

ULTRA-SMALL SCALE-FREE GEOMETRIC NETWORKS

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Abstract

We consider a family of long-range percolation models $(G_p)_{p>0}$ on \mathbb{Z}^d that allow dependence between edges and have the following connectivity properties for $p \in (1/d, \infty)$: (i) the degree distribution of vertices in G_p has a power-law distribution; (ii) the graph distance between points \mathbf{x} and \mathbf{y} is bounded by a multiple of $\log_{pd} \log_{pd} |\mathbf{x} - \mathbf{y}|$ with probability $1 - o(1)$; and (iii) an adversary can delete a relatively small number of nodes from $G_p(\mathbb{Z}^d \cap [0, n]^d)$, resulting in two large, disconnected subgraphs.

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1. Introduction

The statistical properties of large networks have received considerable attention in the recent scientific literature [2], [14], [21], [25]. Of special interest are the power-law random networks in which the fraction of vertices of degree k is proportional to k^{-q} for some $q > 0$. Such networks lack an inherent scale and have been termed ‘scale free’. Scale-free graphs are ubiquitous in random network theory and have been proposed as a way to model the behavior of technological, social, and biological networks [1], [21].

Networks often have a geometric component to them where the vertices have positions in space and geographic proximity plays a role in deciding which vertices get connected. In this context, random geometric graphs are a natural alternative to the classical Erdős–Rényi random graph models. Random connection models [20] provide one way to describe networks with spatial content. In these models the event, $E_{\mathbf{x},\mathbf{y}}$, of a connection between points \mathbf{x} and \mathbf{y} has probability $p_{\mathbf{x},\mathbf{y}} := \mathbb{P}[E_{\mathbf{x},\mathbf{y}}] = g(|\mathbf{x} - \mathbf{y}|)$, where $g: \mathbb{R}^+ \rightarrow [0, 1]$ is a connection function and $|\mathbf{x}|$ denotes the Euclidean norm of \mathbf{x} . The standard long-range percolation model assumes independence of $E_{\mathbf{x},\mathbf{y}}$ and $E_{\mathbf{x},\mathbf{u}}$, $\mathbf{y} \neq \mathbf{u}$, which may not be the case in networked systems. Moreover, the degree distribution in this connection model generally does not follow a power law.

Allowing dependency between edges will in general result in technically more complicated models. In this note we show that a natural edge dependency gives rise to a family of long-range percolation models, $(G_p)_{p>0}$, which is technically tractable and which exhibits three connectivity properties for $p \in (1/d, \infty)$. First, G_p has a power-law distribution. Second, G_p is ultra-small, in the sense that the graph distance between lattice points \mathbf{x} and \mathbf{y} is bounded by a multiple of $\log_{pd} \log_{pd} |\mathbf{x} - \mathbf{y}|$ with probability $1 - o(1)$, where $o(1)$ denotes a quantity tending to 0 as $|\mathbf{x} - \mathbf{y}| \rightarrow \infty$. Ultra-small graph distances imply efficiency, are consistent

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with the ‘small-world phenomenon’ [2], [14], [24], [25], and are relevant in the context of routing, searching, and transport of information. Third, an adversary can delete a relatively small number of nodes from $G_p(\mathbb{Z}^d \cap [0, n]^d)$, after which there are two disconnected subgraphs, each containing nearly one-half of the total number of network nodes.

1.1. A general dependent random connection model

Let $\{U_z\}_{z \in \mathbb{Z}^d}$ be independent, identically distributed uniform[0, 1] random variables indexed by \mathbb{Z}^d . Let $p > 0$ and $\delta \in (0, 1]$. For each $z \in \mathbb{Z}^d$, we take δU_z^{-p} to represent a weight at node z defining the radius of the ‘ball of influence’ at z . Consider the graph $G_{p,\delta} := G_{p,\delta}(\mathbb{Z}^d)$ which puts an edge between nodes $x, y \in \mathbb{Z}^d$ whenever each node is contained in the other’s ball of influence. Thus, this connection rule says that the edge (x, y) appears in $G_{p,\delta}(\mathbb{Z}^d)$ whenever

$$|x - y| \leq \delta \min(U_x^{-p}, U_y^{-p}). \tag{1.1}$$

Let $\delta = 1$. By the independence of the U_z , we have $p_{x,y} := P[E_{x,y}] = |x - y|^{-2/p}$, showing that the probability of (there being) long edges in $G_p := G_{p,1}$ increases with p . Edges in G_p have dependent probabilities: if $|y| < |x|$ then the probability of the edge $(0, y)$ given the edge $(0, x)$ is $|y|^{-1/p}$ instead of $|y|^{-2/p}$.

The family of random connection models $G_{p,\delta}$ is disconnected for general p and δ , but not for $\delta = 1$, since having $U_z^{-p} \geq 1$ for all $z \in \mathbb{Z}^d$ implies that adjacent lattice points are connected in G_p . The main results below show, for all $p \in (1/d, \infty)$, that the components of G_p are of arbitrarily large diameter with arbitrarily large probability. Moreover, in accordance with their Poisson Boolean model counterparts (cf. [20]), it is easy to check, for all $\delta \in (0, 1]$ and large p , that the expected number of nodes in the component of $G_{p,\delta}$ containing 0 is infinite, whereas, for p and δ both small, the expected number of such nodes is finite. Our purpose here is to explore the connectivity properties of $G_p, p \in (1/d, \infty)$.

1.2. Main results

Let $D_p(0)$ denote the degree of the origin in $G_p(\mathbb{Z}^d)$, let ω_d denote the volume of the unit-radius ball in \mathbb{R}^d , and let $\alpha := pd - 1$. Our first result shows that if $p \in (1/d, \infty)$ then the degree of a typical vertex follows a power law, i.e. G_p is scale free.

Theorem 1.1. ($G_p(\mathbb{Z}^d)$ has a power-law degree distribution.) *For all $d = 1, 2, \dots$ and all $p \in (1/d, \infty)$,*

$$\lim_{t \rightarrow \infty} t^{1/\alpha} P[D_p(0) > t] = (pd\omega_d/\alpha)^{1/\alpha}.$$

For all $x, y \in \mathbb{Z}^d$, $d_p(x, y)$ denotes the G_p graph distance (‘chemical distance’) between x and y . Our next result says that G_p is ultra-small (cf. [12]), in that $d_p(x, y)$ is bounded by $4(2 + \log \log |x - y|)$ with probability $1 - o(1)$, where throughout, for all $s > 0$, $\log s$ is short for $\log_{pd} s$. We expect that the upper bound in this result can be improved but have not tried to obtain the sharpest bound.

