

THE AND/OR THEOREM FOR PERCEPTRONS

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Abstract

Minsky and Papert claim that, for any positive integer n , there exist predicates of order 1 whose conjunction and disjunction have order greater than n . Their proof is amended and a stronger result obtained of which their claim is a special case.

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1. Introduction

The *perceptron* is a simple parallel computing device, and its capabilities and limitations have been studied by Minsky and Papert [2]. They claim [2, Section 1.5] to have shown that for any positive integer n there exist predicates ψ_1 and ψ_2 of order 1 for which both $\psi_1 \wedge \psi_2$ and $\psi_1 \vee \psi_2$ have order greater than n . This is called the AND/OR Theorem. Their argument [2, Chapter 4] establishes the existence of predicates ψ_1 and ψ_2 of order 1 for which $\psi_1 \wedge \psi_2$ has order greater than n . A similar argument establishes the existence of predicates ψ'_1 and ψ'_2 of order 1 for which $\psi'_1 \vee \psi'_2$ has order greater than n . However the arguments do not seem to guarantee that $\psi_1 = \psi'_1$ and $\psi_2 = \psi'_2$.

We amend their proof by adapting their techniques to prove a more general result, of which the AND/OR Theorem is a special case.

2. Preliminaries

Let R be a finite set. A *predicate on R* is a function ϕ of subsets of R whose values, depending on the context, may be thought of as TRUE and FALSE, or 1 and 0. To pass conveniently between these two kinds of predicate values, square bracket notation is used: if $\phi(X)$ is a statement about X then

$$[\phi(X)] = \begin{cases} 1 & \text{if } \phi(X) \text{ is true.} \\ 0 & \text{if } \phi(X) \text{ is false.} \end{cases}$$

Let Φ be a family of predicates on R , which is finite since R is finite. A predicate ψ is called a *linear threshold function* with respect to Φ if there are numbers θ , and α_ϕ for each $\phi \in \Phi$ such that

$$\psi(X) = \left[\sum_{\phi \in \Phi} \alpha_\phi \phi(X) > \theta \right].$$

Let $L(\Phi)$ denote the set of linear threshold functions with respect to Φ .

The *support* $S(\phi)$ of a predicate ϕ is the smallest subset S of R for which

$$(*) \quad (\forall X \subseteq R) \phi(X) = \phi(X \cap S)$$

The support exists because (*) is satisfied when $S = R$ and if any (finite) collection of subsets S satisfies (*) then their intersection also satisfies (*). Note however that if R is allowed to be infinite then the notion of support need not make sense.)

The *order* of a predicate ψ is the smallest number k for which there is a set Φ of predicates for which $\psi \in L(\Phi)$ and $|S(\phi)| \leq k$ for all $\phi \in \Phi$.

If $A \subseteq R$ then the *mask of A* , denoted μ_A , is the predicate defined by

$$\mu_A(X) = \begin{cases} 1 & \text{if } A \subseteq X \\ 0 & \text{otherwise.} \end{cases}$$

It is easy to see that $S(\mu_A) = A$. If $A = \{a\}$ denote μ_A also by μ_a .

THEOREM 1 [2, 1.5.3]. *A predicate ψ has order k if and only if k is the smallest number for which there exists a set Φ of masks such that $\psi \in L(\Phi)$ and $|S(\phi)| \leq k$ for all $\phi \in \Phi$.*

Denote the (symmetric) group of all permutations of a set X by $\text{Sym}(X)$. Let Γ be a subgroup of $\text{Sym}(R)$. If ψ is a predicate on R and $\gamma \in \Gamma$, then define the predicate $\psi\gamma$ by

$$\psi\gamma(X) = \psi(\gamma(X)) \quad \text{for } X \subseteq R.$$

Call predicates ψ and ϕ Γ -equivalent, written $\psi \equiv \phi$, if $\psi = \phi\gamma$ for some $\gamma \in \Gamma$. The Relation \equiv is an equivalence because Γ is a group. Note that, for masks, if $A, B \subseteq R$ then $\mu_A\gamma = \mu_{\gamma^{-1}(A)}$ for any $\gamma \in \Gamma$ so that

$$\mu_A \equiv \mu_B \text{ if and only if } A = B\gamma \text{ for some } \gamma \in \Gamma.$$

