same upper bound on  $f'_g(s)$ . It is helpful to note that  $f'_g(s)$  is calculated by a row-by-row traversal of the above table, where the row corresponding to  $L_i$  uses  $c_s = i$  in the calculation of  $f'_g(s)$ . Thus, for  $s = \beta_i + z$ where  $z \in [1, \ell_i]$ ,

(3.4) 
$$f'_g(\beta_i + z) = f'_g(\beta_i + 1) + (z - 1)(2i - 3).$$

Thus, our focus is on estimating  $f'_g(\beta_i + 1)$ , which equals

$$\max\left\{\frac{2i}{i-6}(g-2),2i-3+f'_g(\beta_i)\right\}.$$

In Claim 1 below, we show that  $\frac{2i}{i-6}(g-2)$  is 'close' to  $2i-3+f'_g(\beta_i)$ . To do so, define the following recursive "error" function. First, let  $E_{2g+2} := 0$  and let  $E_k := 0$ . Then for *i* such that  $\ell_i > 0$ , let

$$E_i := \max\left\{0, \left(\sum_{j=i+1}^{i^*-1} 2\gamma_j\right) + (2i-1)\gamma_i - (2i^*-3)\gamma_{i^*} + E_{i^*}\right\}.$$

*Claim 1* For  $s \in [2, g+1]$ , if  $s = \beta_i + z$  where  $z \in [1, \ell_i]$ , then

$$f'_g(\beta_i+z) \leq \frac{2i}{i-6}(g-2) + (z-1)(2i-3) + E_i.$$

**Proof** We proceed by induction on *s*. First consider the base case s = 2. Then with

i = 2g + 2 we have  $s = \beta_{2g+2} + 1 = \frac{2i}{i-6}(g-2)$ , and the claim holds with  $E_{2g+2} = 0$ . Now assume that  $s \ge 3$ , and the claim holds for s - 1. By (3.4), it suffices to consider the *z* = 1 case, and we can assume that  $\ell_i > 0$ . Then  $s - 1 = \beta_i = \beta_{i^*} + \ell_{i^*}$ . By induction,

$$f'_{g}(\beta_{i^{*}} + \ell_{i^{*}}) \leq \frac{2i^{*}}{i^{*} - 6}(g - 2) + (\ell_{i^{*}} - 1)(2i^{*} - 3) + E_{i^{*}}$$
$$= \frac{2i^{*}}{i^{*} - 6}(g - 2) + \ell_{i^{*}}(2i^{*} - 3) - (2i^{*} - 3) + E_{i^{*}}$$

Since  $\ell_j = 0$  for  $j \in [i + 1, i^* - 1]$ ,

$$f'_g(\beta_{i^*} + \ell_{i^*}) \leq \frac{2i^*}{i^* - 6}(g - 2) + \left(\sum_{j=i+1}^{i^*} (2j - 3)\ell_j\right) - (2i^* - 3) + E_{i^*}.$$

By (3.2) and since  $2i - 3 \le 2i^* - 3$ ,

$$\begin{split} f'_g(\beta_{i^*} + \ell_{i^*}) + 2i - 3 \\ &\leqslant \frac{2i^*}{i^* - 6}(g - 2) + \sum_{j=i+1}^{i^*} (2j - 3) \Big( \frac{12(g - 2)}{(j - 7)(j - 6)(2j - 3)} + \gamma_{j-1} - \gamma_j \Big) + E_{i^*} \\ &= \frac{2i^*}{i^* - 6}(g - 2) + \Big( \sum_{j=i+1}^{i^*} \frac{12(g - 2)}{(j - 7)(j - 6)} \Big) + \Big( \sum_{j=i+1}^{i^*} (\gamma_{j-1} - \gamma_j)(2j - 3) \Big) + E_{i^*} \\ &= \frac{2i^*}{i^* - 6}(g - 2) + (g - 2) \Big( \sum_{j=i+1}^{i^*} \frac{2(j - 1)}{(j - 1) - 6} - \frac{2j}{j - 6} \Big) \\ &+ \Big( \sum_{j=i+1}^{i^*} (\gamma_{j-1} - \gamma_j)(2j - 3) \Big) + E_{i^*} \\ &= \frac{2i^*}{i^* - 6}(g - 2) + (g - 2) \Big( \frac{2i}{i - 6} - \frac{2i^*}{i^* - 6} \Big) + \Big( \sum_{j=i+1}^{i^* - 1} 2\gamma_j \Big) \\ &+ (2i - 1)\gamma_i - \gamma_{i^*}(2i^* - 3) + E_{i^*} \\ &= \frac{2i}{i - 6}(g - 2) + \Big( \sum_{j=i+1}^{i^* - 1} 2\gamma_j \Big) + (2i - 1)\gamma_i - (2i^* - 3)\gamma_{i^*} + E_{i^*}. \end{split}$$