Theorem 1.2. ($G_p(\mathbb{Z}^d)$ has small graph distance.) *For all $d = 1, 2, \dots$ and all $p \in (1/d, \infty)$,*

$$\frac{d_p(0, x)}{2 + \log \log |x|} \leq 4$$

with probability $1 - o(1)$, where $o(1)$ tends to 0 as $|x| \rightarrow \infty$.

The network failure of $G_p(\mathbb{Z}^d)$ is easily quantified, as follows.

Theorem 1.3. (Network failure.) *For all $d = 1, 2, \dots$ and all $p \in (1/d, \infty)$, an adversary can delete N nodes from $G_p(\mathbb{Z}^d \cap [0, n]^d)$, where $E[N] = O(n^{d-1}[n^{1-1/p} \vee 1])$, resulting in two disconnected subgraphs on vertex sets of cardinality at least $n^d/2 - N$.*

Theorem 1.3 implies, in particular, that if $p \in (1/d, 1)$ then removing roughly $O(n^{d-1})$ nodes may reduce $G_p(\mathbb{Z}^d \cap [0, n]^d)$ to two large, disconnected subgraphs.

Remarks. 1. *Standard long-range percolation models.* Assume that $p_{x,y} := P[E_{x,y}] = |x - y|^{-s+o(1)}$ as $|x - y| \rightarrow \infty$, for some constant $s \in (0, \infty)$; $E_{x,y}$ and $E_{x,u}$ are independent for all $x, y, u \in \mathbb{Z}^d$. For $s \in (0, d)$, Benjamini *et al.* [4] showed that the graph distance $d(\mathbf{0}, \mathbf{x})$ behaves like the constant $\lceil s/(d - s) \rceil$ as $|\mathbf{x}| \rightarrow \infty$. Here, $\lceil x \rceil$ denotes the greatest integer less than x . For $s = d$, Coppersmith *et al.* [13] showed that $d(\mathbf{0}, \mathbf{x})$ scales as $\log |\mathbf{x}|/\log \log |\mathbf{x}|$, whereas, for $s \in (d, 2d)$, Biskup [7], [8] showed that $d(\mathbf{0}, \mathbf{x})$ scales as $(\log |\mathbf{x}|)^{\Delta+o(1)}$, where $\Delta := \Delta(s, d) := \log 2/\log(2d/s)$. The case $s = 2d$ is open and, for $s \in (2d, \infty)$, $d(\mathbf{0}, \mathbf{x})$ scales at least linearly in $|\mathbf{x}|$, as shown by Berger [5]. The different scalings for the standard long-range percolation model suggest that G_p also has different scalings for $p \in (0, 1/d)$, but we have not determined them. Kleinberg [19] proposed a lattice model where long-range contacts are added in a biased way, there being, however, a uniform bound on the number of such contacts.

2. *Geometric networks in \mathbb{R}^d .* We expect that Theorems 1.1–1.3 extend to analogously defined continuum models on Poisson point sets in \mathbb{R}^d . This would add to the following related results.

- (a) Let $f: \mathbb{R}^d \rightarrow \mathbb{R}^+$ and let \mathcal{P}_f be a Poisson point process on \mathbb{R}^d with intensity f . The *geometric graph*, described in depth by Penrose [23], joins two nodes in \mathcal{P}_f whenever their Euclidean distance is less than a specified cutoff. Hermann *et al.* [18, Section II.B] showed that if $\int_{\mathbb{R}^d} f^r(\mathbf{x}) \, d\mathbf{x} = \infty$ for all $r > r_0$, then the degree distribution is effectively a power law.
- (b) The *on-line nearest-neighbors graph* is defined on randomly ordered point sets in \mathbb{R}^d , and places an edge between each point and its nearest neighbor amongst the points preceding it in the ordering. Such graphs have scale-free properties over certain degree domains [6], [16].
- (c) Franceschetti and Meester [17] developed a scale-free continuum model but did not obtain iterated log bounds on interpoint graph distances.
- (d) The standard *Boolean connection model* puts an edge between \mathbf{x} and \mathbf{y} whenever the respective balls of influence overlap. In the context of (1.1), (\mathbf{x}, \mathbf{y}) is an edge whenever $|\mathbf{x} - \mathbf{y}| \leq \delta(U_{\mathbf{x}}^{-p} + U_{\mathbf{y}}^{-p})$. These models are not in general scale free.

3. *Power exponents $q \in (2, 3)$.* Consider a random graph on n nodes v_1, v_2, \dots, v_n with weight (expected degree) w_i at node v_i . Nodes v_i and v_j are connected with probability $\rho w_i w_j$, where $\rho = (\sum_{i=1}^n w_i)^{-1}$. Chung and Lu [10], [11] provided conditions on the weights under which the degree distribution is proportional to k^{-q} , $q \in (2, 3)$, $k \in \mathbb{Z}$, the average distance between nodes is almost surely $O(\log \log n)$, and the diameter is $O(\log n)$. In unrelated work, Cohen and Havlin [12] argued that whenever the degree distribution of a random graph on n vertices is proportional to k^{-q} , where $q \in (2, 3)$, k is restricted to (m, K) , and where m and $K := K(n)$ are well-defined ‘cutoffs’, then the diameter behaves like $\log \log n$.

4. *Preferential attachment models.* These dynamic graphs evolve with time in such a way that a newly arriving vertex connects to an existing vertex with a probability proportional to the degree of the (latter) vertex. Thus, nodes of high degree tend to acquire more new links than do nodes of low degree. Albert and Barabási [1] showed that such models follow a power law, are not geometry dependent, and, as shown by Bollabás and Riordan [9], are not ultra-small in general.

5. *Degree dependence on p .* Theorem 1.1 tells us that $P[D_p(\mathbf{0}) = k] \sim Ck^{-q}$, where $q := pd/(pd - 1)$. Thus, as p increases on $(1/d, \infty)$, the exponent of the degree distribution, q , decreases to 1.

6. *Further connectivity results.* Theorems 1.1–1.3 describe the connectivity of $G_p(\mathbb{Z}^d)$. Further analysis of the connectivity of $G_p(\mathbb{Z}^d)$, such as thermodynamic and Gaussian limits for the number of three cycles (or other clustering coefficients) on $G_p(\mathbb{Z}^d \cap [0, n]^d)$, is simplified by appealing to the stabilization properties of G_p (see especially [22]). $G_p(\mathbb{Z}^d)$ is *assortative* in that high-degree nodes tend to link to high-degree nodes and low-degree nodes tend to link to low-degree nodes.

