

HYPERMETRIC SPACES AND THE HAMMING CONE

DAVID AVIS

1. Definitions and preliminary results. We denote by $d = (d_{12}, \dots, d_{1n}, d_{23}, \dots, d_{n-1,n})$ a vector of $\binom{n}{2}$ distances between n points. Such a vector d is called a *metric* if it satisfies the triangle inequalities

$$(1) \quad d_{ij} + d_{jk} \geq d_{ik} \quad 1 \leq i, j, k \leq n.$$

The set of all metrics on n points forms a convex polyhedral cone, the extremal properties of which are discussed in [4]. We will be concerned with a sub-cone that is spanned by metrics of the form

$$(2) \quad d_{ij}(t) = \begin{cases} t & i \perp j \in V \\ 0 & \text{otherwise,} \end{cases} \quad 1 \leq i < j \leq n$$

where $t \geq 0$, V is a proper subset of $\{1, 2, \dots, n\}$ and the symbol \perp is used for "exclusive or": $i \perp j \in V$ means $i \in V, j \notin V$ or $i \notin V, j \in V$. The metrics (2) are extreme rays of the metric cone and are called *Hamming rays*. The convex hull of these rays is called the *Hamming cone* H_n and we call d *Hamming*, if $d \in H_n$. Such metrics are also called *L^1 -embeddable* (e.g., [2]) or *addressable* (e.g., [5]).

Let Ω be a finite set and let $\{A_i | 1 \leq i \leq n\}$ be a collection of n subsets of Ω that will be called *addresses*. Then it can be shown that a metric d is Hamming if and only if for some finite set Ω , there exist addresses A_i and non-negative weights w_j ($j \in \Omega$) so that

$$(3) \quad d_{ij} = \sum_{k \in A_i \Delta A_j} w_k,$$

where Δ denotes symmetric difference. In the case where the weights are binary valued, the metrics are just the usual Hamming metrics that appear in coding theory.

Let F_n denote the $(2^{n-1} - 1)$ -tuple of all proper subsets of $\{1, 2, \dots, n\}$ that contain the element 1. Then from the definition, d is Hamming if and only if the following primal problem has a solution:

$$P. \quad \begin{aligned} \sum_{\substack{S \in F_n \\ i \perp j \in S}} \lambda_S &= d_{ij} & 1 \leq i < j \leq n \\ \lambda_S &\geq 0 & S \in F_n. \end{aligned}$$

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This formulation has two consequences. First, if P has a solution it has a solution that uses at most $\binom{n}{2}$ non-zero variables. Thus we may assume that $|\Omega| \leq \binom{n}{2}$ and that the address lengths are similarly bounded. Second, an application of Farkas' Lemma (see [13]) states that P is feasible if and only if the following system is infeasible:

$$D. \quad \sum_{i \perp j \in S} y(i, j) \leq 0 \quad S \in F_n \quad (4) \\ y \cdot d > 0.$$

A vector satisfying (4) is called dual feasible. Let B denote the set of all facets (maximum dimensional faces) of H_n , so that

$$H_n = \{d \mid b \cdot d \leq 0, b \in B\}.$$

If z_S is the Hamming ray that corresponds to some $S \in F_n$, and $b \in B$, then $bz_S \leq 0$ reduces to

$$\sum_{i \perp j \in S} b(i, j) \leq 0.$$

Therefore dual feasible vectors are facets if $\binom{n}{2} - 1$ linearly independent constraints in (4) are satisfied as equations.

We now consider vectors y that have the form $y(i, j) = c_i c_j$ for suitably chosen real numbers c_1, c_2, \dots, c_n . In this case, if

$$\sum_{i=1}^n c_i = t,$$

then row (i, j) of (4) becomes

$$\sum_{i \perp j \in S} c_i c_j = \sum_{i \in S} c_i \left(\sum_{j \notin S} c_j \right) + \sum_{i \notin S} \left(\sum_{j \in S} c_j \right) \\ = 2 \left(\sum_{i \in S} c_i \right) \left(t - \sum_{i \in S} c_i \right) \leq 0.$$

This proves the following theorem.

THEOREM 1.1. *If there exist real numbers c_1, c_2, \dots, c_n with sum $t \geq 0$ satisfy*

$$(i) \quad \sum_{i \in S} c_i \leq 0 \text{ or } \sum_{i \in S} c_i \geq t \text{ for all } S \in F_n, \text{ and}$$

$$(ii) \quad \sum_{1 \leq i < j \leq n} c_i c_j d_{ij} > 0,$$

then $d \notin H_n$.