Edge-Maximal Graphs on Surfaces

Since  $c_s = i$  and by (3.2),

$$\begin{aligned} f'_g(s) \\ &= \max\left\{\frac{2i}{i-6}(g-2), 2i-3+f'_g(\beta_{i^*}+\ell_{i^*})\right\} \\ &\leqslant \max\left\{\frac{2i}{i-6}(g-2), \frac{2i}{i-6}(g-2) \\ &+ \left(\sum_{j=i+1}^{i^*-1} 2\gamma_j\right) + (2i-1)\gamma_i - (2i^*-3)\gamma_{i^*} + E_{i^*}\right\} \\ &= \frac{2i}{i-6}(g-2) + \max\left\{0, \left(\sum_{j=i+1}^{i^*-1} 2\gamma_j\right) + (2i-1)\gamma_i - (2i^*-3)\gamma_{i^*} + E_{i^*}\right\} \\ &= \frac{2i}{i-6}(g-2) + E_i. \end{aligned}$$

This completes the proof of the claim.

We now upper bound the  $E_i$ .

*Claim 2* For  $i \in [7, k]$  such that  $\ell_i > 0$ , there are integers  $\delta_i, \ldots, \delta_k$ , such that

$$E_i \leqslant \sum_{j=i}^k \delta_j \gamma_j,$$

and if  $\Delta_i$  is the multiset  $\{\delta_j \ge 0 : j \in [i, k]\}$  and  $\sum \Delta_i := \sum_{j \in [i, k]} \delta_j$ , then  $\sum \Delta_i \le 2k-3$ . Moreover, if  $E_i > 0$ , then  $\delta_i = 2i - 1$ .

**Proof** We proceed by induction on i = k, k-1, ..., 2. In the base case i = k, we have  $E_k = 0$  and the claim holds with  $\delta_k = 0$  and  $X_k = 0$ . Now assume that  $i \in [7, k-1]$  with  $\ell_i > 0$ , and the claim holds for  $i^*$ . Thus, there are integers  $\delta_{i^*}, ..., \delta_k$  such that  $E_{i^*} \leq \sum_{j=i^*}^k \delta_j \gamma_j$  and  $\sum \Delta_{i^*} \leq 2k - 3$ . Moreover, if  $E_{i^*} > 0$ , then  $\delta_{i^*} = 2i^* - 1$ . By definition,

$$E_{i} = \max\left\{0, (2i-1)\gamma_{i} + \left(\sum_{j=i+1}^{i^{*}-1} 2\gamma_{j}\right) - (2i^{*}-3)\gamma_{i^{*}} + E_{i^{*}}\right\}.$$

If  $E_i = 0$ , then the claim holds with  $\delta_i, \ldots, \delta_k = 0$ . Now assume that  $E_i > 0$ . First suppose that  $E_{i^*} = 0$ . Then

$$E_{i} = (2i-1)\gamma_{i} + \left(\sum_{j=i+1}^{i^{*}-1} 2\gamma_{j}\right) - (2i^{*}-3)\gamma_{i^{*}},$$

and the claim holds with  $\delta_i = 2i-1$  and  $\delta_{i^*} = -(2i^*-3)$  and  $\delta_j = 2$  for  $j \in [i+1, i^*-1]$ , in which case  $\Delta_i = \{2i-1, (i^*-1-i) \times 2\}$  and  $\sum \Delta_i = 2i^*-3 \leq 2k-3$ .