Call a predicate ψ invariant under Γ or Γ -invariant if

$$\psi = \psi\gamma \text{ for all } \gamma \in \Gamma.$$

THEOREM 2 (Group Invariance Theorem) [2, 2.3]. *Let R be a finite set, Γ a subgroup of $\text{Sym}(R)$ and Φ a set predicates on R closed under Γ containing μ_\emptyset the mask of the empty set. Suppose that $\psi \in L(\Phi)$ and that ψ is Γ -invariant. Then*

$$\psi = \left[\sum_{\phi \in \Phi} \beta_\phi \phi > 0 \right]$$

where the coefficients β_ϕ depend only on the Γ -equivalence class of ϕ , that is, if $\phi \equiv \phi'$ then $\beta_\phi = \beta_{\phi'}$.

Note (for when we apply this theorem later) that if Φ is the set of masks of support size $\leq k$ then Φ is closed under Γ and contains μ_\emptyset .

Let ψ_1, \dots, ψ_m be a sequence of predicates or formulae. Define a collection $\text{AND/OR}(\psi_1, \dots, \psi_m)$ of predicates or formulae inductively:

$$\text{AND/OR}(\psi_1) = \{\psi_1\},$$

$$\text{AND/OR}(\psi_1, \dots, \psi_{i+1}) = \{\psi \vee \psi_{i+1}, \psi \wedge \psi_{i+1} \mid \psi \in \text{AND/OR}(\psi_1, \dots, \psi_i)\} \text{ for } i = 1, \dots, m - 1.$$

Thus $\text{AND/OR}(\psi_1, \dots, \psi_m)$ contains 2^{m-1} predicates or formulae, obtained from ψ_1 by successively conjoining or disjoining ψ_i 's. For example $\text{AND/OR}(\psi_1, \psi_2) = \{\psi_1 \vee \psi_2, \psi_1 \wedge \psi_2\}$, and

$$\text{AND/OR}(\psi_1, \psi_2, \psi_3) = \{\psi \vee \psi_2 \vee \psi_3, (\psi_1 \vee \psi_2) \wedge \psi_3, \psi_1 \wedge \psi_2 \wedge \psi_3, (\psi_1 \wedge \psi_2) \vee \psi_3\}.$$

The following is a trivial but important observation: if $\psi_i(X) = \{0^1\}$ for all $i = 1, \dots, m$ then $\psi(X) = \{0^1\}$ for all ψ in $\text{AND/OR}(\psi_1, \dots, \psi_m)$.

The last result we shall need is an adaptation of [2, Lemma 1].

THEOREM 3 (Compactness). *Let $Q_1(x_1, \dots, x_m), Q_2(x_1, \dots, x_m), \dots$ be an infinite sequence of nonzero polynomials of m variables of degree at most N , and let n_1, n_2, \dots be an infinite increasing sequence of positive*

integers. Suppose $\phi: \mathbb{Z}^m \rightarrow \{0, 1\}$ is a function (predicate) such that for each $i = 1, 2, \dots$ we have for all integers x_1, \dots, x_m between $-n_i$ and n_i

$$\phi(x_1, \dots, x_m) = [Q_i(x_1, \dots, x_m) > 0].$$

Then there exists a nonzero polynomial $Q(x_1, \dots, x_m)$ of degree $\leq N$ such that for all integers x_1, \dots, x_m

$$\phi(x_1, \dots, x_m) = \begin{cases} 1 \\ 0 \end{cases} \text{ implies } Q(x_1, \dots, x_m) \begin{cases} \geq 0 \\ \leq 0. \end{cases}$$

PROOF. Let \mathbf{x} be a vector of coordinates ranging over all products of powers of x_1, \dots, x_m in which for each product the sum of the exponents is at most N . Polynomials in x_1, \dots, x_m of degree $\leq N$ are then dot products $\mathbf{x} \cdot \mathbf{c}$ where \mathbf{c} is a vector of constant coefficients. Thus for each i

$$Q_i(x_1, \dots, x_m) = \mathbf{x} \cdot \mathbf{c}_i \text{ for some } \mathbf{c}_i.$$

