HAMILTONIAN CYCLES IN SQUARES OF VERTEX-UNICYCLIC GRAPHS

BY HERBERT FLEISCHNER AND ARTHUR M. HOBBS⁽¹⁾

In this paper we determine necessary and sufficient conditions for the square of a vertex-unicyclic graph to be Hamiltonian. The conditions are simple and easily checked. Further, we show that the square of a vertex-unicyclic graph is Hamiltonian if and only if it is vertex-pancyclic.

We use the terminology and notation of [1]. The square G^2 of a graph G is the graph with $V(G^2) = V(G)$ in which two vertices are adjacent if and only if their distance in G is one or two. A graph G is Hamiltonian if there is a cycle in G, called a Hamiltonian cycle, which includes all of the vertices of G. A graph G is vertex-pancyclic if, for every vertex v of G, there are cycles of every length from 3 through |V(G)| in G which include v. A cactus is a graph in which each edge is in at most one cycle, and a graph is vertex-unicyclic if each vertex is in at most one cycle. Every vertex-unicyclic graph is a cactus. For any i, $V_i(G)$ is the set of vertices of G of degree *i*.

Given a vertex v of a graph G, a v-fragment of G is any maximal connected subgraph of G in which v is not a cut vertex. Clearly, if G is connected and if vis not a cut vertex of G, the only v-fragment is G itself, while if G is connected and v is a cut vertex of G, then there are as many v-fragments of G as there are components of G-v, and containment specifies a one-to-one relation between the components of G-v and the v-fragments of G. (Note: vfragments are a specialization of the J-components of Tutte [4]). If v is in a cycle C of G, the \overline{C} , v-fragments of G are those v-fragments of G which do not include C.

Let S be a v-fragment of a graph G. We say S is a short fragment of G if and only if there is a Hamiltonian path p in $S^2 - v$ whose first and last vertices are both adjacent in S to v; we call p a short path for S. S is a long fragment of G if and only if S is not a short fragment and there is a Hamiltonian path q in $S^2 - v$ whose first vertex is adjacent in S to v and whose last vertex is at distance 2 in S from v; we call q a long path for S. If a v-fragment is neither

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were proved:

THEOREM A. Let G be a connected cactus with cycle $C = x_0, x_1, \ldots, x_n, x_0$. Then G^2 is Hamiltonian if and only if

(1) no \overline{C} , x_i -fragment of G is insufferable, for each $i \in \{0, 1, ..., n\}$;

(2) no more than two \overline{C} , x_i -fragments of G are long, for each $i \in \{0, 1, ..., n\}$; and

(3) if two distinct \overline{C} , x_i -fragments and two distinct \overline{C} , x_j -fragments of G are long, then each non-trivial trail in C beginning on x_i and ending on x_j includes a vertex whose degree in G is $2(x_i$ and x_j may be the same vertex).

THEOREM B. Suppose v-fragment S of a graph G is a tree. Then

(1) S is short if and only if |V(S)| = 2; and

(2) S is long if and only if $|V(S)| \ge 3$ and $S - (V_1(G) \cap V(S))$ is a path.

THEOREM C. Let G be a graph with exactly one cycle $C = x_0, x_1, \ldots, x_n, x_0$. Suppose G is connected, every vertex x of C meets at most two long \overline{C} , x-fragments and no insufferable fragments, and suppose any path in C which joins two vertices, both of which meet two long fragments, includes a vertex whose degree in G is two. Then G^2 is Hamiltonian.

Using these theorems and a few further definitions, we can now give a characterization of vertex-unicyclic graphs whose squares are Hamiltonan. Given sequences $a = r_1, \ldots, r_i, b = s_1, s_2, \ldots, s_j$, and $c = t_1, t_2, \ldots, t_k$, (a), (b), (c) is the sequence $r_1, r_2, \ldots, r_i, s_1, \ldots, s_j, t_1, \ldots, t_k$. We denote the sequence $r_i, r_{i-1}, \ldots, r_2, r_1$ by a^{-1} . Further, we say that a, b, and c are sections of the sequence (a), (b), (c). A bridge of a graph is an endbridge if it meets a vertex of degree 1.