Now assume that  $E_{i^*} > 0$ . Then  $\delta_{i^*} = 2i^* - 1$  and

$$E_{i} \leq (2i-1)\gamma_{i} + \left(\sum_{j=i+1}^{i^{*}-1} 2\gamma_{j}\right) - (2i^{*}-3)\gamma_{i^{*}} + \left(\sum_{j=i^{*}}^{k} \delta_{j}\gamma_{j}\right)$$
  
=  $(2i-1)\gamma_{i} + \left(\sum_{j=i+1}^{i^{*}-1} 2\gamma_{j}\right) + ((2i^{*}-1) - (2i^{*}-3))\gamma_{i^{*}} + \left(\sum_{j=i^{*}+1}^{k} \delta_{j}\gamma_{j}\right)$   
=  $(2i-1)\gamma_{i} + \left(\sum_{j=i+1}^{i^{*}-1} 2\gamma_{j}\right) + 2\gamma_{i^{*}} + \left(\sum_{j=i^{*}+1}^{k} \delta_{j}\gamma_{j}\right).$ 

Let  $\delta_i := 2i - 1$  and  $\delta_{i^*} := 2$  and  $\delta_j := 2$  for  $j \in [i + 1, i^* - 1]$ , Observe that

$$\Delta_{i} = (\Delta_{i+1} \setminus \{2i^{*} - 1\}) \cup \{2i - 1, 2, (i^{*} - 1 - i) \times 2\}.$$

Thus,  $\sum \Delta_{i+1} = \sum \Delta_i$ , which is at most 2k-3 by assumption. Thus the claim is satisfied.

Claim 2 with i = 7 implies that there are integers  $\delta_7, \ldots, \delta_k$  such that

$$E_7 \leqslant \sum_{j=7}^k \delta_j \gamma_j$$

and  $\sum \Delta_i \leq 2k - 3$ . Since  $\gamma_j \in [0, 1)$ ,  $E_7 \leq \sum \Delta_7 \leq 2k - 3$ . Claim 1 and (3.3) then imply that for  $s = g + 1 = \beta_7 + \ell_7$ ,

$$f_g(g+1) \leq f'_g(g+1) \leq 14(g-2) + 11(\ell_7 - 1) + E_7$$
  

$$\leq 14(g-2) + 11((1 - \alpha_7)(g-2) - \gamma_7 + 3 - 1) + (2k - 3)$$
  

$$\leq 14(g-2) + 11((1 - \alpha_7)(g-2) + 2) + (2k - 3)$$
  

$$= (25 - 11\alpha_7)(g-2) + 2k + 19$$
  

$$= (25 - 11(\frac{48332}{114345} + \frac{16}{33}\log 2))(g-2) + 2k + 19$$
  

$$= \lambda(g-2) + 2\left[\sqrt{\frac{3}{2}(g-2)}\right] + 33.$$

This completes the proof.

Note that (2.1) implies that  $f_g(g + 1) \ge 5g - 1$ . Since  $\lambda < \frac{50}{3}$ , this shows that Lemma 3.2 is within a factor of  $\frac{10}{3}$  of optimal.

**Theorem 3.3** For every surface  $\Sigma$  of Euler genus g, every edge-maximal embedding of a graph in  $\Sigma$  is at most 84g edges short of a triangulation of  $\Sigma$ .

**Proof** By Lemma 2.7, it suffices to show that  $5f_g(g+1) - 1 \le 84g$ . For  $g \le 299$ , this is verified by direct calculation of the upper bound on  $f_g(g+1)$  in (3.1). For  $g \ge 300$ , by Lemma 3.2,

$$5f_g(g+1) - 1 \le 5\left(16.6534(g-2) + 2\left(1 + \sqrt{\frac{3}{2}(g-2)}\right) + 33\right) - 1 \le 84g.$$

Note that for each surface  $\Sigma$  of Euler genus g, Proposition 2.1 provides examples of edge-maximal cellular embeddings of graphs in  $\Sigma$  that are 3g edges short of a triangulation of  $\Sigma$ . Thus the 84 in Theorem 3.3 cannot be reduced to less than 3. Also

note that  $K_3$ , which is edge-maximal embeddable on any surface  $\Sigma$ , is 3*g* edges short of a triangulation of  $\Sigma$  (since every 3-vertex pseudograph triangulation of  $\Sigma$  has 3*g*+3 edges).