But the set $\{\hat{\mathbf{c}}_i = \mathbf{c}_i / \|\mathbf{c}_i\| \mid i = 1, 2, \dots\}$ lies on the surface of the unit hypersphere which is compact [1], so has a limit point \mathbf{c} of length 1. In particular $\mathbf{c} \neq \mathbf{0}$. Let $Q(x_1, \dots, x_m) = \mathbf{x} \cdot \mathbf{c}$, so Q is nonzero of degree $\leq N$. Let x_1, \dots, x_m be any integers. Choose n_i larger than each of $|x_1|, \dots, |x_m|$. Suppose $\phi(x_1, \dots, x_m) = 1$, so by our hypothesis $Q_j(x_1, \dots, x_m) > 0$ for $j \geq i$. Thus $\mathbf{x} \cdot \hat{\mathbf{c}}_j > 0$, so $\mathbf{x} \cdot \hat{\mathbf{c}}_j > 0$ for $j \geq i$. If $\mathbf{x} \cdot \mathbf{c} < 0$ then choose $j \geq i$ such that

$$\|\hat{\mathbf{c}}_j - \mathbf{c}\| < \frac{|\mathbf{c} \cdot \mathbf{c}|}{2\|\mathbf{x}\|}$$

and so

$$\begin{aligned} \mathbf{x} \cdot \hat{\mathbf{c}}_j &= \mathbf{x} \cdot \mathbf{c} + \mathbf{x} \cdot (\hat{\mathbf{c}}_j - \mathbf{c}) \leq \mathbf{x} \cdot \mathbf{c} + \|\mathbf{x}\| \|\hat{\mathbf{c}}_j - \mathbf{c}\| \\ &\leq \mathbf{x} \cdot \mathbf{c} + \|\mathbf{x}\| \|\hat{\mathbf{c}}_j - \mathbf{c}\| \quad (\text{by the Cauchy-Schwarz inequality}) \\ &< \mathbf{x} \cdot \mathbf{c} + \frac{\|\mathbf{x}\| \|\mathbf{c}\|}{2} = \frac{\mathbf{x} \cdot \mathbf{c}}{2} < 0, \end{aligned}$$

a contradiction. Hence $\mathbf{x} \cdot \mathbf{c} \geq 0$.

Similarly, if $\phi(x_1, \dots, x_m) = 0$ then $Q_j(x_1, \dots, x_m) \leq 0$ for $j \geq i$, and in the same way one shows that $\mathbf{x} \cdot \mathbf{c} \leq 0$. This completes the proof of Theorem 3.

3. The main theorem

The purpose of this section is to prove

THEOREM 4. *Let N and m be any positive integers, $m \geq 2$. There exist predicates ψ_1, \dots, ψ_m of order 1 such that every predicate in*

$$\text{AND/OR}(\psi_1, \dots, \psi_m)$$

has order greater than N .

Then Minsky and Papert's claim follows.

COROLLARY 5 (the AND/OR theorem). *Let N be any positive integer. There exist predicates ψ_1 and ψ_2 of order 1 such that both $\psi_1 \wedge \psi_2$ and $\psi_1 \vee \psi_2$ have order greater than N .*

PROOF OF THEOREM 4. To prove Theorem 4, let n be a positive integer and R any set containing $2nm$ elements. Express R as the disjoint union

$$R = A_1 \cup \dots \cup A_m \quad \text{where } |A_i| = 2n \text{ for each } i.$$

Define the predicate ψ_i , for $i = 1$ to m , by

$$\psi_i(X) = [|X \cap A_i| \geq n] \quad \text{for } X \subseteq R.$$

Note that each ψ_i has order 1 because

$$\psi_i(X) = \left[\sum_{a \in A_i} \mu_a(X) \geq n \right].$$

Thus for each n we have defined a sequence ψ_1, \dots, ψ_m of order 1 predicates. We will show that for some n each member of $\text{AND/OR}(\psi_1, \dots, \psi_m)$ has order larger than N .

Suppose to the contrary that for each n there is a predicate in

$$\text{AND/OR}(\psi_1, \dots, \psi_m)$$

of order $\leq N$. There are infinitely many n and only 2^{m-1} formulae in $\text{AND/OR}(\psi_1, \dots, \psi_m)$, so there exists at least one formula ψ in

$$\text{AND/OR}(\psi_1, \dots, \psi_m)$$

such that for some infinite sequence of positive integers $n_1 < n_2 < \dots$ the predicates ψ for $n = n_1, n_2, \dots$ have order $\leq N$.