THEOREM 1. Let G be a vertex-unicyclic graph with at least three vertices. Then G^2 is Hamiltonian if and only if

(1) G is connected;

(2) no vertex of G meets more than two non-end bridges of G; and

(3) if v is on a cycle C of G and if v meets two non-end bridges of G, and if P is a trail in C from v to any vertex of G meeting two non-end bridges of G (P may join v to v), then P includes a vertex whose degree in G is two.

Proof. The necessity of (1) is obvious and the necessity of (2) was shown in [2]. Since any x-fragment which includes a non-end bridge of G incident with x is either long or insufferable, the necessity of (3) follows from the third part of Theorem A.

We prove the sufficiency of these conditions by induction. If G has no cycles, then the theorem is a consequence of Theorem B. If G has just one cycle, then

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the theorem follows immediately from Theorem C. Suppose the theorem is true for any vertex-unicyclic graph with k cycles and suppose G has k+1cycles, $k+1 \ge 2$. Then G has a bridge x_1x_2 such that each component G'_i of $G-x_1x_2$ containing x_i has a cycle. Form G_i from G'_i by adding a path $P_i=x_i$, y_i , z_i of length two to G'_i with one end at x_i , such that $\{y_i, z_i\} \cap V(G) = \emptyset$. Since x_1x_2 is a non-end bridge of G, both graphs G_i satisfy the conditions of this theorem if G does, and each has no more than k cycles. Thus G_i^2 contains a Hamiltonian cycle h_i . Since z_i has degree one in G_i there is a vertex w_i of G'_i adjacent to x_i in G such that the sequence w_i, y_i, z_i, x_i is a section of h_i or h_i^{-1} (or of a rotation of one of these). We may suppose $h_i = w_i, y_i, z_i, x_i, (p_i), w_i$ for some sequence p_i . But w_1x_2 and w_2x_1 are edges of G^2 . Thus $x_1, (p_1), w_1, x_2, (p_2), w_2, x_1$ is a Hamiltonian cycle in G^2 . This proves the theorem.

Because of the relatively simple structure of a vertex-unicyclic graph, Theorem 1 can be improved very easily, as follows:

THEOREM 2. Let G be a vertex-unicyclic graph. Then G^2 is Hamiltonian if and only if G^2 is vertex-pancyclic.

Proof. The result is immediate from the definitions if G^2 is vertex-pancyclic. Suppose G^2 is Hamiltonian and suppose v is a vertex of G. Since the theorem is immediate from Theorem B if G is a tree, we may suppose the theorem holds for all vertex-unicyclic graphs smaller than G. By Theorem 1, no vertex of G meets more than two non-end bridges of G, and given any two vertices on a cycle, both of which meet two non-end bridges of G, and a trail joining the vertices, the trail includes a vertex whose degree in G is 2. If G has a vertex x of degree 1 other than v, these properties are also satisfied by G-x, which is also connected and vertex-unicyclic. Thus $(G-x)^2$ is vertex-pancyclic, so G^2 has cycles of every length from 3 through |V(G)| which include v.

Suppose G has no vertices of degree one other than (possibly) v. If G is a cycle, we can remove any edge from G to obtain a tree whose square is vertex-pancyclic. Otherwise, G includes a cycle which does not include v. Among all cut vertices x' of G and all x'-fragments F' not including v, choose a cut vertex x and x-fragment F which includes a smallest number of cycles. Because G is vertex-unicyclic, if F includes more than one cycle, it includes a bridge which meets a cut vertex y' for which a y'-fragment contained in F includes fewer cycles than F. Thus F includes just one cycle C. Let y be the only vertex of C which meets an edge not in C. Then y has degree three in G. Let z be a vertex of C adjacent to y. Then G-z is clearly a graph which satisfies the conditions of Theorem 1, includes v, and has fewer vertices than G. Thus $(G-z)^2$ includes cycles of all lengths from 3 through |V(G-z)| and containing v, so G^2 is vertex-pancyclic.

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Institut für Informationsverarbeitung Wien, Austria DEPARTMENT OF MATHEMATICS TEXAS A & M UNIVERSITY COLLEGE STATION, TX77843 U.S.A.