## A NOTE ON ROOT DEGISION PROBLEMS IN GROUPS

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1. Introduction. Consider a positive integer $r>1$. We say that the $r$ th root problem is solvable for a group $G$ if we can decide for any $W \in G$ whether or not $W$ has an $r$ th root, i.e. whether or not there exists $V \in G$ such that $W=V^{r}$.

Baumslag, Boone and Neumann [1] proved that there exists a finitely presented group with all root problems unsolvable. Here we are concerned with the relationship between the different root problems. We prove:

Theorem 1. Let $r$ and $s$ be positive integers such that neither divides the other. Then the corresponding root problems are independent.

Theorem 2. Let $r_{1}, r_{2}, \ldots, r_{n}$ be positive integers such that each has a prime divisor which does not divide any of the others. Then the corresponding root problems are independent.

Recall that a group $G$ with generators $x_{1}, x_{2}, \ldots$ and defining relations $R_{1}, R_{2}, \ldots$ is said to be recursively presented if there exists an effective process which lists the words $R_{n}$. We say that the decision problems $D_{1}, D_{2}, \ldots, D_{n}$ for groups are independent if for any subset $\left\{i_{1}, \ldots, i_{m}\right\}$ of $\{1, \ldots, n\}$ there exists a recursively presented group with $D_{i 1}, \ldots, D_{i_{m}}$ solvable, but the remaining $D_{i}$ 's unsolvable.

Our proofs use a one-to-one recursive function $\phi: \mathbf{N} \rightarrow \mathbf{N}$ whose image, $\operatorname{Im} \phi$, is non-recursive, i.e. given a positive integer $k$, we can compute $\phi(k)$ but we cannot decide if $k$ belongs to $\operatorname{Im} \phi$. Such functions are known to exist; c.f. Britton [2, Lemma 2.31]. We will also use the following theorem which follows from elementary properties of free products.

Theorem A. Let $G$ be the free product of groups $A_{i}$. Then the $r$ th root problem is solvable for $G$ if the word problem and $r$ th root problem are solvable for the factors $A_{i}$.
2. Proof of Theorem 2. Let $p$ be a prime and let $\phi: \mathbf{N} \rightarrow \mathbf{N}$ be a one-to-one recursive function with a non-recursive image (see introduction). Let $G_{p}$ be the group with generators

$$
x_{1}, x_{2}, x_{3}, \ldots \text { and } y_{1}, y_{2}, y_{3}, \ldots
$$

[^0]and defining relations
$$
x_{\phi(1)}=y_{1}^{p}, \quad x_{\phi(2)}=y_{2}^{p}, \quad x_{\phi(3)}=y_{3}^{p}, \ldots
$$

Clearly $G_{p}$ is recursively presented since $\phi$ is recursive. We claim:
Lemma. The word problem is solvable for $G_{p}$. The $r$ th root problem is unsolvable for $G_{p}$ if $p$ divides $r$, and solvable if $p$ does not divide $r$.

Proof of Lemma. Observe that $G_{p}$ is the free product of the infinite cyclic groups generated by the $y_{i}$ and the infinite cyclic groups generated by the $x_{j}$ for $j \notin \operatorname{Im} \phi$. Since $\phi$ is recursive, we can always write $W \in G_{p}$ as a word in the $x$ 's and $y$ 's so that $y_{i}$ does not appear next to $x_{j}$ whenever $\phi(i)=j$. This gives the syllable (free product) length of $W$. Thus we can solve the word problem for $G_{p}$.

Suppose $p$ divides $r$, say $r=p t$. Then $x_{i}{ }^{t}$ has an $r$ th root if and only if $i \in \operatorname{Im} \phi$; if $\phi(j)=i$ then $y_{j}$ is an $r$ th root of $x_{i}{ }^{t}$. But $\operatorname{Im} \phi$ is non-recursive; hence the $r$ th root problem is unsolvable for $G_{p}$ if $p$ divides $r$.

On the other hand, suppose $p$ does not divide $r$. Since the word problem is solvable for $G_{p}$ and since $G_{p}$ is a free product, it suffices to solve the $r$ th root problem for a factor $H$ of $G_{p}$. Let $V \in H$. Then $V=y_{i}{ }^{m}$ or $V=x_{i}{ }^{m}$. We claim that $V$ has an $r$ th root if and only if $r$ divides $m$. This is clearly true in the case that $V=y_{i}{ }^{m}$, or $V=x_{i}{ }^{m}$ for $i \notin \operatorname{Im} \phi$. Suppose $\phi(j)=i$; then $V=x_{i}{ }^{m}=$ $y_{j}{ }^{m p}$. But $p$ does not divide $r$; hence $V$ has an $r$ th root if and only if $r$ divides $m$. Thus the $r$ th root problem is solvable for $H$, and hence it is solvable for $G_{p}$. Accordingly, the lemma is proved.