7. *The case $p \in (0, 1/d)$.* If $p \in (0, 1/d)$ then G_p has few long edges and the proofs of the scale-free and ultra-small properties break down. The scalar $1/d$ thus represents the boundary between graphs that are ultra-small scale free and those which are not.

2. Proof of Theorem 1.1

Throughout, we adopt the following notation: $B_r(\mathbf{x})$ denotes the Euclidean ball of radius r centered at $\mathbf{x} \in \mathbb{R}^d$, $L_r(\mathbf{x}) := B_r(\mathbf{x}) \cap \mathbb{Z}^d \setminus \{\mathbf{x}\}$ denotes the lattice points a distant at most r from \mathbf{x} , and C denotes a generic positive constant whose value may change from line to line. The underlying probability space is $\Omega := [0, 1]^{\mathbb{Z}^d}$ and is equipped with the product probability measure $P := \mu^{\mathbb{Z}^d}$, where μ is the uniform probability measure on $[0, 1]$. Conditional on $U_{\mathbf{0}} = u$, $D_p(\mathbf{0})$ is the number of points \mathbf{y} in $L_{u^{-p}}(\mathbf{0})$ with weight, $U_{\mathbf{y}}^{-p}$, exceeding $|y|$; hence, $U_{\mathbf{y}} \in [0, |y|^{-1/p}]$. Writing $D(u^{-p})$ for the value of $D_p(\mathbf{0})$ conditioned on $\mathbf{0}$ having weight u^{-p} , we have

$$D(u^{-p}) = \sum_{\mathbf{y} \in L_{u^{-p}}(\mathbf{0})} 1_{\{U_{\mathbf{y}} \leq |y|^{-1/p}\}}.$$

Thus, to prove Theorem 1.1 we condition on $U_{\mathbf{0}}$ and show that

$$\lim_{t \rightarrow \infty} t^{1/\alpha} \int_0^1 P[D(u^{-p}) > t] du = \left(\frac{pd\omega_d}{\alpha}\right)^{1/\alpha}, \tag{2.1}$$

where, recall, $\alpha := pd - 1$. The next lemma will be useful in establishing (2.1). Let $\beta := pd\omega_d/\alpha$.

Lemma 2.1. *For all $p \in (1/d, \infty)$, we have*

$$E[D(u^{-p})] = \beta u^{-\alpha} + O(\max(1, u^{-pd+p+1})), \tag{2.2}$$

where the error on the right-hand side of (2.2) holds as $u \rightarrow 0^+$.

Proof. Note that $E[D(u^{-p})]$ is approximated by

$$\int_{|\mathbf{x}| \leq u^{-p}} |\mathbf{x}|^{-1/p} d\mathbf{x} = d\omega_d \int_0^{u^{-p}} t^{d-1-1/p} dt = \beta u^{-\alpha}.$$

Let $R := R(u)$ be the maximal collection of grid cubes (cubes centered at points in \mathbb{Z}^d with edge length 1) contained within $B_{u^{-p}}(\mathbf{0})$. The approximation error

$$\left| E[D(u^{-p})] - \int_{|\mathbf{x}| \leq u^{-p}} |\mathbf{x}|^{-1/p} d\mathbf{x} \right|$$

is bounded by the sum of the following three errors:

$$\begin{aligned} E_1 &:= \left| E[D(u^{-p})] - \sum_{\mathbf{y} \in R(u) \cap \mathbb{Z}^d, \mathbf{y} \neq \mathbf{0}} |\mathbf{y}|^{-1/p} \right|, \\ E_2 &:= \left| \sum_{\mathbf{y} \in R(u) \cap \mathbb{Z}^d, \mathbf{y} \neq \mathbf{0}} |\mathbf{y}|^{-1/p} - \int_{R(u)} |\mathbf{x}|^{-1/p} d\mathbf{x} \right|, \\ E_3 &:= \left| \int_{R(u)} |\mathbf{x}|^{-1/p} d\mathbf{x} - \int_{|\mathbf{x}| \leq u^{-p}} |\mathbf{x}|^{-1/p} d\mathbf{x} \right|. \end{aligned}$$

Now,

$$E_1 = \sum_{\mathbf{y} \in (B_{u^{-p}}(\mathbf{0}) \setminus R(u)) \cap \mathbb{Z}^d, \mathbf{y} \neq \mathbf{0}} |\mathbf{y}|^{-1/p}$$

and, so, is bounded by the product of

$$\text{card}\{(B_{u^{-p}}(\mathbf{0}) \setminus R(u)) \cap \mathbb{Z}^d\} \quad \text{and} \quad \sup\{|\mathbf{y}|^{-1/p} : \mathbf{y} \in (B_{u^{-p}}(\mathbf{0}) \setminus R(u)) \cap \mathbb{Z}^d\}.$$

Since the first factor is bounded by $Cu^{-p(d-1)}$ and the second by Cu , it follows that $E_1 \leq Cu^{-pd+p+1}$. A similar method shows that $E_3 \leq Cu^{-pd+p+1}$.

We estimate E_2 as follows. For all $\mathbf{y} \in \mathbb{Z}^d$, let $Q_{\mathbf{y}}$ denote the grid cube with center \mathbf{y} . For all $s = 1, 2, \dots$, let $M(s) := \text{card}\{\mathbf{y} \in \mathbb{Z}^d : |\mathbf{y}| \in [s, s+1)\}$. Since there is a constant $C > 0$ such that, for all $\mathbf{x} \in Q_{\mathbf{y}}$ and all $\mathbf{y} \in \mathbb{Z}^d$,

$$||\mathbf{y}|^{-1/p} - |\mathbf{x}|^{-1/p}| \leq C|\mathbf{y}|^{-1/p-1},$$

it follows that

$$E_2 \leq C \sum_{s=1}^{u^{-p}} s^{-1/p-1} M(s) \leq C \sum_{s=1}^{u^{-p}} s^{-1/p+d-2} \leq C \max(1, u^{-pd+p+1}),$$

since $M(s) \leq Cs^{d-1}$. Combining the bounds for $E_1, E_2,$ and E_3 yields Lemma 2.1.