For the time being fix $n = n_j$. Let Φ be the set of masks of support size $\leq N$. Then $\psi \in L(\Phi)$ by Theorem 1. Consider the group

$$\Gamma = \{\gamma \in \text{Sym}(R) \mid \gamma(A_i) = A_i \text{ for } i = 1, \dots, m\}.$$

Then each predicate ψ_i is Γ -invariant since

$$|X \cap A_i| = |\gamma(X \cap A_i)| = |\gamma(X) \cap \gamma(A_i)| = |\gamma(X) \cap A_i|$$

for each $\gamma \in \Gamma$ and $X \subseteq R$. By a simple induction all members of

$$\text{AND/OR}(\psi_1, \dots, \psi_m)$$

are Γ -invariant. Hence ψ is Γ -invariant, so by Theorem 2,

$$\psi = \left[\sum_{\phi \in \Phi} \beta_\phi \phi > 0 \right]$$

for some coefficients β_ϕ which depend only on the Γ -equivalence class of ϕ . Let

$$V = \{\mathbf{v} = (v_1, \dots, v_m) \in \mathbb{Z}^m \mid 0 \leq v_i \leq 2n \text{ for } i = 1 \text{ to } m \text{ and } v_1 + \dots + v_m \leq N\}.$$

Then

$$\Phi = \bigcup_{\mathbf{v} \in V} \Phi_{\mathbf{v}}$$

where $\Phi_{\mathbf{v}} = \{\text{masks } \mu_A \mid A \cap A_i = v_i \text{ for } i = 1 \text{ to } m\}$. Since $\Gamma|_{A_i} = \text{Sym}(A_i)$, the $\Phi_{\mathbf{v}}$ are the equivalence classes of Φ under Γ as \mathbf{v} ranges over V . Hence

$$\psi(X) = \left[\sum_{\phi \in \Phi} \beta_\phi \phi(X) > 0 \right] = \left[\sum_{\mathbf{v} \in V} \beta_{\mathbf{v}} N_{\mathbf{v}}(X) > 0 \right]$$

where $\beta_{\mathbf{v}} = \beta_\phi$ for any $\phi \in \Phi_{\mathbf{v}}$ and $N_{\mathbf{v}}(X)$ is the number of masks in $\phi_{\mathbf{v}}$ whose support is contained in X . But then

$$N_{\mathbf{v}}(X) = \binom{|X \cap A_1|}{v_1} \dots \binom{|X \cap A_m|}{v_m}$$

where $\binom{m}{r}$ is by definition $m(m-1)\dots(m-r+1)/r!$ so that $\binom{m}{r} = 0$ if $m < r$.

Hence each $N_{\mathbf{v}}(X)$ is a polynomial in the m variables $y_1 = |X \cap A_1|, \dots, y_m = |X \cap A_m|$ of degree $v_1 + \dots + v_m \leq N$. But a linear combination of polynomials of degree $\leq N$ is also a polynomial of degree $\leq N$, so

$$\psi(X) = [P(y_1, \dots, y_m) > 0]$$

for some polynomial $P(y_1, \dots, y_m)$ of degree $\leq N$. Note that $0 \leq y_i \leq 2n$ since $|A_i| = 2n$ for each i . Put $x_i = y_i - n$, so that $-n \leq x_i \leq n$ for each i . Now put

$$Q_j(x_1, \dots, x_m) = P(x_1 + n, \dots, x_m + n)$$

which is a polynomial in x_1, \dots, x_m of degree $\leq N$. The subscript j is to remind us that $n = n_j$. Thus

$$\psi(X) = [Q_j(x_1, \dots, x_m) > 0]$$

where $-n_j \leq x_1, \dots, x_m \leq n_j$.