#### 3.1 Orientable Surfaces

Further improvements are possible if we restrict our attention to orientable surfaces. Let G be an edge-maximal graph embeddable in an orientable surface  $\Sigma$ . Recall from Lemma 2.4 that among six consecutive vertices on a face of G, there are at least four distinct vertices, as otherwise a facial walk would contain *abcabc*, implying deg(b) = 2. When  $\Sigma$  is orientable, among five consecutive vertices on a face of G, there are at least four distinct vertices, as otherwise a facial walk would contain *abcab*, repeating ab. This enables us to add more edges to G' in the proof of Lemma 2.7. Consider a face F of G of length  $t \ge 4$ . By Lemma 2.4, F contains four distinct consecutive vertices. Let  $(v_0, v_1, v_2, \dots, v_{t-1})$  be a facial walk of *F*, where  $v_0, v_1, v_2, v_3$  are distinct. Add the edge  $v_0v_i$  to G' whenever  $i \equiv 3 \pmod{4}$  and  $3 \leq i \leq t - 4$ ; this divides F into  $\lfloor \frac{t+1}{4} \rfloor$  faces in G' each containing four distinct vertices (since  $v_0, v_1, v_2, v_3$  are distinct, and every other face contains five consecutive vertices in F, and thus has at least four distinct vertices). Define the graph G'' and set B as above. Consider a face *F* of *G* of length  $t \ge 4$ . Then *B* contains exactly  $\lfloor \frac{t+1}{4} \rfloor$  vertices corresponding to *F*, and  $t-3 \leq 4\lfloor \frac{t+1}{4} \rfloor - 1$ . Thus *G* can be triangulated by adding at most 4|B| - 1 edges. By the same argument used in the proof of Lemma 2.7, G is at most  $4f_g(g+1) - 1$  edges short of a triangulation. This leads, by (3.1), to the results shown in Table 2 and the following theorem.

**Theorem 3.4** For every orientable surface  $\Sigma$  of Euler genus g, every edge-maximal embedding of a graph in  $\Sigma$  is at most 67 g edges short of a triangulation of  $\Sigma$ .

**Proof** By the above discussion it suffices to show that  $4f_g(g + 1) - 1 \le 67g$ . For  $g \le 670$ , this is verified by direct calculation of the upper bound on  $f_g(g + 1)$  in (3.1). For  $g \ge 671$ , by Lemma 3.2,

$$4f_g(g+1) - 1 \leq 4\left(16.6534(g-2) + 2\left(1 + \sqrt{\frac{3}{2}(g-2)}\right) + 33\right) - 1 \leq 67g.$$

#### 4 Multigraph Triangulations

In this section we give the result mentioned as an aside in Section 2 about not needing to introduce loops.

**Proposition 4.1** For every simple graph G with at least three vertices, embedded in a surface  $\Sigma$ , we can add non-loop edges to extend this embedding to a multigraph triangulation.

**Proof** Suppose that *H* is a multigraph containing *G*, and embedded in  $\Sigma$  such that each face has at least three distinct vertices in each boundary component. We claim

g	surface	impurity ≤	$ E(G)  \ge$
2	$\mathbb{S}_1$	67	3n - 67
4	$\mathbb{S}_2$	179	3n - 173
6	$\mathbb{S}_3$	307	3n - 295
8	$\mathbb{S}_4$	427	3n - 409
10	$\mathbb{S}_5$	559	3n - 535
12	$\mathbb{S}_6$	691	3n - 661
14	$\mathbb{S}_7$	819	3n - 783
16	$\mathbb{S}_8$	951	3n - 909
18	S9	1087	3n - 1039
20	$\mathbb{S}_{10}$	1215	3n - 1161
22	$\mathbb{S}_{11}$	1339	3n - 1279
24	$\mathbb{S}_{12}$	1483	3n - 1417
26	$\mathbb{S}_{13}$	1607	3n - 1535
28	$\mathbb{S}_{14}$	1743	3n - 1665
30	$\mathbb{S}_{15}$	1875	3n - 1791
32	$\mathbb{S}_{16}$	2007	3n - 1917
34	$\mathbb{S}_{17}$	2139	3n - 2043
36	$\mathbb{S}_{18}$	2275	3n - 2173
38	$\mathbb{S}_{19}$	2411	3n - 2303
40	$\mathbb{S}_{20}$	2539	3n - 2425

 Table 2:
 Number of edges in an edge-maximal graph embeddable in an orientable surface.

that the embedding extends to a multigraph triangulation. Suppose that this is false, and suppose that *H* is a counterexample with as many edges as possible.