Letting $s := u^{-p}$ in (2.1), note that, to prove Theorem 1.1, it suffices to show that

$$\lim_{t \rightarrow \infty} t^{1/\alpha} \int_1^\infty \mathbb{P}[D(s) > t] \frac{1}{p} s^{-1/p-1} ds = \beta^{1/\alpha}. \tag{2.3}$$

We observe that (2.3) is plausible because Lemma 2.1 suggests that $\mathbb{P}[D(s) > t]$ is close to 1 for $t \ll \beta s^{\alpha/p}$ and close to 0 for $t \gg \beta s^{\alpha/p}$, indicating that the left-hand side of (2.3) behaves like

$$\lim_{t \rightarrow \infty} t^{1/\alpha} \int_{(t/\beta)^{p/\alpha}}^\infty \frac{1}{p} s^{-1/p-1} ds = \beta^{1/\alpha}.$$

To put this heuristic argument on a rigorous footing, we rewrite the integral in (2.3) as a sum of two integrals. The first integral is estimated via Bernstein’s inequality and the second is handled using Poisson approximation arguments. We do this as follows.

For all $v > 0$, let $m(v) := \sup\{s : E[D(s)] \leq v\}$. Recalling that $\alpha := pd - 1$, from Lemma 2.1 we obtain

$$E[D(s)] = \beta s^{\alpha/p} + O(\max(1, s^{d-1-1/p})) = \beta s^{\alpha/p}(1 + \max(O(s^{1/p-d}), O(s^{-1}))). \tag{2.4}$$

It follows, for large v and $p \in (1/d, \infty)$, that

$$m(v) = \left(\frac{v}{(1 + o(1))\beta} \right)^{p/\alpha},$$

where $o(1)$ tends to 0 as $v \rightarrow \infty$. Given fixed $t \geq \beta$ and $\varepsilon \in (0, \frac{1}{2})$, define the following two integration domains:

$$I_1 := [1, m(t - t^{1/2+\varepsilon})], \quad I_2 := [m(t - t^{1/2+\varepsilon}), \infty).$$

Rewrite the left-hand side of (2.3) as

$$\lim_{t \rightarrow \infty} t^{1/\alpha} \int_{I_1} P[D(s) > t] \frac{1}{p} s^{-1/p-1} ds + \lim_{t \rightarrow \infty} t^{1/\alpha} \int_{I_2} P[D(s) > t] \frac{1}{p} s^{-1/p-1} ds =: S_1 + S_2,$$

provided that both limits exist.

To prove Theorem 1.1 it suffices to show that $S_1 = 0$ and $S_2 = \beta^{1/\alpha}$. We first show that $S_1 = 0$. Bernstein’s inequality [15, p. 12] for sums of independent, bounded random variables yields, for all $s \in I_1$,

$$P[D(s) > t] \leq \exp\left(\frac{-(t - E[D(s)])^2}{2 E[D(s)] + 4t/3}\right).$$

Using the bounds $\inf_{s \in I_1} (t - E[D(s)]) \geq t^{1/2+\varepsilon}$ and $\sup_{s \in I_1} E[D(s)] \leq t - t^{1/2+\varepsilon} < t$, for all $s \in I_1$ we thus obtain

$$P[D(s) > t] \leq \exp\left(\frac{-(t^{1/2+\varepsilon})^2}{10t/3}\right) = \exp\left(-\frac{3t^{2\varepsilon}}{10}\right).$$

It follows that

$$S_1 \leq \limsup_{t \rightarrow \infty} t^{1/\alpha} \exp\left(-\frac{3t^{2\varepsilon}}{10}\right) \int_1^\infty \frac{1}{p} s^{-1/p-1} ds = 0.$$

We next show that $S_2 = \beta^{1/\alpha}$. By approximating $D(s)$ with a Poisson random variable we establish the following simplified expression for S_2 . Here and elsewhere, $Po(\lambda)$ denotes a Poisson random variable with mean λ .

Lemma 2.2. *For all $p \in (1/d, \infty)$, we have*

$$S_2 = \lim_{t \rightarrow \infty} t^{1/\alpha} \int_{m(t-t^{1/2+\varepsilon})}^\infty P[Po(E[D(s)]) > t] \frac{1}{p} s^{-1/p-1} ds.$$

Proof. For all $y \in \mathbb{Z}^d$, let $p_y := E[1_{\{U_y \leq |y|^{-1/p}\}}] = |y|^{-1/p}$. Letting d_{TV} be the total variation distance, it follows from well-known Poisson approximation bounds (e.g. Equation (1.23)

of [3]) that

$$d_{TV}(D(s), \text{Po}(E[D(s)])) \leq \left(\sum_{y \in L_s(\mathbf{0})} p_y \right)^{-1} \sum_{y \in L_s(\mathbf{0})} p_y^2.$$

By an analysis similar to that in the proof of Lemma 2.1 and (2.4), for $d > 2/p$ we obtain

$$\sum_{y \in L_s(\mathbf{0})} p_y^2 = \frac{pd\omega_d}{pd - 2} s^{d-2/p} (1 + o(1)),$$

whereas, for $1/p < d \leq 2/p$, we have

$$\sum_{y \in L_s(\mathbf{0})} p_y^2 = O(1).$$

It follows from Lemma 2.1 that, for $d > 2/p$, we obtain

$$\begin{aligned} d_{TV}(D(s), \text{Po}(E[D(s)])) &\leq (\beta s^{d-1/p} (1 + o(1)))^{-1} \beta \left(\frac{pd - 1}{pd - 2} \right) s^{d-2/p} (1 + o(1)) \\ &= O(s^{-1/p}), \end{aligned}$$

whereas, for $1/p < d \leq 2/p$, we have

$$d_{TV}(D(s), \text{Po}(E[D(s)])) = O(s^{-d+1/p}).$$

Letting

$$e(s, t) := P[D(s) > t] - P[\text{Po}(E[D(s)]) > t],$$

it follows that, uniformly in $t \in (0, \infty)$, we have $|e(s, t)| = O(s^{-\xi})$, where $\xi = 1/p$ for $d > 2/p$ and $\xi = d - 1/p$ for $1/p < d \leq 2/p$. We now rewrite S_2 as

$$S_2 = \lim_{t \rightarrow \infty} t^{1/\alpha} \int_{m(t-t^{1/2+\varepsilon})}^{\infty} (P[\text{Po}(E[D(s)]) > t] + e(s, t)) \frac{1}{p} s^{-1/p-1} ds$$

and show that the term containing $e(s, t)$ is negligible.