We now prove Theorem 2. Let $\left\{i_{1}, \ldots, i_{m}\right\}$ be a subset of $\{1,2, \ldots, n\}$ and let $G$ be the direct product

$$
G=G_{z 1} \times G_{z 2} \times \ldots \times G_{z_{m}}, \quad z_{j}=p_{i j}
$$

By the above lemma, the $r_{i}$ th root problem is solvable for $G$ if and only if $i \notin\left\{i_{1}, \ldots, i_{m}\right\}$. Thus, Theorem 2 is proved.
3. Proof of Theorem 1. Let $d=\operatorname{gcd}(r, s)$; say $r=d a$ and $s=d b$. Then $\operatorname{gcd}(a, b)=1$; also $a \neq 1$ and $b \neq 1$ since neither $r$ nor $s$ divides the other. Let $G$ be the group with generators

$$
x_{1}, x_{2}, x_{3}, \ldots \quad \text { and } \quad y_{1}, y_{2}, y_{3}, \ldots
$$

and defining relations

$$
x_{\phi(1)}{ }^{d}=y_{1}{ }^{\tau}, \quad x_{\phi(2)}{ }^{d}=y_{2}{ }^{\tau}, \quad x_{\phi(3)^{d}}{ }^{d}=y_{3}{ }^{r}, \ldots
$$

Clearly $G$ is recursively presented since $\phi$ is recursive. We claim that the $s$ th root problem is solvable for $G$, but the $r$ th root problem is not. This will prove our theorem since we can interchange $r$ and $s$ to obtain a group for which the $r$ th root problem is solvable but the $s$ th root problem is not.

Note that $G$ is the free product of the groups

$$
H_{\imath}=\left\langle x_{\phi(i)}, y_{i} ; x_{\phi(i)}{ }^{d}=y_{i}{ }^{r}\right\rangle, \text { and } K_{j}=\left\langle x_{j}\right\rangle, j \notin \operatorname{Im} \phi .
$$

Observe that $x_{k}{ }^{d}$ has an $r$ th root if and only if $k \in \operatorname{Im} \phi$; hence the $r$ th root problem is unsolvable for $G$ because $\operatorname{Im} \phi$ is non-recursive. It remains to show that the $s$ th root problem is solvable for $G$. We claim, first of all, that the word problem is solvable for $G$. Let $W \in G$, i.e. let $W$ be a word in the $x$ 's and $y$ 's. Now if a $y_{i}$ appears next to an $x_{j}$ in $W$, then we can decide whether or not they belong to the same factor of $G$ because we can decide whether or not $\phi(i)=j$. Moreover, we can solve the word problem for the factors of $G$. Thus we can determine the syllable length of $W$, and hence solve the word problem for $G$. Accordingly, it suffices to solve the $s$ th problem for the factors of $G$. That is, given an element $V$ in a factor of $G$, we have to decide whether or not $V$ has an $s$ th root. There are two cases.

Case I. $V$ is a power of $y_{i}$, or a word in $Y_{i}$ and $x_{j}$ with $\phi(i)=j$. Then $V$ belongs to $H_{i}$. How $H_{i}$ is a free product of two infinite cyclic groups with a cyclic amalgamation. The sth root problem is solvable for such a group; cf. [3]. Thus we can decide whether or not $V$ has an $s$ th root.

Case II. $V$ is a power of $x_{i}$; say $V=x_{i}{ }^{n}$. We claim that $V$ has an $s$ th root if and only if $s$ divides $n$. If $i \notin \operatorname{Im} \phi$ then $V$ would belong to the infinite cyclic group $K_{i}$ generated by $x_{i}$, and the claim clearly holds. On the other hand, suppose $i \in \operatorname{Im} \phi ;$ say $\phi(k)=i$. Then $V$ belongs to the group

$$
H_{k}=\left\langle x_{\imath}, y_{k} ; x_{\imath}{ }^{d}=y_{k}{ }^{\tau}\right\rangle
$$

We view $H_{k}$ as a free product with an amalgamation. If $s$ divides $n$ then clearly $V$ has an sth root. Suppose, however, that $V$ has an sth root. We claim that one such sth root $U$ has syllable length one. If $d$ does not divide $n$, then $U$ must have syllable length one. If $d$ does divide $n$, then $V$ belongs to the center of $H_{k}$. It follows that $V$ has an $s$ th root $U$ which is cyclically reduced. This $s$ th root $U$ has syllable length one.

We now have that $V=U^{s}$ where $U$ has syllable length one. There are two possibilities:

Case A. $U=x_{i}{ }^{c}$. Then $x_{2}{ }^{c s}=U^{s}=V=x_{i}{ }^{n}$, whence $s$ divides $n$.
Case B. $U=y_{k}{ }^{e}$. Then $y_{k}{ }^{e s}=U^{s}=V=x_{i}{ }^{n}$. Then $V$ lies in the amalgamated subgroup, whence $y_{k}{ }^{e s}$ is a power of $y_{k}{ }^{r}$; say

$$
e s=r f
$$

Recall $r=a d$ and $s=b d$ where $\operatorname{gcd}(a, b)=1$. It follows that $b$ divides $f$; say $f=b g$. Using the relation $x_{i}{ }^{d}=y_{k}{ }^{r}$, we have

$$
V=U^{s}=y_{k}^{e s}=y_{k}{ }^{\tau f}=x_{i}{ }^{d f}=x_{i}{ }^{d b g}=x_{i}{ }^{s g} .
$$

But $V=x_{i}{ }^{n}$; hence $s$ divides $n$.
We have shown in Case II that $V=x_{i}{ }^{n}$ has an $s$ th root if and only if $s$ divides $n$. Thus we can decide whether or not $V$ has an $s$ th root.

Accordingly, the sth root problem is solvable for $G$, and therefore the theorem is proved.

## References

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2. J. L. Britton, Solution of the word problem for certain types of groups. I, Glasgow Math. J. 3 (1956), 45-54.
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[^0]:    Received February 25, 1972 and in revised form, May 2, 1972.