Define functions $\phi_i: \mathbb{Z}^m \rightarrow \{0, 1\}$ for $i = 1$ to m by

$$\phi_i(z_1, \dots, z_m) = \begin{cases} 1 & \text{if } z_i \geq 0 \\ 0 & \text{if } z_i < 0. \end{cases}$$

Then $\psi(X) = \phi_i(x_1, \dots, x_m)$ where as before $x_i = |X \cap A_i| - n$ for $i = 1$ to m . Let $\phi \in \text{AND/OR}(\phi_1, \dots, \phi_m)$ be built using \wedge and \vee in exactly the same way as ψ . Then $\phi: \mathbb{Z}^m \rightarrow \{0, 1\}$ and for the x_i defined above $\phi(x_1, \dots, x_m) = \psi(X)$. Thus

$$\phi(x_1, \dots, x_m) = [Q_j(x_1, \dots, x_m) > 0]$$

where $-n_j \leq x_1, \dots, x_m \leq n_j$.

This holds for each j , so by Theorem 3 there exists a nonzero polynomial $Q(x_1, \dots, x_m)$ of degree $\leq N$ such that

$$\phi(x_1, \dots, x_m) = \begin{cases} 1 \\ 0 \end{cases} \text{ implies } Q(x_1, \dots, x_m) \begin{cases} \geq 0 \\ \leq 0 \end{cases}$$

for all integers x_1, \dots, x_m .

Put $Q(x_1, \dots, x_m) = x_m^d q(x_1, \dots, x_{m-1}) + r(x_1, \dots, x_m)$ where $q(x_1, \dots, x_{m-1})$ is a nonzero polynomial in x_1, \dots, x_{m-1} and $r(x_1, \dots, x_m)$ is either the zero polynomial or a nonzero polynomial such that the highest power of x_m appearing is less than d . Note that $\phi = \phi' \wedge \phi_m$ or $\phi = \phi' \vee \phi_m$ for some $\phi' \in \text{AND/OR}(\phi_1, \dots, \phi_{m-1})$, and that for all x_m

$$\phi'(x_1, \dots, x_m) = \begin{cases} 1 & \text{if } x_1, \dots, x_{m-1} \geq 0 \\ 0 & \text{if } x_1, \dots, x_{m-1} < 0. \end{cases}$$

If $\phi = \phi' \wedge \phi_m$ choose $x_1, \dots, x_{m-1} \geq 0$, whilst if $\phi = \phi' \vee \phi_m$ choose $x_1, \dots, x_{m-1} < 0$ for which $q(x_1, \dots, x_{m-1}) \neq 0$. In both cases

$$\phi(x_1, \dots, x_m) = \begin{cases} 1 & \text{if } x_m \geq 0 \\ 0 & \text{if } x_m < 0 \end{cases}$$

so that

$$Q(x_1, \dots, x_m) = \begin{cases} \geq 0 & \text{if } x_m \geq 0 \\ \leq 0 & \text{if } x_m < 0. \end{cases}$$

If d is even then

$$\begin{aligned} 0 \leq \lim_{x_m \rightarrow \infty} Q(x_1, \dots, x_m) &= \lim_{x_m \rightarrow \infty} x_m^d q(x_1, \dots, x_{m-1}) \\ &= \lim_{x_m \rightarrow -\infty} x_m^d q(x_1, \dots, x_{m-1}) = \lim_{x_m \rightarrow -\infty} Q(x_1, \dots, x_m) \leq 0 \end{aligned}$$

so that $\lim_{x_m \rightarrow \infty} Q(x_1, \dots, x_m) = 0$, which contradicts that $q(x_1, \dots, x_{m-1}) \neq 0$.

Hence d is odd. Now, if $\phi = \phi' \wedge \phi_m$ choose $x_1, \dots, x_{m-1} < 0$, whilst if $\phi = \phi' \vee \phi_m$ choose $x_1, \dots, x_{m-1} \geq 0$ for which $q(x_1, \dots, x_{m-1}) \neq 0$. In the first case $\phi(x_1, \dots, x_m) = 0$ so that $Q(x_1, \dots, x_m) \leq 0$ for all x_m , whilst in the second case $\phi(x_1, \dots, x_m) = 1$ so that $Q(x_1, \dots, x_m) \geq 0$ for all x_m . But, since d is odd,

$$\lim_{x_m \rightarrow \infty} Q(x_1, \dots, x_m) = - \lim_{x_m \rightarrow -\infty} Q(x_1, \dots, x_m)$$

so that in both cases $\lim_{x_m \rightarrow \infty} Q(x_1, \dots, x_m) = 0$, which again contradicts that $q(x_1, \dots, x_{m-1}) \neq 0$. This completes the proof of Theorem 4.

References

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