Recall that

$$m(t - t^{1/2+\varepsilon}) = \left(\frac{t - t^{1/2+\varepsilon}}{(1 + o(1))\beta} \right)^{p/\alpha},$$

where, here and in the remainder of this section, $o(1)$ tends to 0 as $t \rightarrow \infty$. It follows that

$$\int_{m(t-t^{1/2+\varepsilon})}^{\infty} e(s, t) s^{-1/p-1} ds = O\left(\int_{m(t-t^{1/2+\varepsilon})}^{\infty} s^{-\xi-1/p-1} ds \right) = O(t^{-p/\alpha(\xi+1/p)})$$

and, therefore, that

$$\lim_{t \rightarrow \infty} t^{1/\alpha} \int_{m(t-t^{1/2+\varepsilon})}^{\infty} e(s, t) s^{-1/p-1} ds = 0.$$

We thus obtain Lemma 2.2.

It is now straightforward to show that $S_2 = \beta^{1/\alpha}$. Letting $z := \beta s^{d-1/p}/t$, whence $s = (tz/\beta)^{p/\alpha}$ and $E[D(s)] = tz(1 + O((tz)^{-\rho}))$ with $\rho := \rho(p, d) > 0$, we obtain, via Lemma 2.2,

$$S_2 = \lim_{t \rightarrow \infty} \frac{\beta^{1/\alpha}}{\alpha} \int_{1+o(1)}^{\infty} P[\text{Po}(tz(1 + O((tz)^{-\rho}))) > t] z^{-1/\alpha-1} dz.$$

The integrability of the integrand on $[1 + o(1), \infty)$ gives, for all $\gamma > 0$,

$$S_2 = \lim_{t \rightarrow \infty} \frac{\beta^{1/\alpha}}{\alpha} \int_{1+\gamma}^{\infty} \mathbb{P}[\text{Po}(tz(1 + O((tz)^{-\rho}))) > t] z^{-1/\alpha-1} dz + \gamma \cdot O(1).$$

For all $z \in [1 + \gamma, \infty)$, we have $\mathbb{P}[\text{Po}(tz(1 + O((tz)^{-\rho}))) > t] \rightarrow 1$ as $t \rightarrow \infty$. The dominated convergence theorem yields

$$S_2 = \frac{\beta^{1/\alpha}}{\alpha} \int_1^{\infty} z^{-1/\alpha-1} dz + \gamma O(1) = \beta^{1/\alpha} + \gamma \cdot O(1).$$

Now let $\gamma \rightarrow 0$ to obtain $S_2 = \beta^{1/\alpha}$, as desired.

3. Proof of Theorem 1.2

We prove Theorem 1.2 by showing, for all $\mathbf{x} \in \mathbb{Z}^d$, the existence of an event $E := E(\mathbf{x}) \subset \Omega$, with $\mathbb{P}[E] = 1 - o(1)$, such that on E there is a path π consisting of N edges in $G_p(\mathbb{Z}^d)$ joining $\mathbf{0}$ to \mathbf{x} , where $N \leq 4(2 + \log \log |\mathbf{x}|)$. Here and in the sequel, $o(1)$ denotes a quantity tending to 0 as $|\mathbf{x}| \rightarrow \infty$.

Constructing the path π would be easy if the balls of influence at $\mathbf{0}$ and \mathbf{x} both had radius at least $|\mathbf{x}|$, for then π would consist merely of the single edge $(\mathbf{0}, \mathbf{x})$. In general, the balls of influence at $\mathbf{0}$ and \mathbf{x} have much smaller radii and the path π thus needs to join a sequence of balls such that consecutive balls contain each other's centers.

The heart of the proof will consist of constructing a sequence of nodes of cardinality roughly $2 \log \log |\mathbf{x}|$ with these properties: the first node, $\mathbf{0}'$, is at distance at most $\frac{1}{2} \log \log |\mathbf{x}|$ from $\mathbf{0}$; the last node, \mathbf{x}' , is at distance at most $\frac{1}{2} \log \log |\mathbf{x}|$ from \mathbf{x} ; and the edges defined by consecutive nodes are in G_p , i.e. the balls of influence at consecutive nodes contain each other's centers. Since $\mathbf{0}$ and $\mathbf{0}'$ can be joined with a path of at most $\log \log |\mathbf{x}|$ edges, and likewise for \mathbf{x} and \mathbf{x}' , we can obtain a path π consisting of roughly $4 \log \log |\mathbf{x}|$ edges. The construction of this sequence of nodes depends critically on an intermediate node, denoted here by \mathbf{P}_0 , that has an unusually large ball of influence. Before defining $\mathbf{0}'$, \mathbf{P}_0 , and \mathbf{x}' , we need some terminology.

For all $\mathbf{x} \in \mathbb{R}^d$ and $r > 0$, let $L_r^+(\mathbf{x})$ and $L_r^-(\mathbf{x})$ denote the lattice points in the upper and lower hemispheres of radius r centered at \mathbf{x} . That is, $L_r^+(\mathbf{x}) := B_r(\mathbf{x}) \cap (\mathbb{Z}^{d-1} \times \mathbb{Z}^+)$ and, similarly, $L_r^-(\mathbf{x}) := B_r(\mathbf{x}) \cap (\mathbb{Z}^{d-1} \times \mathbb{Z}^-)$. Here $\mathbb{Z}^+ := \{1, 2, \dots\}$ and $\mathbb{Z}^- := \{-1, -2, \dots\}$.

3.1. Definition of $\mathbf{0}'$, \mathbf{P}_0 , and \mathbf{x}'

Throughout, we appeal to the following elementary fact. Recall that $\log s$ is short for $\log_{pd} s$.

Lemma 3.1. *Let U_1, \dots, U_n be independent and identically uniformly distributed on $[0, 1]$. Then, for all $n > pd$, we have*

$$\min_{i \leq n} U_i \leq \frac{K \log n}{n}$$

with probability at least $1 - n^{-K}$.

In the sequel, we fix K to be large, with a value to be determined later.