Kelly [11] calls a metric space, d , *hypermetric* if for all integers c_1, \dots, c_n which sum to one,

$$(5) \quad \sum_{1 \leq i < j \leq n} c_i c_j d_{ij} \leq 0.$$

The notion of hypermetricity, but not the term, seems to have been first introduced by Deza [8]. Theorem 1.1 has the following corollary, which has been proved independently by Deza [7], [8] and Kelly [11].

COROLLARY 1.2. *If $d \in H_n$ then d is hypermetric.*

Observe that the triangle inequality is obtained from (5) by setting $c_i = c_j = 1$ and $c_k = -1$ for $i, j, k \in \{1, 2, \dots, n\}$. This is the first of a series of inequalities, the next being the pentagon inequality, see [7] and [11], where three indices are set to $+1$ and two indices are set to -1 . The reader interested in a full treatment of hypermetric spaces is referred to [12]. In the next section we show that the converse of Corollary 1.2 is false in general and give some specific instances when it is true. We also exhibit a facet of H_n that does not have the form of Theorem 1.1.

2. Main results. This section deals with the converse of Corollary 1.2. First, Deza [7] has shown that every 5 point hypermetric is Hamming. Second, as we now demonstrate, a theorem of Djoković [9] can be used to prove that all hypermetric bi-partite graphs are Hamming under the normal shortest distance metric for graphs. Indeed, following Djoković, a subset V_0 of vertices of a bi-partite graph G is *closed* if for every a and b contained in V_0 and any vertex w satisfying

$$d_G(a, w) + d_G(w, b) = d_G(a, b)$$

we have $w \in V_0$. Here, d_G is the shortest distance metric induced by the graph G . For every edge ab , let $G(a, b)$ denote the set of points closer to a than b . Note that the fact that G is bipartite implies that if $w \notin G(a, b)$, then $w \in G(b, a)$, for any vertex w .

THEOREM 2.1. (Djoković [9], see also [5]) *For a connected bi-partite graph G , d_G is Hamming if and only if $G(a, b)$ is closed for adjacent vertices a and b .*

We now show that this theorem has the following corollary.

COROLLARY 2.2. *For a connected bipartite graph $G = (V, E)$, d_G is Hamming if and only if the pentagon inequality is satisfied.*

Proof. (\Rightarrow) Assume $d_G \notin H_n$. By Theorem 2.1 there exist adjacent vertices $a, b \in V$ with $G(a, b)$ not closed. Therefore there exists $u, v \in G(a, b)$ and $w \in G(b, a)$ with

$$d_G(u, w) + d_G(w, v) = d_G(u, v).$$

Now $u \neq a$, for otherwise

$$d_G(a, v) = d_G(a, w) + d_G(w, v) = 1 + d_G(b, w) + d_G(w, v) \geq 1 + d_G(b, v)$$

and so $v \notin G(a, b)$, a contradiction. Similarly $v \neq a$. Also $w \neq b$, for otherwise

$$d_G(u, v) = d_G(u, b) + d_G(b, v) > d_G(u, a) + d_G(a, v),$$

violating the triangle inequality. We must therefore have the situation in Figure 2.1. Let

$$c_i = \begin{cases} 1 & i = u, b, v \\ -1 & i = a, w \\ 0 & \text{otherwise.} \end{cases}$$

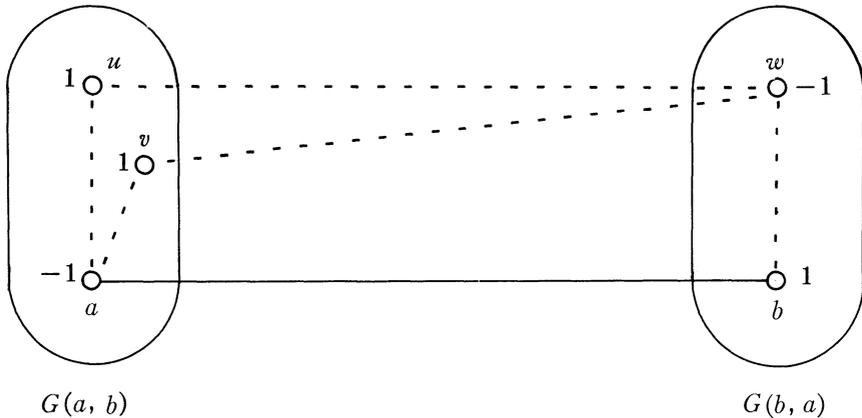


FIGURE 2.1

Note that $d_G(u, b) = 1 + d_G(u, a)$ since $u \in G(a, b)$. Using similar simplifications we have

$$\begin{aligned} \sum c_i c_j d_G(i, j) &= [3 + d_G(u, a) + d_G(v, a) + d_G(b, w)] \\ &\quad - [1 - d_G(u, a) + d_G(v, a) + d_G(b, w)] \\ &= 2 > 0. \end{aligned}$$

Therefore d_G violates the pentagon inequality and is not hypermetric.