Since *H* is not a triangulation, there is a face *F* for which either (a) there is a single facial walk *f* of length at least 4, or (b) there are at least two boundary components *B* and *B'* of *F*, or (c) the boundary of *F* is a single triangle but the face is not plane.

(a) It suffices to show that there are distinct vertices a and b in the walk f such that both the segment of f from a to b and the return segment of f from b to a contain a vertex not in  $\{a, b\}$ ; for then we can add the edge ab in the face. This replaces F by either (i) a single new face with the same vertices in the boundary, or (ii) two new faces both with a single boundary component containing at least 3 distinct vertices, yielding a contradiction.

If some vertex *a* occurs exactly once in *f*, then we can take *b* to be any vertex other than the predecessor or successor of *a*. So assume that no vertex occurs exactly once in *f*. Consider three consecutive vertices uvw in the walk *f* (where  $v \notin \{u, w\}$ , since there are no loops). If u = w, we can let a = v and let *b* be any vertex in *f* not in  $\{v, w\}$ : for then  $a \neq b$ ,  $w \notin \{a, b\}$ , and *w* occurs in both relevant segments of *f*. If  $u \neq w$ , we may let a = u and b = w: for then  $a \neq b$ ,  $v \notin \{a, b\}$ , and *v* occurs in both relevant segments.

(b) Pick a vertex x in B and a (distinct) vertex y in B', and add the edge xy across the face. This replaces F by a single new face F', replaces B and B' by their union as a boundary component, and leaves any other boundary component untouched, yielding a contradiction.

(c) The face *F* is either  $\mathbb{S}_g$  or  $\mathbb{N}_g$  for some g > 1, with a triangular disk cut out. To handle the orientable case, take a 3-vertex multigraph triangulation of  $\mathbb{S}_g$ , and glue a face onto *F*. Similarly, in the non-orientable case, take a 3-vertex multigraph triangulation of  $\mathbb{N}_g$ , and glue a face onto *F*. In both cases we obtain a contradiction.

# 5 Open Problems

We conclude the paper with a few open problems.

• Let  $c_1$  be the infimum of all numbers c such that every edge-maximal graph embeddable in a surface  $\Sigma$  of Euler genus g is at most cg edges short of a triangulation of  $\Sigma$ . Let  $c_2$  be the infimum of all numbers c such that every edge-maximal embedding in a surface  $\Sigma$  of Euler genus g is at most cg edges short of a triangulation of  $\Sigma$ . Trivially,  $c_1 \leq c_2$ . We have proved that  $3 \leq c_1 \leq c_2 < 84$ . Can these inequalities be improved?

• Are projective planar graphs pure? Are there examples, other than  $K_8 - E(C_5)$ , showing that the class of graphs embeddable in a given surface is impure? Note that [2] conjectured that apart from the sphere, no orientable surface yields a pure class of graphs. This conjecture remains wide open.

• For a surface  $\Sigma$  of Euler genus g, what is the least number k such that for every edge-maximal graph G embeddable in  $\Sigma$ , there is a (simple) triangulation G' of  $\Sigma$  with the same vertex set as G such that E(G) and E(G') have symmetric difference of size at most k? Is k in O(g)?

Acknowledgments This research was initiated at the 2016 Barbados Graph Theory Workshop and the 2016 Workshop on Probability, Combinatorics and Geometry, both held at Bellairs Research Institute in Barbados. Thanks to the workshop organisers, and to the other participants for creating a stimulating working environment. Thanks to Vida Dujmović and Carsten Thomassen for helpful discussions about this research.

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