3.1.1. *Definition of $\mathbf{0}'$.* Let $E_0 := E_0(\mathbf{x})$ be the event that there is a node $\mathbf{z} \in L_{(1/2) \log \log |\mathbf{x}|}^-(\mathbf{0})$ such that

$$U_{\mathbf{z}} \leq \frac{K \log(\log \log |\mathbf{x}|)^d}{(\log \log |\mathbf{x}|)^d}.$$

Clearly, E_0 depends only on $U_z, z \in L_{(1/2)\log\log|\mathbf{x}|}^-(\mathbf{0})$. In case more than one node in $L_{(1/2)\log\log|\mathbf{x}|}^-(\mathbf{0})$ satisfies the last bound, we choose z to be that node with smallest lexicographical order.

By Lemma 3.1, $P[E_0] \geq 1 - C(\log\log|\mathbf{x}|)^{-dK}$. Given E_0 we let $\theta' := z$. Note that θ' is random and that, since $pd > 1$, for all large $|\mathbf{x}|$ we have

$$U_{\theta'}^{-p} \geq 2 \log\log|\mathbf{x}|. \tag{3.1}$$

Inequality (3.1) will be important in the sequel. For now note that, since $G_p(\mathbb{Z}^d)$ connects adjacent lattice points, it follows that $d_p(\mathbf{y}, \mathbf{x}) \leq 2|\mathbf{y} - \mathbf{x}|$ for all $\mathbf{x}, \mathbf{y} \in \mathbb{Z}^d$, i.e. that

$$d_p(\mathbf{0}, \theta') \leq \log\log|\mathbf{x}|. \tag{3.2}$$

3.1.2. *Definition of \mathbf{x}' .* Similarly, given \mathbf{x} , with probability at least $1 - C(\log\log|\mathbf{x}|)^{-dK}$ there is an event E_x such that there is a node $\mathbf{x}' \in L_{(1/2)\log\log|\mathbf{x}|}^-(\mathbf{x})$ on E_x with weight

$$U_{\mathbf{x}'}^{-p} \geq 2 \log\log|\mathbf{x}|.$$

Clearly $d_p(\mathbf{x}, \mathbf{x}') \leq \log\log|\mathbf{x}|$ and E_x depends only on $U_z, z \in L_{(1/2)\log\log|\mathbf{x}|}^-(\mathbf{x})$.

3.1.3. *Definition of P_0 .* Assume without loss of generality that the components of \mathbf{x} have even parity, meaning that $\mathbf{x}/2 \in \mathbb{Z}^d$. Consider the event, $E_{\mathbf{x}/2}$, that there is a node $P_0 \in L_{|\mathbf{x}|/10}(\mathbf{x}/2)$ with

$$U_{P_0} \leq \frac{K \log(|\mathbf{x}|)^d}{|\mathbf{x}|^d}.$$

Lemma 3.1 implies that $P[E_{\mathbf{x}/2}] \geq 1 - C(|\mathbf{x}|^{-dK})$. We note that, for large $|\mathbf{x}|$,

$$U_{P_0}^{-p} \geq 2|\mathbf{x}| \tag{3.3}$$

since $pd > 1$.

3.2. Construction of the path π via θ', P_0 , and \mathbf{x}'

It will suffice to show that there is an event $E := E(\mathbf{x})$, with $P[E(\mathbf{x})] = 1 - o(1)$, such that on E there are two paths, each having at most $2 + 2\lceil\log\log|\mathbf{x}|\rceil$ edges, with one path joining P_0 to $\mathbf{0}$ and the other joining P_0 to \mathbf{x} . It will be enough to show the existence of a path between P_0 and $\mathbf{0}$, for the method can be repeated verbatim to yield the path between P_0 and \mathbf{x} . We first introduce some additional terminology.

We abbreviate our notation by letting $b := pd$. Note that $b > 1$ by assumption. Fix $\varepsilon \in (0, 1)$ and $\mathbf{x} \in \mathbb{Z}^d$ with $|\mathbf{x}|$ large. For all $j = 1, 2, \dots$, let

$$r_j := r_j(\mathbf{x}, \varepsilon) := |\mathbf{x}|^{b^{-j(1-\varepsilon)}}$$

and note that $r_j \downarrow 1$ and $1 < r_j < |\mathbf{x}|$ for all $j = 1, 2, \dots$. We record an elementary fact.

Lemma 3.2. $r_{j+1} = r_j^{\beta(p,d,\varepsilon)}$, where $\beta(p, d, \varepsilon) := b^{-1+\varepsilon}$.

For all $j = 1, 2, \dots$, consider the following disjoint ‘semi-annular’ regions of lattice points:

$$A_j := [(L_{r_j}^+(\theta') - L_{r_{j+1}}^+(\theta')) \setminus L_{|\mathbf{x}|/10}^+(\mathbf{x}/2)].$$

The construction of the path joining P_0 to $\mathbf{0}$ is facilitated by the following four lemmas. The first three lemmas show that, for all $j, 1 \leq j \leq \lceil\log\log|\mathbf{x}|\rceil + 1$, there are points $P_j \in A_j$ such that (P_j, P_{j-1}) and $(P_{\lceil\log\log|\mathbf{x}|\rceil+1}, \theta')$ belong to $G_p(\mathbb{Z}^d)$. The fourth lemma shows that this happens on an event with probability $1 - o(1)$. By consecutively linking $P_j, 0 \leq j \leq \lceil\log\log|\mathbf{x}|\rceil + 1$, and θ' , we construct a path joining P_0 to θ' with $\lceil\log\log|\mathbf{x}|\rceil + 2$

edges. Since $\mathbf{0}'$ is within $\frac{1}{2} \log \log |\mathbf{x}|$ of $\mathbf{0}$, we need at most $\lceil \log \log |\mathbf{x}| \rceil$ edges to join $\mathbf{0}'$ to $\mathbf{0}$ (recall (3.2)). This gives a path joining \mathbf{P}_0 to $\mathbf{0}$ with at most $2\lceil \log \log |\mathbf{x}| \rceil + 2$ edges. Since $2 + 2\lceil \log \log |\mathbf{x}| \rceil \leq 4 + 2 \log \log |\mathbf{x}|$, we obtain Theorem 1.2, as desired. We now turn to our four lemmas.