Complete results for graphs in general are known only for $|G| \leq 6$. An examination of graphs with 6 or fewer vertices (listed in [10]) produced the five minimal nonhypermetric graphs shown in Figure 2.2. The integers attached to the vertices of the graphs in Figure 2.2 correspond to the integers c_1, c_2, \dots, c_n that form the coefficients of the hypermetric inequality (5) that is violated. All nonHamming graphs with 6 or fewer

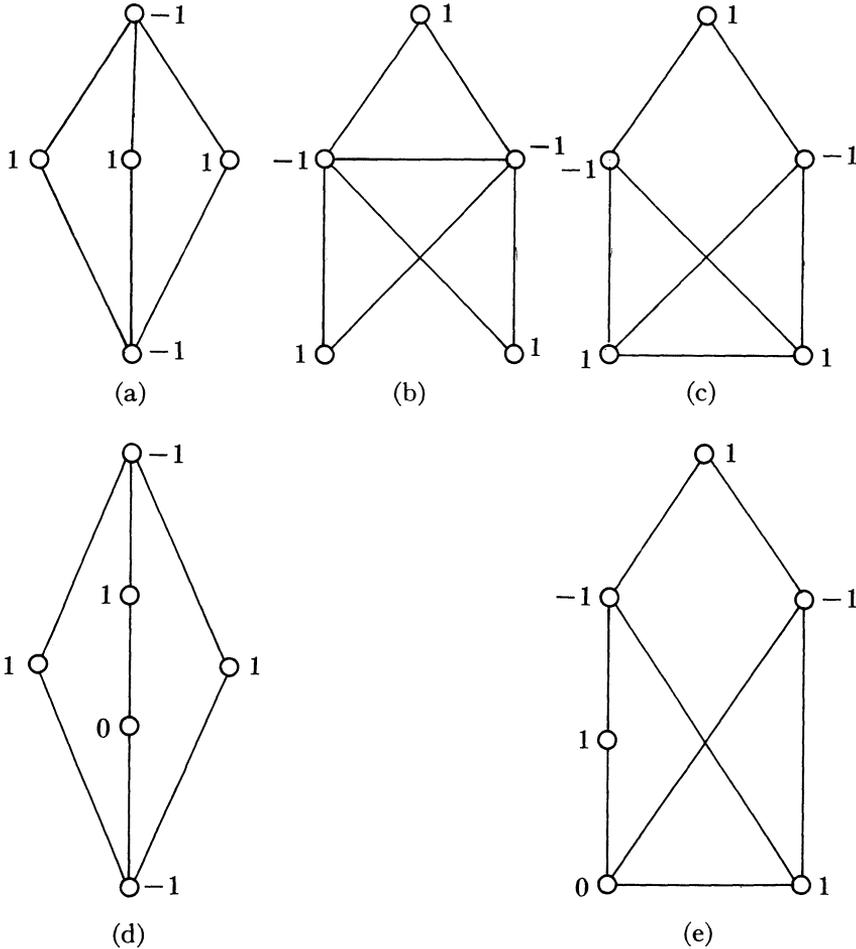


FIGURE 2.2. All minimal nonhypermetric graphs on 6 or fewer vertices

vertices contain an isometric nonhypermetric subgraph. Thus the converse of Corollary 1.2 is true for graphs G with $|G| \leq 6$. For $|G| \leq 5$, this also follows from results in [7].

We now show that the converse to Corollary 1.2 is false for $|G| \geq 7$, by exhibiting a nonHamming hypermetric graph. Consider the graph G formed from K_7 by deleting edges v_1v_2 and v_1v_3 .

THEOREM 2.3. G is a nonHamming hypermetric graph.