Lemma 3.3. *There exists an event E_1 , with $P[E_1] = 1 - O(r_1^{-dK})$, such that on E_1 there is a node $\mathbf{P}_1 \in A_1$ which is linked to \mathbf{P}_0 , i.e. the edge $(\mathbf{P}_0, \mathbf{P}_1)$ is in $G_p(\mathbb{Z}^d)$.*

Proof. The number of lattice points in A_1 is $\Theta(|\mathbf{x}|^{db^{-1+\varepsilon}})$, i.e. there is a constant $C > 0$ such that the number of lattice points is bounded from above by $C|\mathbf{x}|^{db^{-1+\varepsilon}}$ and bounded from below by $C^{-1}|\mathbf{x}|^{db^{-1+\varepsilon}}$. Lemma 3.1 implies that there is an event E_1 , depending only on $\{U_z\}_{z \in A_1}$ and with

$$P[E_1] = 1 - O(|\mathbf{x}|^{-dKb^{-1+\varepsilon}}),$$

such that, for large $|\mathbf{x}|$, E_1 implies the existence of $\mathbf{P}_1 \in A_1$ with

$$U_{\mathbf{P}_1} \leq \frac{K \log(|\mathbf{x}|^{db^{-1+\varepsilon}})}{|\mathbf{x}|^{db^{-1+\varepsilon}}}.$$

Again, if there is more than one node in A_1 satisfying this inequality, we choose the one with smallest lexicographical order. Since $b := pd$ it follows for large $|\mathbf{x}|$ that \mathbf{P}_1 has weight

$$U_{\mathbf{P}_1}^{-p} \geq \frac{|\mathbf{x}|^{b\varepsilon}}{(K \log(|\mathbf{x}|^{db^{-1+\varepsilon}}))^p} \geq 2|\mathbf{x}|. \tag{3.4}$$

We now show that \mathbf{P}_1 is linked to \mathbf{P}_0 . It suffices to show that

$$|\mathbf{P}_0 - \mathbf{P}_1| \leq \min(U_{\mathbf{P}_0}^{-p}, U_{\mathbf{P}_1}^{-p}).$$

However, $|\mathbf{P}_0 - \mathbf{P}_1| \leq |\mathbf{P}_0| + |\mathbf{P}_1| \leq 2|\mathbf{x}|$, so Lemma 3.3 follows from (3.3) and (3.4).

Given \mathbf{x} , let $m := m(\mathbf{x})$ denote the largest integer such that $r_m \geq \log \log |\mathbf{x}|$; m is well defined since $r_j \downarrow 1$. If $t := \lceil 1/(1 - \varepsilon) \rceil \log \log |\mathbf{x}|$ then

$$|\mathbf{x}|^{b^{-t(1-\varepsilon)}} = |\mathbf{x}|^{1/\log |\mathbf{x}|} = b,$$

showing that m is bounded by t . The next lemma extends the arguments of Lemma 3.3 and builds a path of m edges from \mathbf{P}_0 to a node $\mathbf{P}_m \in A_m$.

Lemma 3.4. *For all j , $1 \leq j \leq m$, there is an event E_j , depending only on $\{U_z\}_{z \in A_j}$, such that*

- (i) $P[E_j] = 1 - O(r_j^{-dK})$, and
- (ii) *on each E_j there is a node $\mathbf{P}_j \in A_j$ such that, on $E_{j-1} \cap E_j$, the edge $(\mathbf{P}_{j-1}, \mathbf{P}_j)$ is in G_p .*

Proof. Since $\text{card}\{A_j\} = \Theta(r_j^d)$, Lemma 3.1 implies that for large $|\mathbf{x}|$ there is an event E_j , with $P[E_j] = 1 - O(r_j^{-dK})$, which depends only on $\{U_z\}_{z \in A_j}$ and which implies the existence of a $\mathbf{P}_j \in A_j$ satisfying

$$U_{\mathbf{P}_j} \leq \frac{K \log(r_j^d)}{r_j^d} =: W_j.$$

It remains to show that

$$|\mathbf{P}_j - \mathbf{P}_{j-1}| \leq \min(U_{\mathbf{P}_j}^{-p}, U_{\mathbf{P}_{j-1}}^{-p}) \tag{3.5}$$

for all $j, 1 \leq j \leq m$. Lemma 3.3 shows (3.5) for $j = 1$. The maximal distance between points in A_j and A_{j-1} is at most twice r_{j-1} , i.e. $|\mathbf{P}_j - \mathbf{P}_{j-1}| \leq 2r_{j-1}$. It thus suffices to show that

$$2r_{j-1} \leq \min(W_j^{-p}, W_{j-1}^{-p}) = W_j^{-p}, \tag{3.6}$$

which holds since $W_{j-1}^{-p} \geq W_j^{-p}$ for all $j, 1 \leq j \leq m$.

However, by Lemma 3.2,

$$W_j^{-p} = \frac{r_j^{pd}}{(Kd \log r_j)^p} = \frac{((r_{j-1})^{b^{-1+\epsilon}})^{pd}}{(Kdb^{-1+\epsilon} \log(r_{j-1}))^p}.$$

Thus, for all $j, 1 \leq j \leq m$,

$$\frac{W_j^{-p}}{r_{j-1}} = \frac{(r_{j-1})^{b^\epsilon - 1}}{(Kdb^{-1+\epsilon} \log(r_{j-1}))^p} \geq \frac{(r_m)^{b^\epsilon - 1}}{(Kdb^{-1+\epsilon} \log(r_m))^p},$$

since the r_j are decreasing. By definition of r_m and since $b^\epsilon - 1 > 0$, the last ratio clearly exceeds 2 for large $|\mathbf{x}|$, showing (3.6) and completing the proof of Lemma 3.4.

The next lemma shows that we may link \mathbf{P}_m and $\mathbf{0}'$ via a node $\mathbf{P}_{m+1} \in A_{m+1}$. Combined with Lemmas 3.2 and 3.3, this builds a path between \mathbf{P}_0 and $\mathbf{0}'$ which contains $m + 2$ edges.