Proof. We begin by proving that d_G is hypermetric. Assume conversely that there exists integers c_1, c_2, \dots, c_7 such that

$$\sum_{i=1}^7 c_i = 1 \quad \text{and} \quad \sum_{1 \leq i < j \leq 7} c_i c_j d_G(i, j) \geq 1.$$

Using the identity

$$2 \sum_{1 \leq i < j \leq 7} c_i c_j = \left(\sum_{i=1}^7 c_i \right)^2 - \sum_{i=1}^7 c_i^2$$

and the structure of G we have

$$2 \sum_{1 \leq i < j \leq 7} c_i c_j d_G(i, j) = 2c_1(c_2 + c_3) + 1 - \sum_{i=1}^7 c_i^2.$$

Therefore c_1, \dots, c_7 satisfy

$$(6) \quad 2c_1(c_2 + c_3) \geq \sum_{i=1}^7 c_i^2.$$

Consider the function f defined over integer vectors $c = (c_1, \dots, c_7)$ by

$$f(c) = 2c_1(c_2 + c_3) - \sum_{i=1}^7 c_i^2.$$

Set $k = c_1 + c_2 + c_3$ and define g by

$$g(c_1, c_2, c_3) = 2c_1(c_2 + c_3) - c_1^2 - c_2^2 - c_3^2 - (1 - k)^2/4.$$

Then g bounds f from above for fixed c_1, c_2 and c_3 since it sets $c_4 = c_5 = c_6 = c_7$ (their optimal values). We are interested in the values of k that allow g to be nonnegative. To this end we seek to solve the quadratic program

$$\max g(c_1, c_2, c_3)$$

subject to

$$c_1 + c_2 + c_3 = k.$$

Inspection shows that, for fixed k , the maximum occurs for $c_2 = c_3$. Making the indicated substitutions,

$$g(c_1, c_2, c_3) = 8c_2k - 14c_2^2 - k^2 - (1 - k)^2/4.$$

The maximum (c_1^*, c_2^*, c_3^*) therefore occurs at $(3k/7, 2k/7, 2k/7)$ and

$$g(c_1^*, c_2^*, c_3^*) = k^2/7 - (1 - k)^2/4.$$

By inspection this maximum is negative outside of the range $1 \leq k \leq 4$. Therefore we need only look at values of k in this range to seek a solution to (6). For these values it is easy to compute the maximum of the left side and the minimum of the right side of (6) independently. The details are omitted. Therefore (6) is never satisfied and d_G is hypermetric.

To show that d_G is nonHamming we exhibit an inequality that is satisfied by all Hamming extreme rays, but is not satisfied by d_G . We define

$c \in R^{\binom{7}{2}}$ by

$$(7) \quad c(i, j) = \begin{cases} 5 & i = 1 & j = 2, 3 \\ -3 & i = 1 & j \geq 4 \\ 3 & i = 2 & j = 3 \\ -2 & i = 2, 3 & j \geq 4 \\ 1 & i \geq 4 & j \geq 5. \end{cases}$$

We now check that

$$(8) \quad \sum_{i, j \in T} c(i, j) \leq 0 \quad \text{for all } T \in F_n.$$

Recall that $1 \in T$ for every $T \in F_n$. The notation $[s, t]$ will refer to elements T of F_n that contain s members of $\{2, 3\}$ and t members of $\{4, 5, 6, 7\}$. The left side of (8) is given in Table 2.1.

Table 2.1

		t					LHS of (8)
		0	1	2	3	4	
s	0	-2	0	0	-2	-6	$-2 + 3t - t^2$
	1	-12	-6	-2	0	0	$-12 + 7t - t^2$
	2	-28	-18	-10	-4	-	$-28 + 11t - t^2$

Therefore (3) is satisfied for all $T \in F_n$ hence $cx \leq 0$ for all extreme rays x of H_n and hence for all $x \in H_n$. Since $c \cdot d_G = 1$, d_G is nonHamming.

It can be shown that the vector c defined in (7) is in fact a facet of H_7 . Inspection of Table 2.1 shows that equality holds in (3) for the subsets T corresponding to $[s, t] = [0, 1], [0, 2], [1, 3]$ and $[1, 4]$. Now there are 4 subsets of the first type, namely, $\{1, 4\}, \{1, 5\}, \{1, 6\}$ and $\{1, 7\}$. Similarly there are 6 of the second type, 8 of the third type and 2 of the fourth for a total of $20 = \binom{7}{2} - 1$ subsets. The corresponding 20 extreme rays can be shown to be independent.

Thus the facets of the Hamming cone are not all of the type $b(i, j) = c_i c_j$ ($1 \leq i < j \leq n$). This answers a question posed by Deza, who has given a list of facets of this type for $n \leq 8$ [6].

3. Concluding remarks. The results of this paper first appeared in the author's Ph.D. thesis that is reprinted in part as [3]. Independently, P. Assouad [1] has shown that the Hamming cone is properly contained in the hypermetric cone using a different proof based on the corresponding dual cones. Many related topics may be found in the monograph in

preparation by Assouad and Deza [2] that contains a survey of all known results on embeddability in L^1 .

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*McGill University,
Montreal, Quebec*