Lemma 3.5. *There is an event E_{m+1} , depending only on $\{U_z\}_{z \in A_{m+1}}$, such that $P[E_{m+1}] = 1 - O(r_{m+1}^{-dK})$, and on $E_0 \cap E_m \cap E_{m+1}$ there is a point $\mathbf{P}_{m+1} \in A_{m+1}$ such that the edges $(\mathbf{P}_m, \mathbf{P}_{m+1})$ and $(\mathbf{P}_{m+1}, \mathbf{0}')$ both belong to $G_p(\mathbb{Z}^d)$.*

Proof. First, by definition of m and by Lemma 3.2 we have

$$(\log \log |\mathbf{x}|)^\beta \leq r_m^\beta = r_{m+1} < \log \log |\mathbf{x}|.$$

By Lemma 3.1, for large $|\mathbf{x}|$ there is an event E_{m+1} , with $P[E_{m+1}] = 1 - O(r_{m+1}^{-dK})$, which depends only on $\{U_z\}_{z \in A_{m+1}}$ and which implies the existence of a point $\mathbf{P}_{m+1} \in A_{m+1}$ with

$$U_{\mathbf{P}_{m+1}} \leq \frac{K \log(r_{m+1}^d)}{r_{m+1}^d} \leq \frac{K \log(\log \log |\mathbf{x}|)^d}{(\log \log |\mathbf{x}|)^{\beta d}} \leq \frac{K \log(\log \log |\mathbf{x}|)^d}{(\log \log |\mathbf{x}|)^{(pd)^\epsilon / p}}$$

since $\beta d = (pd)^\epsilon / p$. Since $(pd)^\epsilon > 1$, it follows that, for large $|\mathbf{x}|$, on E_{m+1} we have

$$U_{\mathbf{P}_{m+1}}^{-p} \geq 2 \log \log |\mathbf{x}|. \tag{3.7}$$

Following the arguments in the proof of Lemma 3.4 (with j equal to $m + 1$ there), we find that, on $E_m \cap E_{m+1}$, $(\mathbf{P}_m, \mathbf{P}_{m+1})$ is an edge in $G_p(\mathbb{Z}^d)$. Furthermore, on $E_0 \cap E_m \cap E_{m+1}$, the edge $(\mathbf{P}_{m+1}, \mathbf{0}')$ belongs to $G_p(\mathbb{Z}^d)$ if and only if

$$|\mathbf{0}' - \mathbf{P}_{m+1}| \leq \min(U_{\mathbf{0}'}^{-p}, U_{\mathbf{P}_{m+1}}^{-p}). \tag{3.8}$$

However,

$$|\mathbf{0}' - \mathbf{P}_{m+1}| \leq |\mathbf{0}' - \mathbf{0}| + |\mathbf{0} - \mathbf{P}_{m+1}| \leq \log \log |\mathbf{x}| + r_{m+1} \leq 2 \log \log |\mathbf{x}|,$$

showing that (3.8) follows, using (3.7) and (3.1).

The last lemma completes the proof of Theorem 1.2.

Lemma 3.6. *For all $x \in \mathbb{Z}^d$, there is an event $E(x)$, with $P[E(x)] = 1 - o(1)$, such that on $E(x)$ there exists a path joining P_0 to $\mathbf{0}$ with $4 + 2 \log \log |x|$ edges.*

Proof. Let $E(x) := E_0 \cap E_{x/2} \cap (\bigcap_{j=1}^{m+1} E_j)$. On $E(x)$ we have shown that there is a path, π , joining P_0 to $\mathbf{0}$ via the successive nodes $P_1, P_2, \dots, P_m, P_{m+1}, \mathbf{0}', \mathbf{0}$. The number of edges in π is bounded by $m + 2 + \lceil \log \log |x| \rceil$, where $\lceil \log \log |x| \rceil$ denotes an upper bound on the number of edges between $\mathbf{0}'$ and $\mathbf{0}$. Since ε is arbitrary in the definition of t , it follows that $m \leq \lceil \log \log |x| \rceil$. Thus, $\text{card}\{\pi\} \leq 4 + 2 \log \log |x|$.

Finally, we show that $P[E(x)] = 1 - o(1)$. For all $j, 1 \leq j \leq m + 1$, E_j depends only on $\{U_z\}_{z \in A_j}$ and, since the A_j are disjoint, the $\{E_j\}_{1 \leq j \leq m+1}$ are independent. Clearly, since E_0 depends on $\{U_z\}_{z \in \mathbb{Z}^{d-1} \times \mathbb{Z}^-}$, we have independence of $E_0, E_1, E_2, \dots, E_{m+1}$. Similarly, $E_{x/2}, E_0, E_1, E_2, \dots, E_{m+1}$ are independent.

By independence, we have

$$P[E(x)] = P\left[\bigcap_{j=1}^{m+1} E_j\right] P[E_0] P[E_x] P[E_{x/2}] = (1 - o(1))^3 \prod_{j=1}^{m+1} P[E_j].$$

Now, m is bounded by $C \log \log |x|$ and the definition of r_m shows, for large K , that $m r_{m+1}^{-dK} \rightarrow 0$ as $|x| \rightarrow \infty$. Since $1 - 2s \leq \exp(-s) \leq 1 - s/2$ for small, positive s , it follows that

$$\begin{aligned} \prod_{j=1}^{m+1} P[E_j] &= \prod_{j=1}^{m+1} (1 - O(r_j^{-dK})) \\ &\geq \exp\left(-C \sum_{j=1}^{m+1} r_j^{-dK}\right) \\ &\geq 1 - C \sum_{j=1}^{m+1} r_j^{-dK} \\ &\geq 1 - C \sum_{j=1}^{m+1} r_{m+1}^{-dK}. \end{aligned}$$

This yields $P[E(x)] = 1 - o(1)$, as desired, completing the proof of Lemma 3.6.

4. Proof of Theorem 1.3

Assume without loss of generality that n has even parity. Partition $[0, n]^d \cap \mathbb{Z}^d$ into $Q_1 := [0, \frac{1}{2}n] \times [0, n]^{d-1} \cap \mathbb{Z}^d$ and $Q_2 := (\frac{1}{2}n, n] \times [0, n]^{d-1} \cap \mathbb{Z}^d$. For all $k = 0, 1, 2, \dots, n/2$, write $Q_{1,k} := \{n/2 - k\} \times [0, n]^{d-1} \cap \mathbb{Z}^d$ and note that $Q_1 = \bigcup_{k=0}^{n/2} Q_{1,k}$.

The number of nodes in Q_1 whose balls of influence have nonempty intersection with Q_2 is

$$N := \sum_{k=0}^{n/2} \sum_{i \in Q_{1,k}} 1_{\{U_i^{-p} \geq k+1\}}.$$

If we remove these N nodes from Q_1 then $G_p(Q_1)$ and $G_p(Q_2)$ are disconnected, i.e. the graphs have no edges joining them. Moreover, as the number of nodes in $Q_{1,k}$ equals n^{d-1} , we obtain

$$E[N] = \sum_{k=0}^n n^{d-1} \mathbb{P}[U_0^{-p} \geq k+1] = n^{d-1} \sum_{k=0}^n (k+1)^{-1/p} \leq Cn^{d-1} [n^{1-1/p} \vee 1],$$

which is exactly the desired upper bound.

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